

### Exercise 3.7.

Find a regular expression corresponding to each of the following subsets of  $\{a, b\}^*$ .

- a. The language of all strings containing exactly two  $a$ 's.
- c. The language of all strings that do not end with  $ab$ .
- e. The language of all strings not containing the substring  $aa$ .
- f. The language of all strings in which the number of  $a$ 's is even.
- g. The language of all strings containing no more than one occurrence of the string  $aa$ . (The string  $aaa$  should be viewed as containing two occurrences of  $aa$ .)

### Exercise 3.7.

Find a regular expression corresponding to each of the following subsets of  $\{a, b\}^*$ .

- i. The language of all strings containing both  $bb$  and  $aba$  as substrings.
- j. The language of all strings not containing the substring  $aaa$ .
- k. The language of all strings not containing the substring  $bba$ .
- l. The language of all strings containing both  $aba$  and  $bab$  as substrings.
- m. The language of all strings in which the number of  $a$ 's is even and the number of  $b$ 's is odd.

**Exercise 3.1.** In each case below, find a string of minimum length in  $\{a, b\}^*$  **not** in the language corresponding to the given regular expression.

**a.**  $b^*(ab)^*a^*$

**b.**  $(a^* + b^*)(a^* + b^*)(a^* + b^*)$

**Exercise 3.2.** Consider the two regular expressions

$$r = a^* + b^* \quad s = ab^* + ba^* + b^*a + (a^*b)^*$$

- a. Find a string corresponding to  $r$  but not to  $s$ .
- b. Find a string corresponding to  $s$  but not to  $r$ .
- c. Find a string corresponding to both  $r$  and  $s$ .
- d. Find a string in  $\{a, b\}^*$  corresponding to neither  $r$  nor  $s$ .

### Exercise 3.10.

- a.** If  $L$  is the language corresponding to the regular expression  $(aab + bbaba)^*baba$ , find a regular expression corresponding to  $L^r = \{x^r \mid x \in L\}$ .
- b.** Using the example in part (a) as a model, give a recursive definition (based on Definition 3.1) of the reverse  $e^r$  of a regular expression  $e$ .
- c.** Show that for every regular expression  $e$ , if the language  $L$  corresponds to  $e$ , then  $L^r$  corresponds to  $e^r$ .

**Exercise 3.41.** For each of the following regular expressions, draw an NFA accepting the corresponding language, so that there is a recognizable correspondence between the regular expression and the transition diagram.

e.  $(a^*bb)^* + bb^*a^*$

**Exercise 3.42.** For part (e) of Exercise 3.41, draw the NFA that is obtained by a literal application of Kleene's theorem, without any simplifications.

**Exercise 3.51.** Use the algorithm of **Brzozowski and McCluskey** to find a regular expression corresponding to the FA below.

**a.**

