



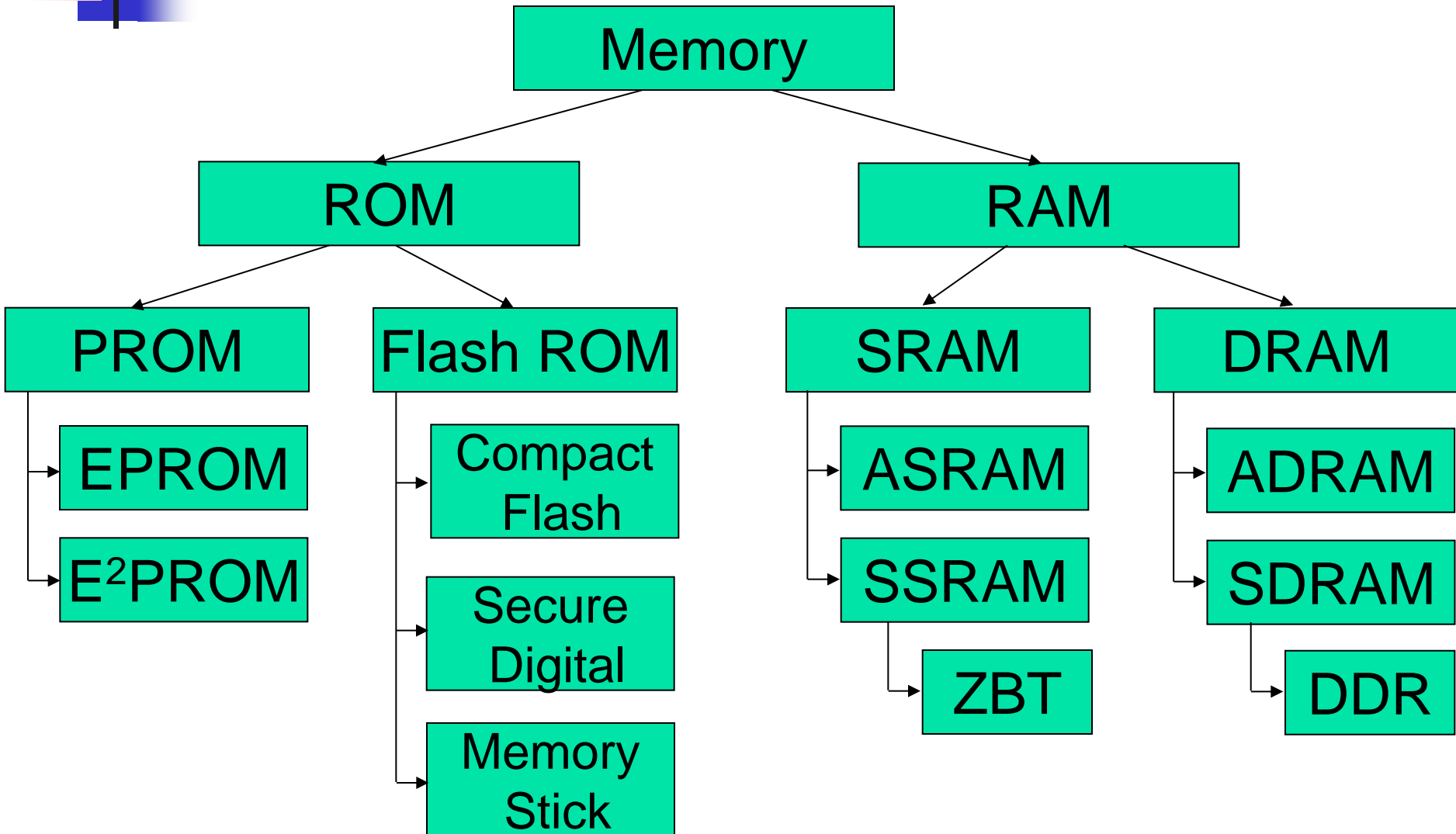
Memories



Overview

- Memory Classification
- Read-Only Memory (ROM)
 - Types of ROM
 - PROM, EPROM, E²PROM
 - Flash ROMs (Compact Flash, Secure Digital, Memory Stick)
- Random Access Memory (RAM)
 - Types of RAM
 - Static RAM (SRAM) - ASRAM, SSRAM, ZBT
 - Dynamic RAM (DRAM) - SDRAM, DDR RAM, RDRAM
- Functional Behavior of RAM
 - Introduction, Block Diagram, Memory Size
 - Reading/Writing from/to RAM
- Implementing Static RAM
 - Making Larger and Wider Memory from Smaller Memories
- Functional Behavior of ROM
 - Memories and Boolean Functions

Types of Memories





Read – Only Memory (ROM)

- A read-only memory, or ROM, is a special kind of memory whose contents cannot be easily modified.
 - Writing (storing) data into a ROM chip is a relatively slow process.
 - Data is stored onto a ROM chip using special hardware tools.
- ROMs can store data even without power!
- ROMs are useful for holding data that “almost never” changes.
 - Arithmetic circuits might use tables to speed up computations of logarithms or divisions.
 - Many computers use a ROM to store important programs that should not be modified, such as the system BIOS.
 - PDAs, game machines, cell phones, vending machines and other electronic devices may also contain non-modifiable programs.



Types of ROMs

- PROM (Programmable ROM)
 - It can be programmed (written) only once.
- EPROM (Electrical Programmable ROM)
 - It can be programmed a limited number of times.
 - Before reprogramming the memory you must erase it with ultraviolet waves.
- E²PROM (Electrical Erasable Programmable ROM)
 - The same as EPROM but
 - You can erase it by applying electrical pulses on special pins.
- Programming (writing) and erasing the above memories is a very slow process. It may take seconds or minutes!

Types of ROMs (Flash ROMs)

- Some newer types of ROMs do allow for easier writing, although the speeds still do not compare with regular RAM memories.
 - MP3 players, digital cameras and other toys use CompactFlash, Secure Digital, or MemoryStick cards for non-volatile storage.
 - Many devices allow you to upgrade programs stored in “flash ROM.





Random Access Memory (RAM)

- RAM is a kind of memory whose contents can be easily modified.
 - Writing (storing) data into a RAM chip is as fast as reading data.
 - No special hardware tools are needed to store the data in the RAM.
- RAMs cannot store data without power!
- RAMs are useful for storing temporary data that has to be modified very often.
 - Many computers use RAM for storing the currently executed program and/or intermediate data.
 - The CPU cache memory is a RAM memory.



Types of RAMs

- SRAM (Static RAM)
 - It is build using Latches
 - Expensive in terms of hardware
 - For each bit of storage you need 6 transistors
 - Very fast memory
- DRAM (Dynamic RAM)
 - **Dynamic memory** is built with capacitors.
 - A stored charge on the capacitor represents a logical 1.
 - No charge represents a logic 0.
 - However, capacitors lose their charge after a few milliseconds. The memory requires constant **refreshing** to recharge the capacitors. (That's what's “dynamic” about it.)
 - Dynamic RAMs tend to be physically smaller than static RAMs.
 - A single bit of data can be stored with just one capacitor and one transistor, while static RAM cells typically require 4-6 transistors.
 - This means dynamic RAM is cheaper and denser—more bits can be stored in the same physical area.



Dynamic vs. Static Memory

- In practice, dynamic RAM is used for a computer's main memory, since it is cheap and you can pack a lot of storage into a small space
 - These days you can buy 1 GB of memory for as little as \$15
 - You can have a system with 8 GB or more of memory (up to 128GB)
- The disadvantage of dynamic RAM is its speed
 - Up to 1.6GHz, which can be slower than the processor itself
 - You also have to consider **latency**, or the time it takes data to travel from RAM to the processor
- Real systems augment dynamic memory with small but fast sections of static memory called **caches**.
 - Typical processor caches range in size from 128KB to 4MB
 - That is small compared to a 8 GB main memory, but it is enough to significantly increase a computer's overall speed



Types of SRAM

- ASRAM (Asynchronous Static RAM)
 - no clock signal is used when reading/writing data
- SSRAM (Synchronous Static RAM)
 - read and write from/to the memory is synchronized by a clock signal
 - ZBT (Zero-Bus Turnaround) SSRAM
 - Very fast memory
 - See more details on Internet



Types of DRAM

- ADRAM (Asynchronous Dynamic RAM)
 - older type of DRAM used in the early personal computers
 - memory access is not synchronized with the system clock
- SDRAM (Synchronous Dynamic RAM)
 - widely used nowadays
 - responds to read/write operations in synchrony with the signal of the system clock
 - DDR SDRAM (Double Data Rate SDRAM)

SDRAM

- **Synchronous DRAM**, or **SDRAM**, is a very common type of PC memory
- Memory chips are organized into “modules” that are connected to the CPU via a 64-bit (8-byte) bus
- The bus speeds are rated in megahertz: PC66, PC100 and PC133 memory bus run at 66MHz, 100MHz and 133MHz, respectively
- The memory **bandwidth** can be computed by multiplying the number of transfers per second (T/s) by the size of each transfer
 - PC100 bus works at 100MHz => 100MT/s
 - PC100 can transfer up to 800MB per second (100MT/s x 8 Bytes/Transfer)





DDR SDRAM

- A newer type of memory is **Double Data Rate**, or **DDR-SDRAM**
 - Several generations from DDR2 to DDR4
- It is very similar to regular SDRAM, except data can be transferred on both the positive *and* negative clock edges
- DDR4 memory bus runs at 800-1600MHz but the effective data transfer rate is doubled, i.e., 1600-3200MT/s
- These memories are called PC4-12800 and PC4-25600 with bandwidth
 - Up to 12800MB/s = 1600MT/s x 8 Bytes/Transfer
 - Up to 25600MB/s = 3200MT/s x 8 Bytes/Transfer
- DDR-SDRAM has lower power consumption, using 1.2-2.5V instead of 3.3V like SDRAM. This makes it good for notebooks and other mobile devices



ROMs vs. RAMs

- There are some important differences between ROM and RAM.
 - ROMs are “non-volatile”
 - data is preserved even without power
 - RAM contents disappear once power is lost.
 - ROMs require special (and slower) techniques for writing, so they are considered to be “read-only” devices.



Random Access Memory (RAM)

- Sequential circuits all depend upon the presence of memory.
 - A flip-flop can store one bit of information.
 - A register can store a single “word,” typically 32-64 bits.
- **Random access memory**, or **RAM**, allows us to store even larger amounts of data. Today you will see:
 - The basic interface to RAM memory.
 - How you can implement static RAM chips hierarchically.



Introduction to RAM

- **Random-access memory**, or **RAM**, provides large quantities of temporary storage in a computer system.
- Remember the basic properties of a memory:
 - It should be able to store a value.
 - You should be able to read the value that was stored.
 - You should be able to change the stored value.
- A RAM is similar, except that it can store *many* values.
 - An **address** will specify which memory value we are interested in.
 - Each value can be a multiple-bit **word** (e.g., 32 bits).
- We will refine the memory properties as follows:

A RAM should be able to:

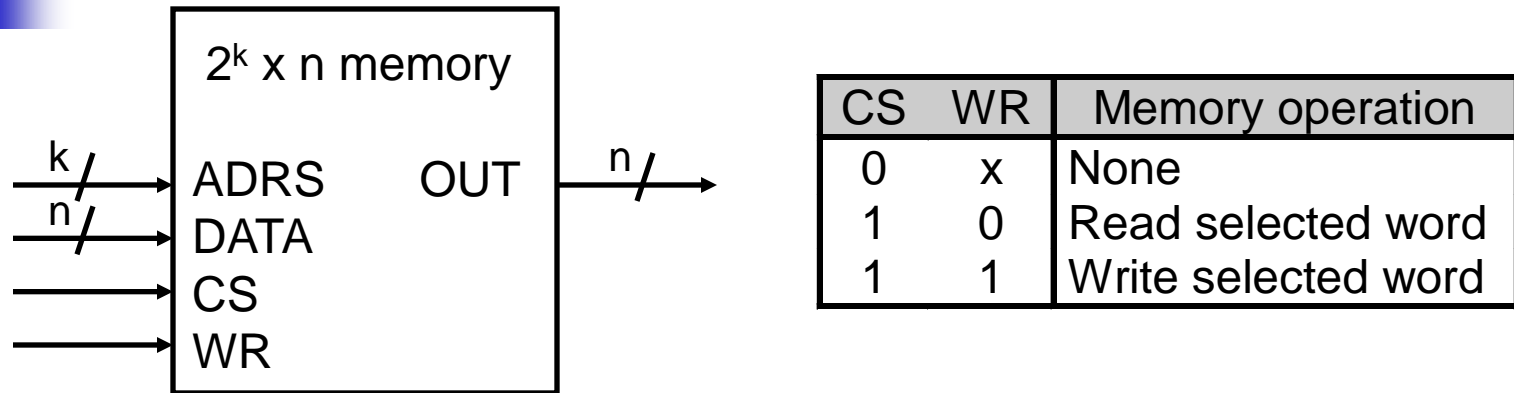
- Store many words (values), one per address
- Read the word that was stored at a particular address
- Change the word that was stored at a particular address

Picture of Memory

- You can think of computer memory as being one big array of data.
 - The address serves as an array index.
 - Each address refers to one word of data.
- You can read or modify the data at any given memory address, just like you can read or modify the contents of an array at any given index.
- If you have worked with pointers in C or C++, then you have already worked with memory addresses.

Address	Data
00000000	
00000001	
00000002	
.	
.	
.	
.	
.	
.	
.	
.	
.	
.	
FFFFFFFFD	
FFFFFFFE	
FFFFFFF	

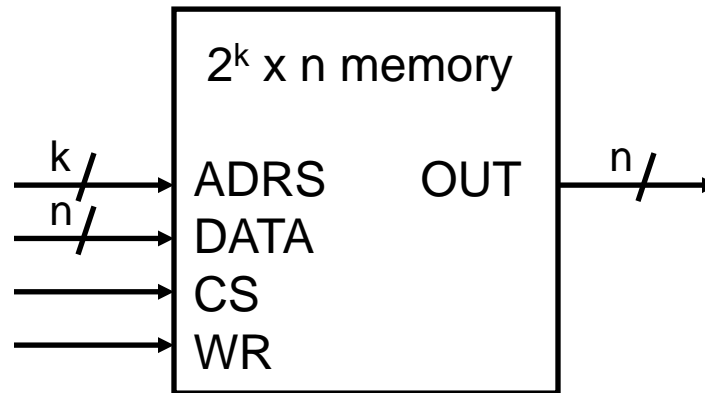
Block Diagram of RAM



- This block diagram introduces the main interface to RAM.
 - A Chip Select, **CS**, enables or disables the RAM.
 - **ADDRESS** specifies the address or location to read from or write to.
 - **WR** selects between reading from or writing to the memory.
 - To read from memory, WR should be set to 0.
OUT will be the n-bit value stored at ADDRESS.
 - To write to memory, we set $WR = 1$.
DATA is the n-bit value to save in memory.
- This interface makes it easy to combine RAMs together, as you will see.

Memory Sizes

- We refer to this as a $2^k \times n$ memory.
 - There are k *address lines*, which can specify one of 2^k addresses.
 - Each address contains an n -bit word.



- For example, a $2^{24} \times 16$ RAM contains $2^{24} = 16\text{M}$ words, each 16 bits long.
 - The total *storage capacity* is $2^{24} \times 16 = 2^{28}$ bits.

Size Matters!

- Memory sizes are usually specified in numbers of **bytes** (8 bits).
- The 2^{28} -bit memory on the previous page translates into:

$$2^{28} \text{ bits} / 8 \text{ bits per byte} = 2^{25} \text{ bytes}$$

- With the abbreviations below, this is equivalent to 32 megabytes.

	Prefix	Base 2	Base 10
K	Kilo	$2^{10} = 1,024$	$10^3 = 1,000$
M	Mega	$2^{20} = 1,048,576$	$10^6 = 1,000,000$
G	Giga	$2^{30} = 1,073,741,824$	$10^9 = 1,000,000,000$

- To confuse you, RAM size is measured in base 2 units, while hard drive size is measured in base 10 units.
 - In this class, we'll only concern ourselves with the base 2 units.

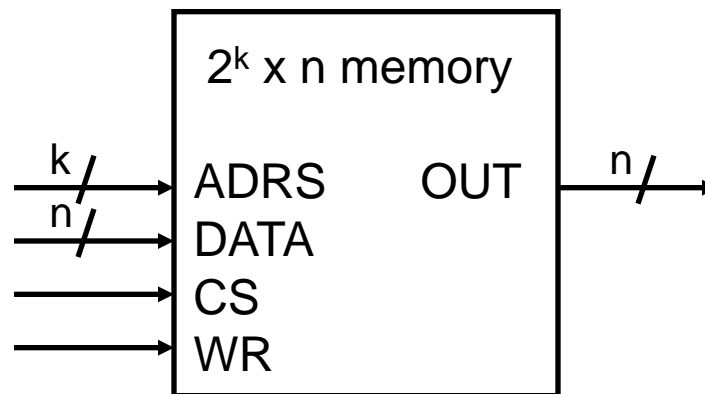
Typical Memory Sizes

- Some typical memory capacities:
 - PCs usually come with 4-8 GB RAM.
 - Digital cameras and MP3 players can have 4 GB or more of storage.
- Many operating systems implement **virtual memory**, which makes the memory seem larger than it really is.
 - Most systems allow up to 32-bit addresses. This works out to 2^{32} , or about four billion, different possible addresses.
 - With a data size of one byte, the result is apparently a 4GB memory!
 - The operating system uses hard disk space as a substitute for “real” memory.

Address	Data
00000000	
00000001	
00000002	
.	
.	
.	
.	
.	
.	
.	
.	
.	
.	
.	
FFFFFFFFD	
FFFFFFFE	
FFFFFFF	

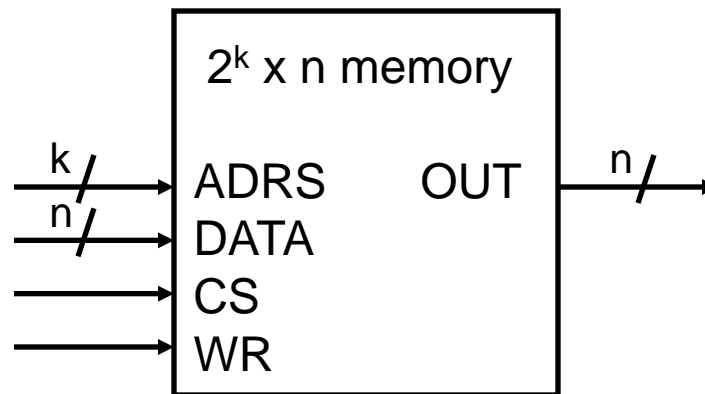
Reading RAM

- To *read* from this RAM, the controlling circuit must:
 - Enable the chip by ensuring $CS = 1$.
 - Send the desired address to the ADRS input.
 - Select the read operation, by setting $WR = 0$.
 - The contents of that address appear on OUT after a little while.
- Notice that the DATA input is unused for read operations.
- Each memory has a specific timing diagram that specifies the correct time sequence of the events described above!
 - The timing diagram depends on the memory implementation.
 - It is given by the producer of the memory chip.



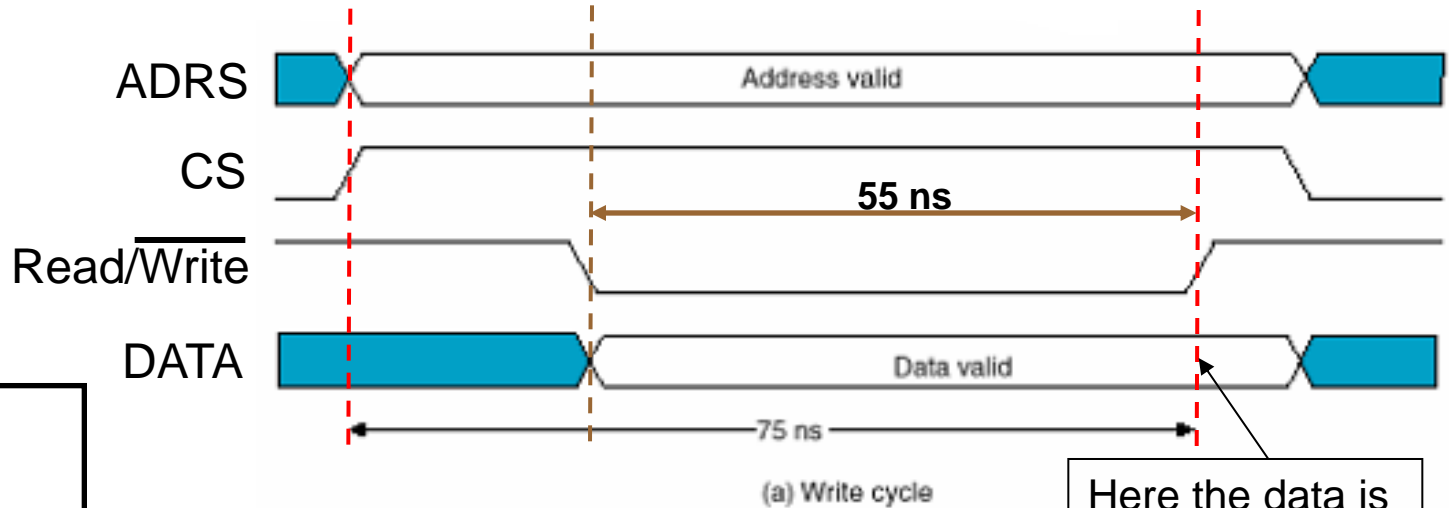
Writing RAM

- To *write* to this RAM, you need to:
 - Enable the chip by setting $CS = 1$.
 - Send the desired address to the ADRS input.
 - Send the word to store to the DATA input.
 - Select the write operation, by setting $WR = 1$.
- The output OUT is not needed for memory write operations.
- Again, each memory has a specific timing diagram that specifies the correct time sequence of the events described above!



Example of Read/Write Timing Diagrams

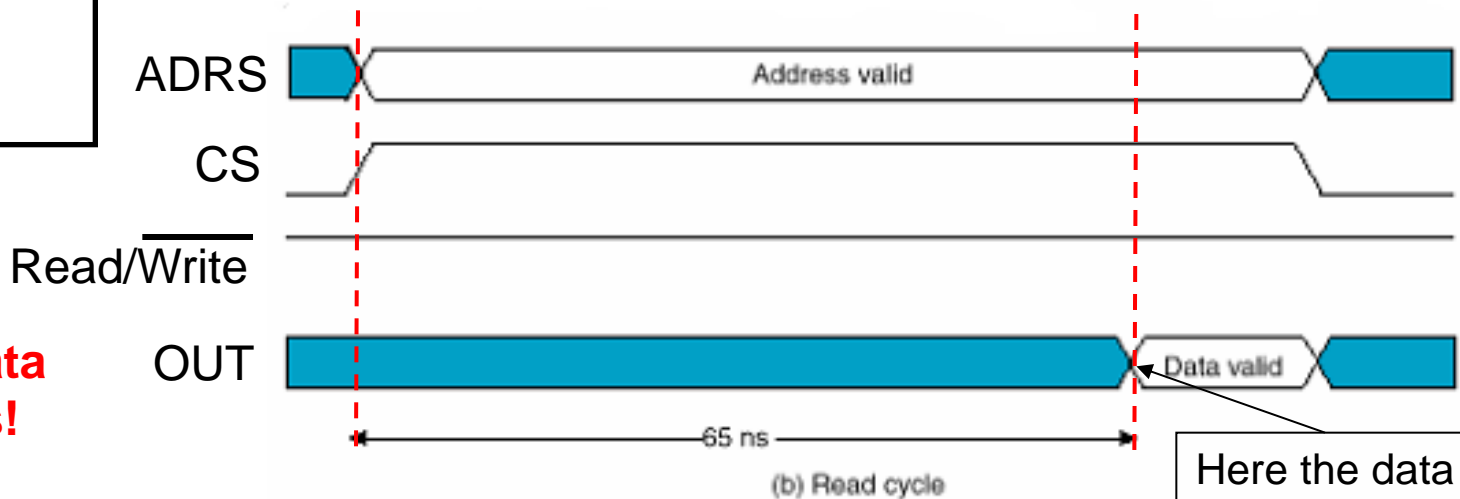
Writing data takes 75 ns!



(a) Write cycle

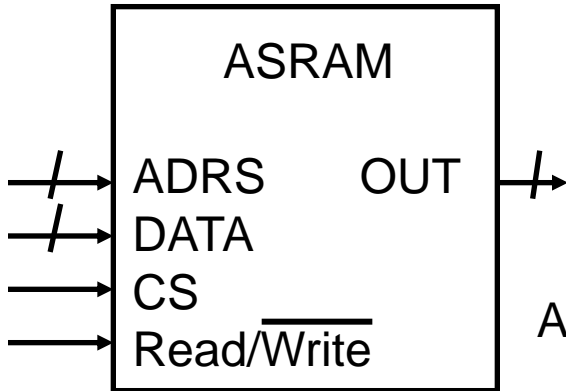
Here the data is stored in RAM

Reading data takes 65 ns!



(b) Read cycle

Here the data can be read



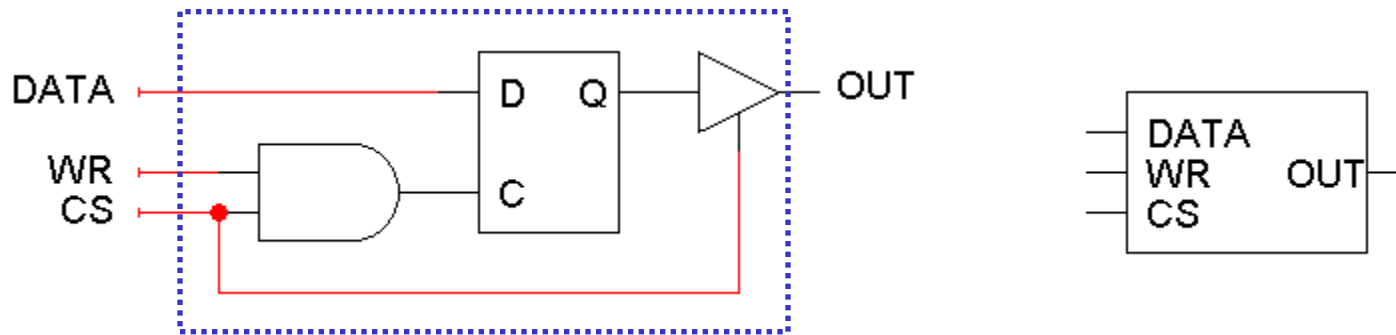


Static Memory

- How can we implement the memory chip?
- There are many different kinds of RAM.
 - We will discuss only **static memory**, which is most commonly used in caches and video cards.
- Static memory is modeled using one *latch* for each bit of storage.
- Why use latches instead of flip-flops?
 - A latch can be made with only two NAND or two NOR gates, but a flip-flop requires at least twice that much hardware.
 - In general, smaller is faster, cheaper and requires less power.
 - The tradeoff is that getting the timing exactly right is a pain.

Starting with Latches

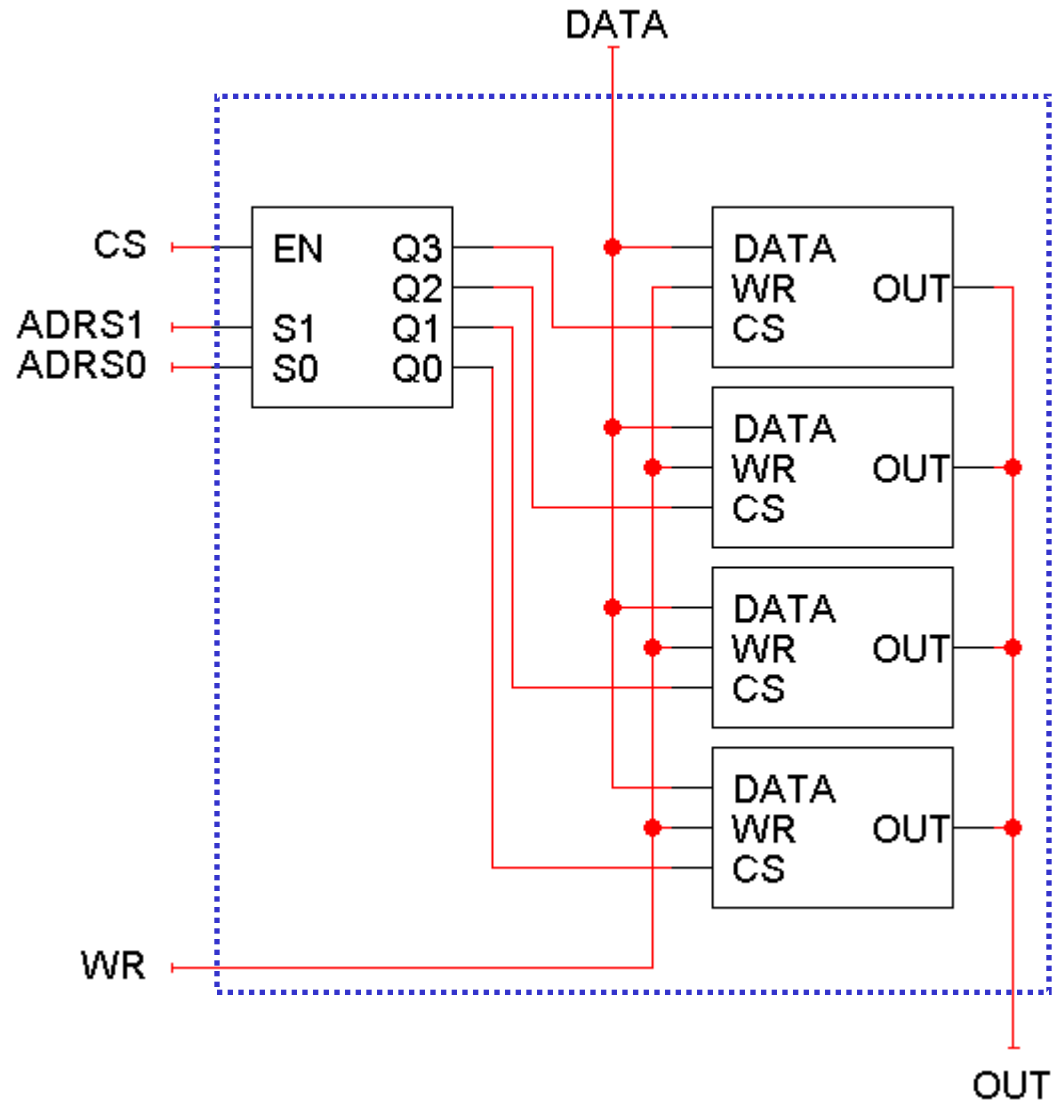
- To start, we can use one latch to store each bit. A one-bit **RAM cell** is shown here.



- Since this is just a one-bit memory, an ADDRESS input is not needed.
- Writing to the RAM cell:
 - When $CS = 1$ and $WR = 1$, the latch control input will be 1.
 - The DATA input is thus saved in the D latch.
- Reading from the RAM cell and maintaining the current contents:
 - When $CS = 0$ or when $WR = 0$, the latch control input is 0, so the latch just maintains its present state.
 - When $CS = 1$ and when $WR = 0$, the latch control input is also 0 and the current latch contents will appear on OUT.

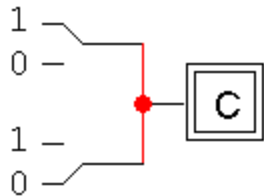
My First RAM

- We can use these cells to make a 4 x 1 RAM.
- Since there are four words, ADRS is two bits.
- Each word is only one bit, so DATA and OUT are one bit each.
- Word selection is done with a decoder attached to the CS inputs of the RAM cells. Only one cell can be read or written at a time.
- Notice that the outputs are connected together with a *single line*!



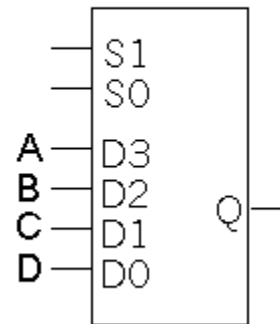
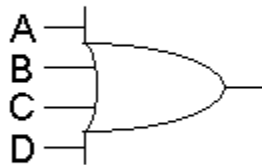
Connecting Outputs Together

- In normal practice, it is bad to connect outputs together. If the outputs have different values, then a conflict arises (**short circuit**).



The "C" means "conflict",
i.e., short circuit

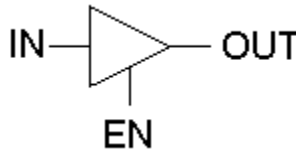
- The standard way to "combine" outputs is to use OR gates or MUXs.



- This can get expensive, with many wires and gates with large Fan-Ins.

Those Funny Triangles

- The triangle represents a **three-state buffer**.
- Unlike regular logic gates, the output can be one of *three* different possibilities, as shown in the table.

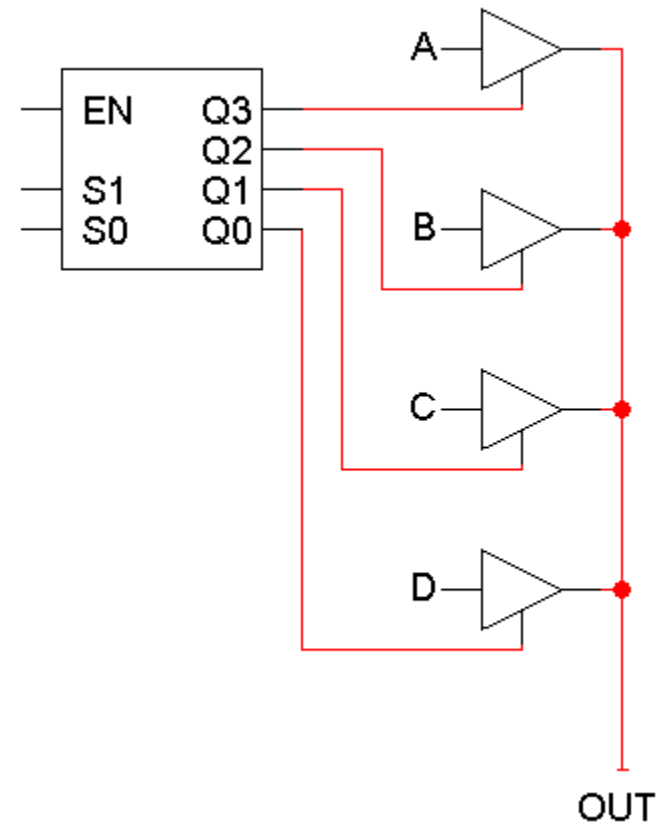


EN	IN	OUT
0	x	Disconnected
1	0	0
1	1	1

- “**Disconnected**” means no output appears at all, in which case it is safe to connect OUT to another output signal.
- The disconnected value is also sometimes called **high impedance** or **Hi-Z**.

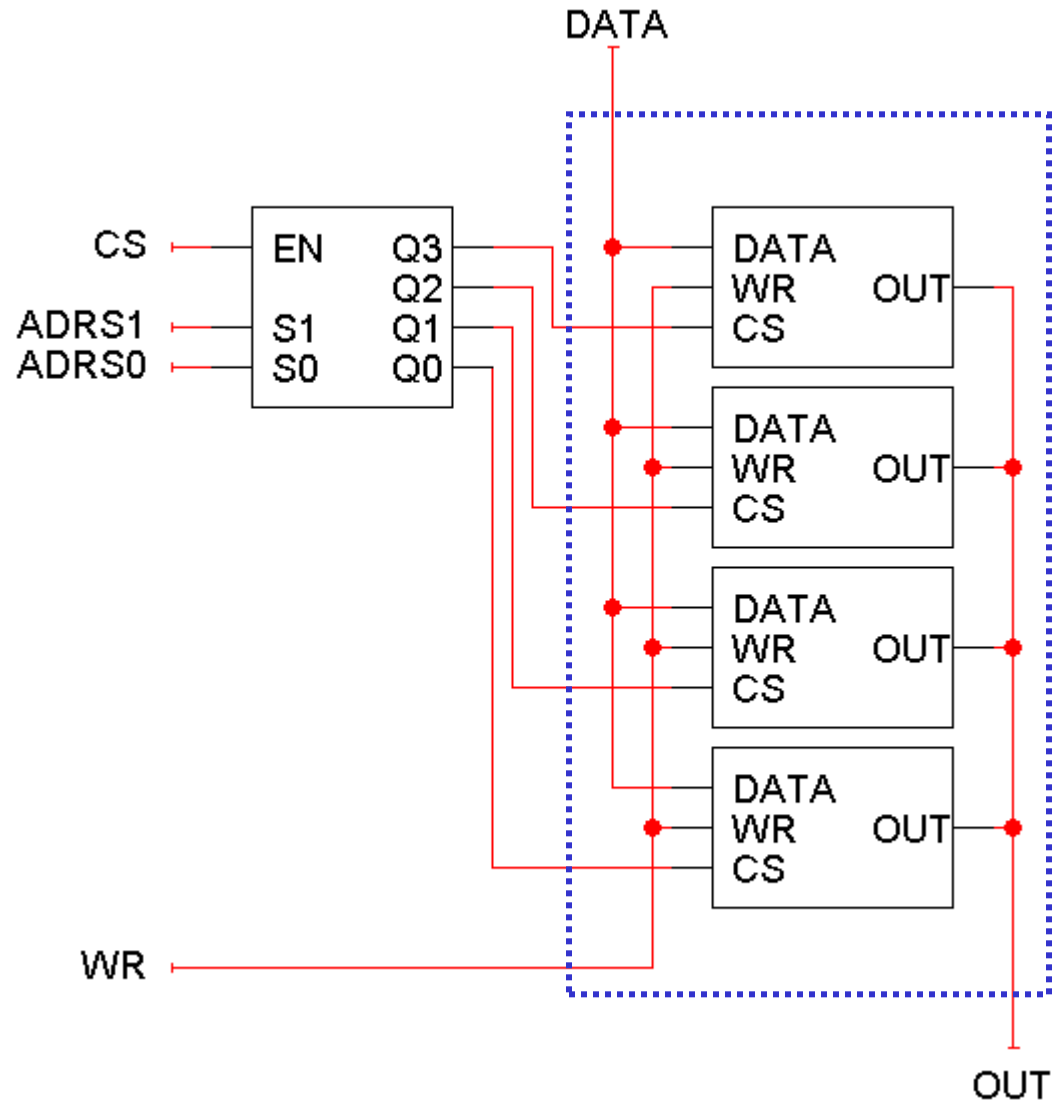
Connecting Three-State Buffs Together

- You can connect several three-state buffer outputs together if you can *guarantee* that only one of them is enabled at any time.
- The easiest way to do this is to use a decoder!
- If the decoder is disabled, then all the three-state buffers will appear to be disconnected, and OUT will also appear disconnected.
- If the decoder is enabled, then exactly one of its outputs will be true, so only one of the tri-state buffers will be connected and produce an output.
- The net result is we can save some wire and gate costs. We also get a little more flexibility in putting circuits together.



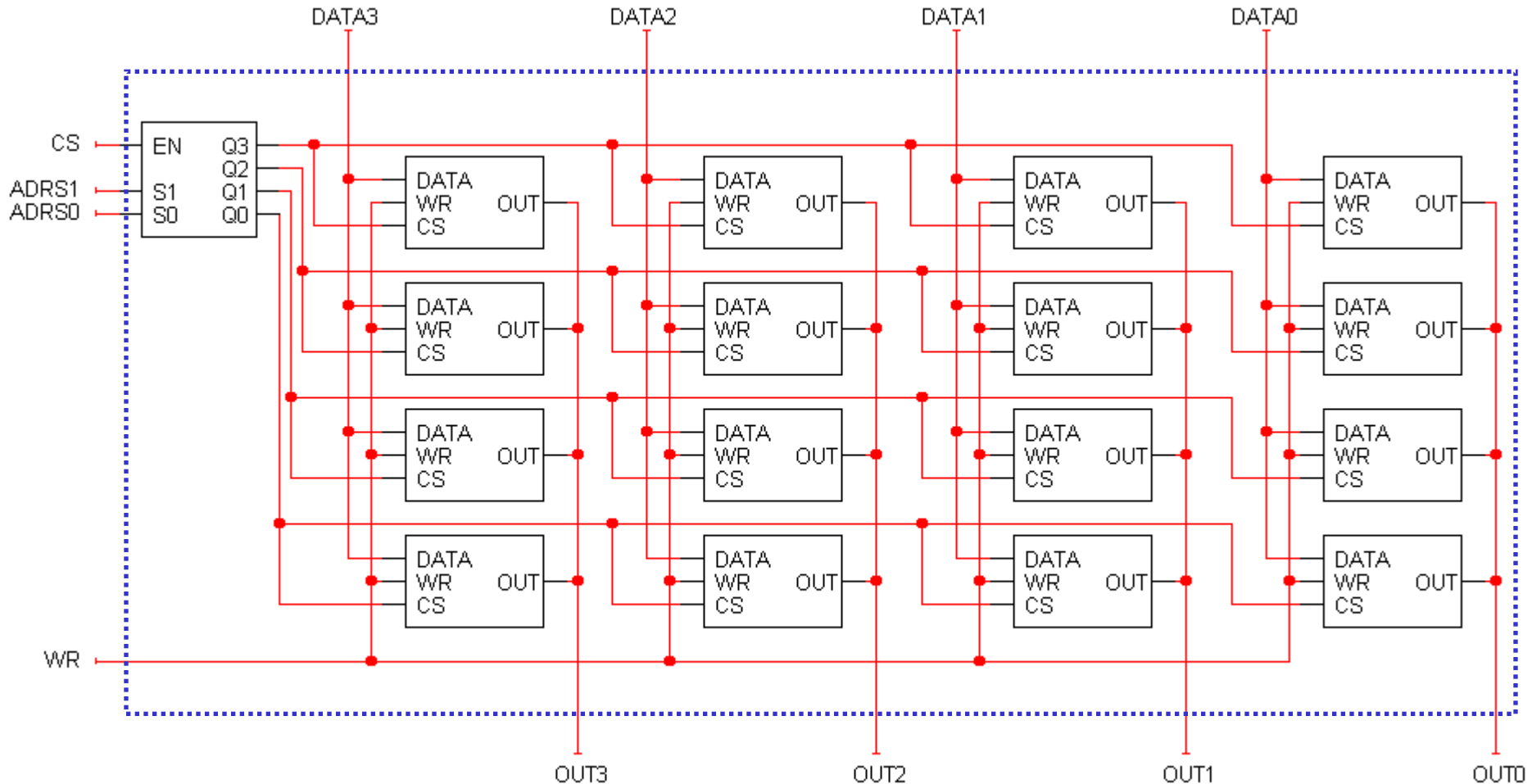
Bigger and Better

- Here is the 4 x 1 RAM once again.
- How can we make a “wider” memory with more bits per word, like maybe a 4 x 4 RAM?
- Duplicate the stuff in the dashed box!



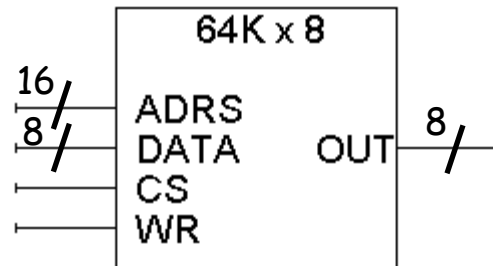
A 4 x 4 RAM

- DATA and OUT are now each *four* bits long, so you can read and write four-bit words.



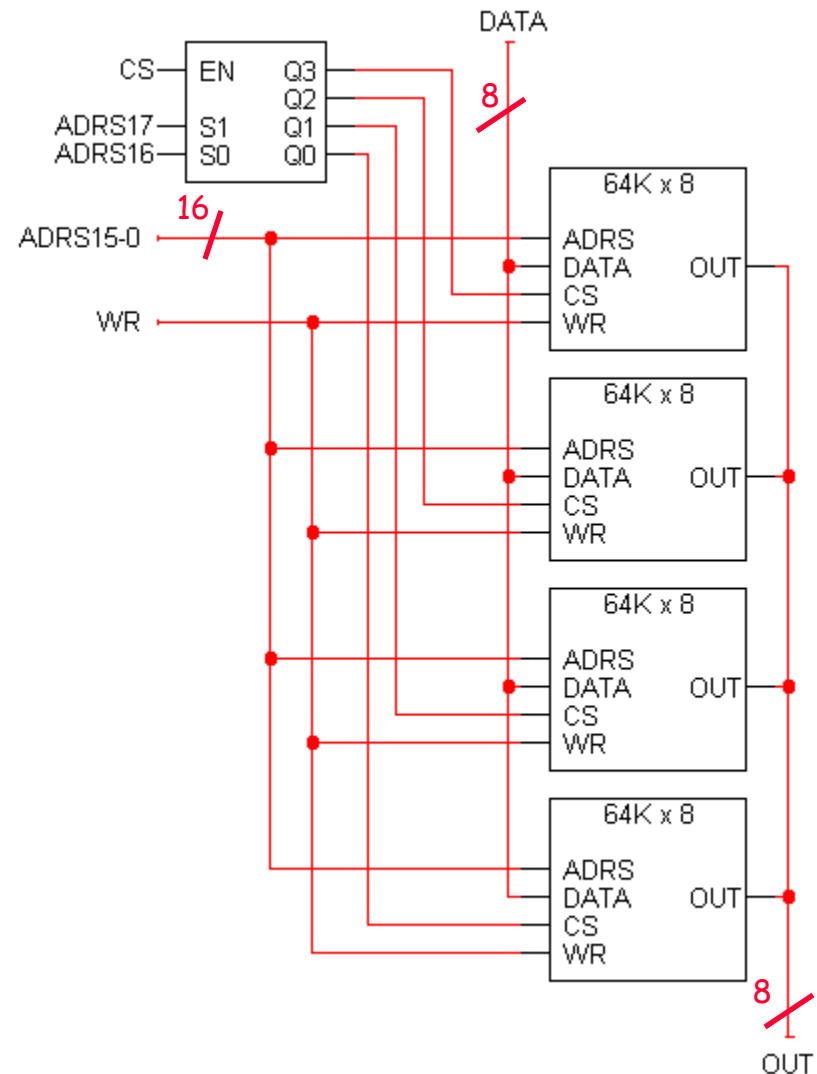
Larger RAMs from Smaller RAMs

- We can use small RAMs as building blocks for making larger memories, by following the same principles as in the previous examples.
- As an example, suppose we have some 64K x 8 RAMs to start with:
 - $64K = 2^6 \times 2^{10} = 2^{16}$, so there are 16 address lines.
 - There are 8 data lines.



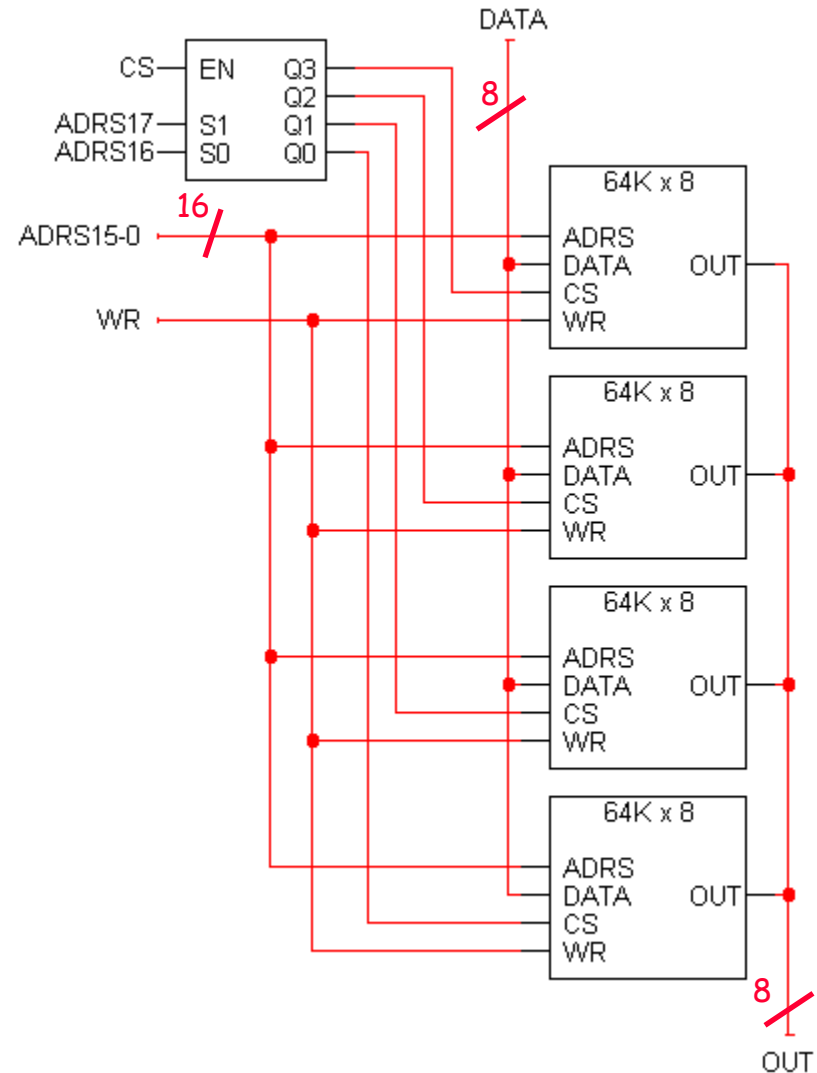
Making a Larger Memory

- We can put four 64K x 8 chips together to make a 256K x 8 memory.
- For 256K words, we need 18 address lines.
 - The two most significant address lines go to the decoder, which selects one of the four 64K x 8 RAM chips.
 - The other 16 address lines are shared by the 64K x 8 chips.
- The 64K x 8 chips also share WR and DATA inputs.
- This assumes the 64K x 8 chips have three-state outputs.

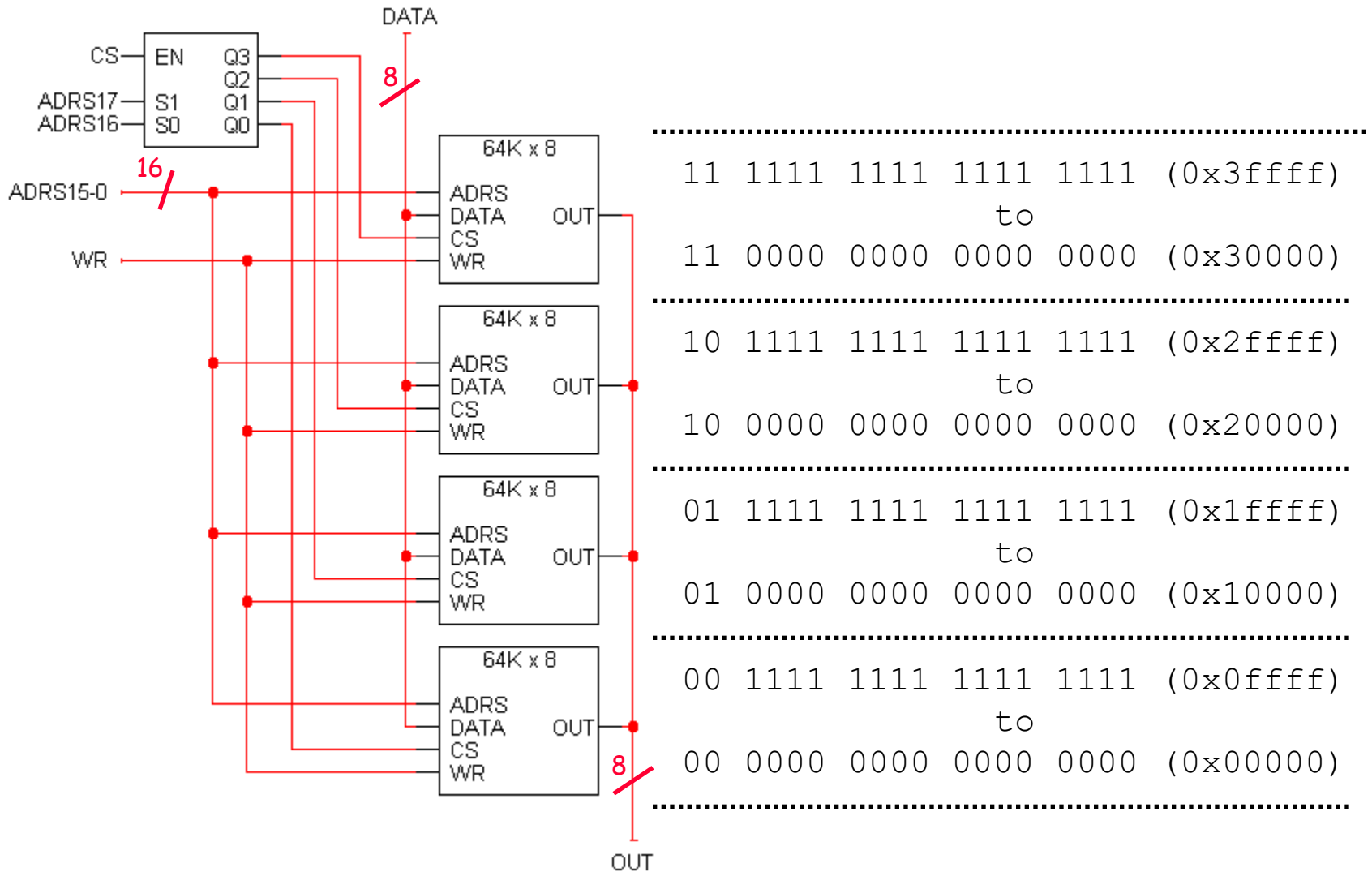


Analyzing the 256K x 8 RAM

- There are 256K words of memory, spread out among the four smaller 64K x 8 RAM chips.
- When the two most significant bits of the address are 00, the bottom RAM chip is selected. It holds data for the first 64K addresses.
- The next chip up is enabled when the address starts with 01. It holds data for the second 64K addresses.
- The third chip up holds data for the next 64K addresses.
- The final chip contains the data of the final 64K addresses.

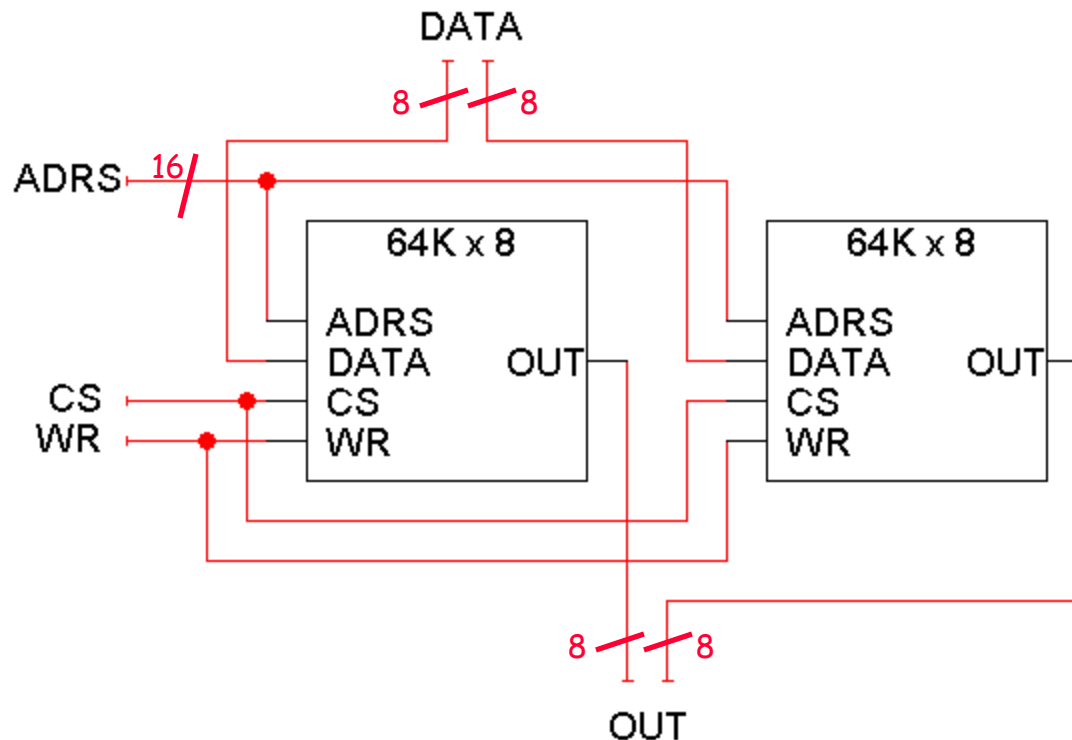


Address Ranges



Making a Wider Memory

- You can also combine smaller chips to make wider memories, with the same number of addresses but more bits per word.
- Here is a 64K x 16 RAM, created from two 64K x 8 chips.
 - The left chip contains the most significant 8 bits of the data.
 - The right chip contains the lower 8 bits of the data.

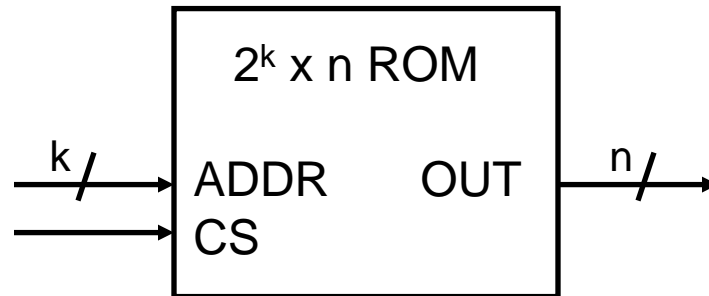




Summary

- A RAM looks like a bunch of registers connected together, allowing users to select a particular address to read or write.
- Much of the hardware in memory chips supports this selection process:
 - Chip select inputs
 - Decoders
 - Tri-state buffers
- By providing a general interface, it is easy to connect RAMs together to make “longer” and “wider” memories.

Read-Only Memory



- A read-only memory, or ROM, is a special kind of memory whose contents cannot be easily modified
 - The WR and DATA inputs that we saw in RAMs are not needed.
 - Data is stored onto a ROM chip using special hardware tools.
- ROMs are useful for holding data that almost never changes
 - Arithmetic circuits might use tables to speed up computations of logarithms or divisions.
 - Many computers use a ROM to store important programs that should not be modified, such as the system BIOS.

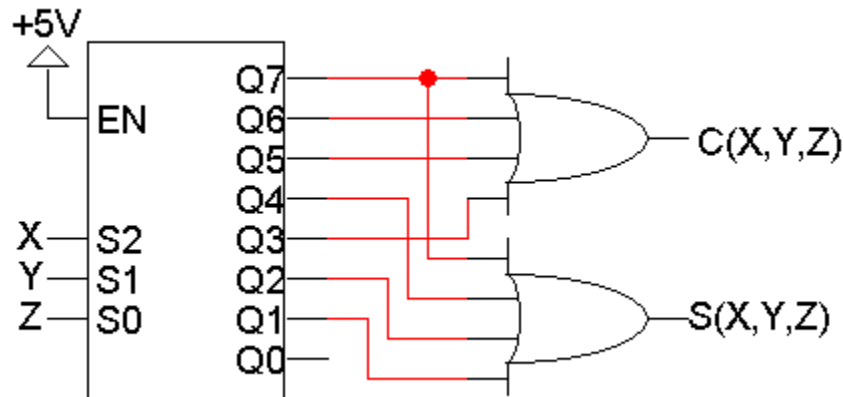
Memories and Functions

- ROMs are actually combinational devices, not sequential ones!
 - You can think of a ROM as a combinational circuit that takes an address as input, and produces some data as the output.
- A **ROM table** is basically just a truth table.
 - The table shows what data is stored at each ROM address.
 - You can generate that data combinationaly, using the address as the input.

Address $A_2A_1A_0$	Data $V_2V_1V_0$
000	000
001	100
010	110
011	100
100	101
101	000
110	011
111	011

Decoders

- We can already convert truth tables to circuits easily, with decoders.

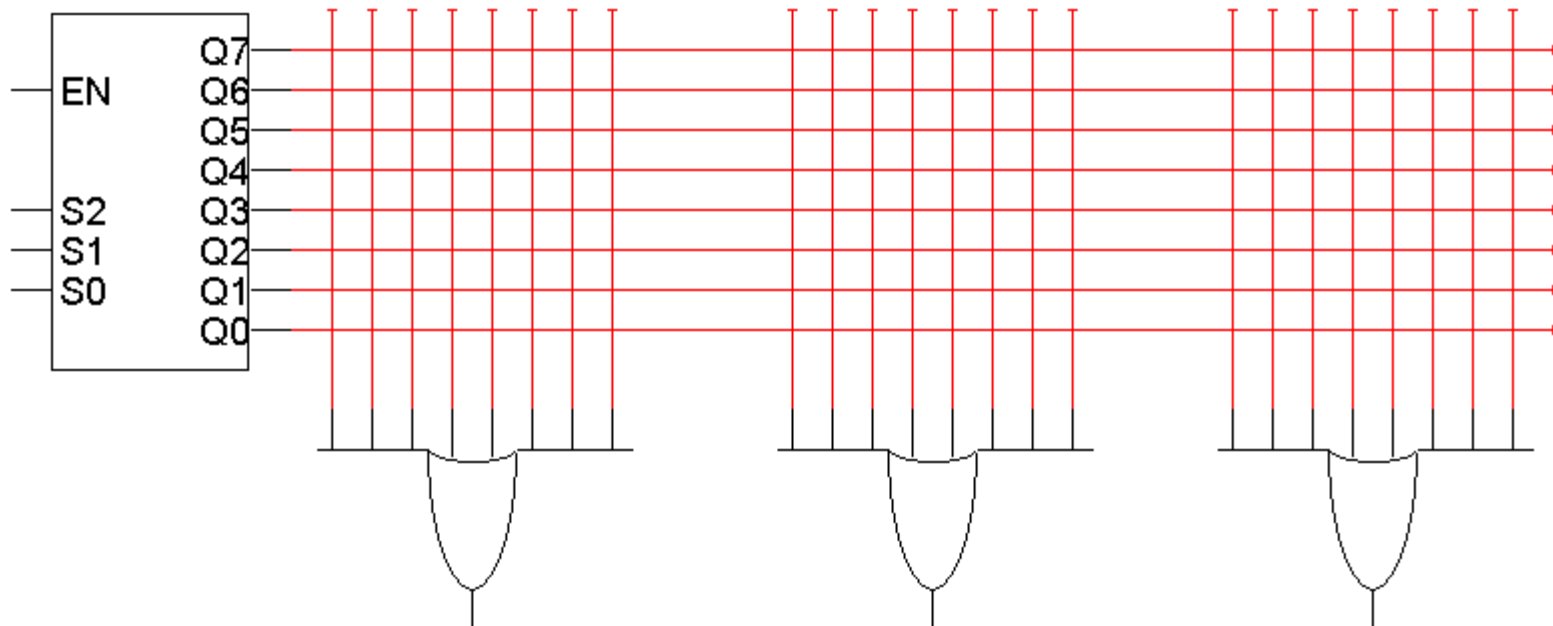


X	Y	Z	C	S
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1

- For example, you can think of this Full-Adder circuit as a memory that “stores” the sum and carry outputs from the truth table on the right.

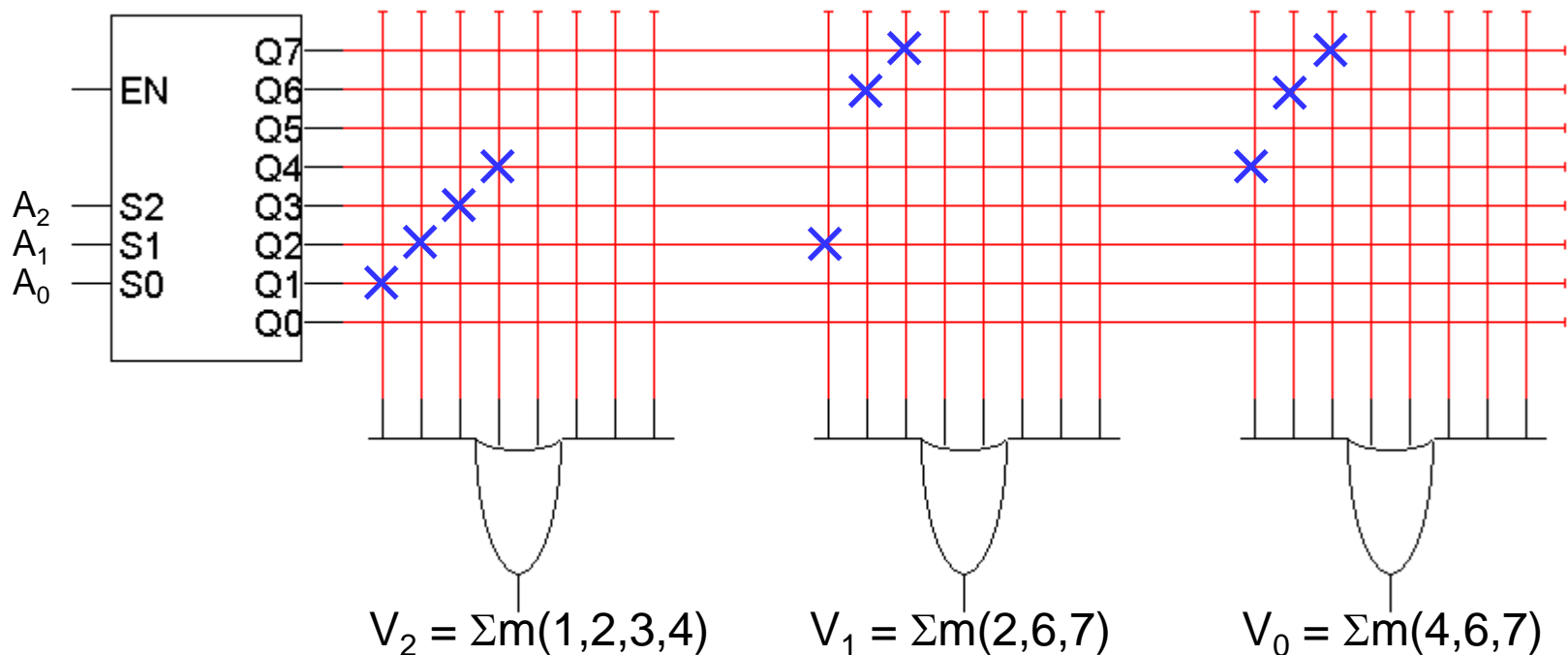
ROM Setup

- ROMs are based on this decoder implementation of functions.
 - A blank ROM just provides a decoder and several OR gates.
 - The connections between the decoder and the OR gates are “programmable,” so different functions can be implemented.
- To program a ROM, you just make the desired connections between the decoder outputs and the OR gate inputs.



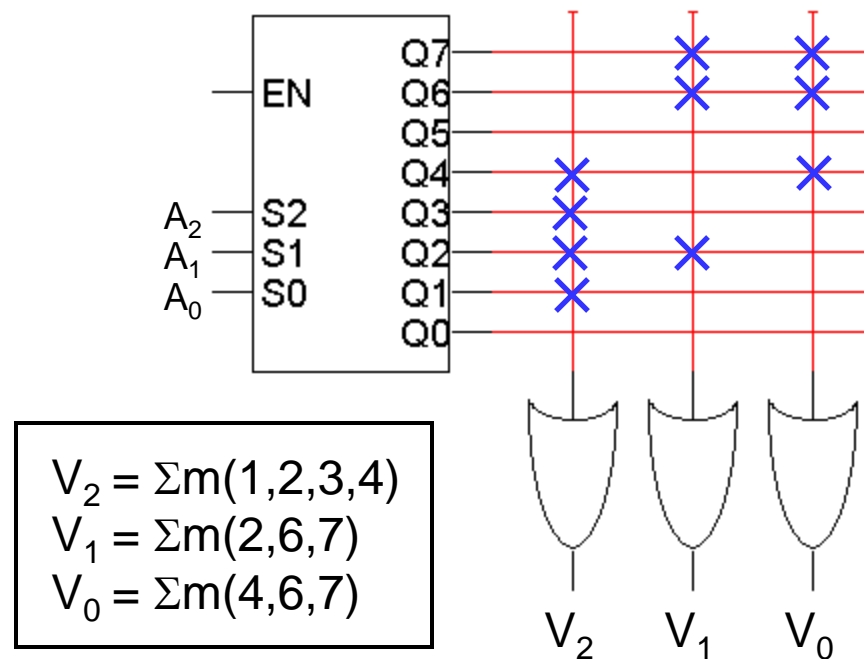
ROM Example

- Here are three functions, $V_2V_1V_0$, implemented with an **8 x 3 ROM**.
- Blue crosses (X) indicate connections between decoder outputs and OR gates. Otherwise there is no connection.



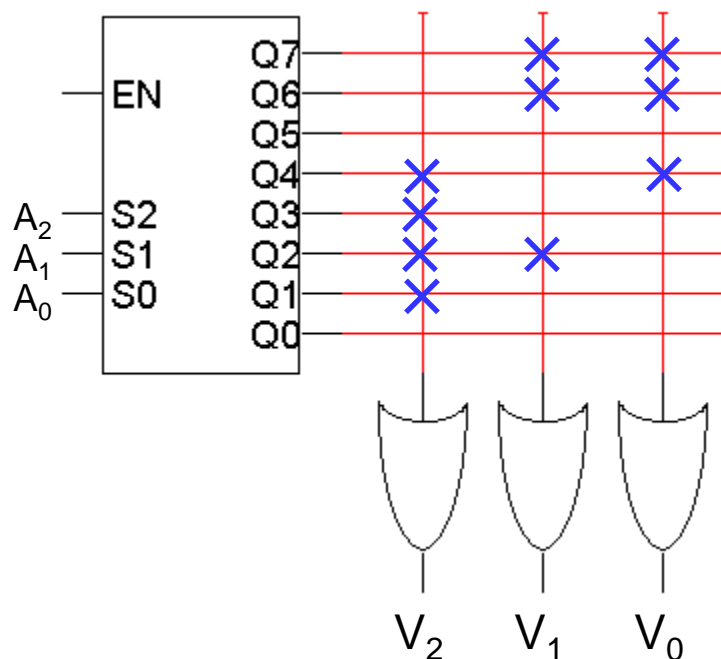
The Same Example Again

- Here is an alternative presentation of the same 8 x 3 ROM, using “abbreviated” OR gates to make the diagram neater.



Why Is This a “Memory”?

- This combinational circuit can be considered a read-only memory.
 - It stores eight words of data, each consisting of three bits.
 - The decoder inputs form an **address**, which refers to one of the eight available words.
 - So, every input combination corresponds to an address, which is “read” to produce a 3-bit data output.



Address $A_2A_1A_0$	Data $V_2V_1V_0$
000	000
001	100
010	110
011	100
100	101
101	000
110	011
111	011



Functions and Memories

- ROMs give us another way to implement functions.
- The idea behind using a ROM to implement a function is to “store” the function’s truth table, so we do not have to do any (well, very little) computation.
- This is like “memorization” or “caching” techniques in programming.



Summary

- We discussed RAM and ROM memories.
- There are two main kinds of RAM memory.
 - Static RAM
 - costs more in terms of HW, but the memory is faster
 - often used to implement cache memories
 - Dynamic RAM
 - costs less HW and requires less physical space
 - making it ideal for larger-capacity memories
 - access times are slower
- ROMs are programmable devices that can implement arbitrary functions, which is equivalent to acting as a read-only memory.