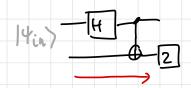


## SOLUTIONS TO

## TAKE HOME ASSIGNMENTS

PROBLEM 1.1.



(a) Matrix: (NOZ) CNOT (HOM) (Kn) NOTE ORDER.

In block FORM :

 $U = \frac{1}{5} \begin{pmatrix} 2 & 0 \\ 0 & 2 \end{pmatrix} \begin{pmatrix} 1 & 0 \\ 0 & \times \end{pmatrix} \begin{pmatrix} T & 1 \\ 1 & -1 \end{pmatrix} = \frac{1}{5} \begin{pmatrix} 2 & 0 \\ 0 & 2X \end{pmatrix} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} = \frac{1}{5} \begin{pmatrix} 2 & 2 \\ 2X & -2X \end{pmatrix}$   $ZX = \begin{pmatrix} 0 & 1 \\ -1 & 0 \end{pmatrix}$ 

$$(b) \quad \bigcup(o(7) = \bigcup \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix} = 2^{n} \operatorname{column} of U = \frac{1}{12} \begin{pmatrix} 0 \\ -1 \\ 0 \end{pmatrix}$$

$$(c) |\langle ob | U | lo \rangle|^{2} + |\langle lo b \rangle|^{2} = \frac{1}{2}$$

PROBLEM 1.2.

(a) H is hermitian & unbary => eigenvalues much be 
$$H_1 - 1$$
 products  $(1 + 1) = \frac{1}{16} (\frac{1}{16}) = \frac{1}{1$ 

PROBLEM 1.2 Continuel

NB: many different-looking 
$$\binom{i(1+Vi)}{1} = \frac{Vi}{2} (1+i) \binom{1}{i-Vi}$$
  
expressions are  $\binom{i(1+Vi)}{1} = \frac{Vi}{2} (1+i) \binom{1}{i-Vi}$   
equal  $\binom{i(1+Vi)}{1} = -\frac{Vi}{2} (1+i) \binom{1}{i+Vi}$ 

PROBLEM 1.3

$$(a) \quad \sqrt{(a|a)} = 1 \quad (b|b) = 1 \quad (a|b) = \frac{F_2}{3} - \frac{F_2}{3} = 0.$$

$$(2) | (< 0 | H 2 \pi H | 0) | =$$
  $| < -1 | 1 + ) |^{2} = \frac{1}{2}$ 

$$| < +|s|+ >|^{2} = | \frac{1}{2} (4+i)|^{2} = |\frac{1}{12}|^{2} = \frac{1}{2}$$
$$| < -|s|+ >|^{2} = |\frac{1}{2} (4-i)|^{2} + - = \frac{1}{2}$$

(C )

$$\left|\left<\alpha\right|\left(\frac{(1+i)}{2}\left(0\right)+\frac{(1-i)}{2}\left(1\right)\right)^{2}\right|$$

$$= \left| \frac{1}{2\sqrt{3}} \left( 1 + \sqrt{2} + (1 - \sqrt{2})i \right) \right|^{2}$$

$$= \frac{1}{12} \left[ 2 \left( \Lambda + \sqrt{2} \right)^{2} \right] = \frac{1}{2}$$

Since 1a) & 15) are the complex basis the other inner product must be 1/2 as hell.

PROBLEM 2 Full credit was given for any solution, only mistales penalized A move general statement proven below FY1: (a) A general statement (not required for credit) Let 147,147 Se any two n-gubit states and let { V1 ... V; }, - { V1, - Vim } be m sets of Unituries s.E. & ISKEN  $V_1^{k} \otimes \cdots \otimes V_{jk}^{k}$  acts on u gulits, and let  $\{\theta_1 \cdots \theta_{j_1}^{k}\}, \dots, \{\theta_n \cdots \theta_{j_m}^{m}\}$ O; E[0,21] be (arbitrary) phases. Then  $|\langle \Psi| \begin{bmatrix} \Im & i \Psi_{k}^{n} \\ \Im & e^{i\Psi_{k}} \end{bmatrix} \cdots \begin{bmatrix} \Im & i \Psi_{k}^{n} \\ \Im & e^{i\Psi_{k}} \end{bmatrix} |\Psi\rangle [$  $= \left| \begin{array}{c} e^{\sum_{k=1}^{j_{k}} \theta_{k}} \\ = \left| \begin{array}{c} e^{\sum_{k=1}^{j_{k}} \theta_{k}} \\ e^{\sum_{k=1}^{j_{k}} k} \end{array} \right| \left| \begin{array}{c} e^{\sum_{k=1}^{j_{k}} \theta_{k}} \\ e^{\sum_{k=1}^{j_{k}} k} \\ e^{\sum_{k=1}^{j_{k}} k} \end{array} \right| \left| \begin{array}{c} e^{\sum_{k=1}^{j_{k}} \theta_{k}} \\ e^{\sum_{k=1}^{j_{k}} k} \\ e^{\sum_{k=1}^{j_{k}} k} \end{array} \right| \left| \begin{array}{c} e^{\sum_{k=1}^{j_{k}} \theta_{k}} \\ e^{\sum_{k=1}^{j_{k}} k} \\ e^{\sum_{k=1}^{j_{k}} k} \\ e^{\sum_{k=1}^{j_{k}} k} \\ e^{\sum_{k=1}^{j_{k}} \theta_{k}} \\ e^{\sum_{k$  $|-> |e^{i\varphi} \propto |= |e^{i\varphi} || \propto |= |\alpha|, \# \varphi...$ 

(b) Full credit was given for any solution, only mistales penalized A move general statement proven below FYI:

GENERIC RESULT: DETECTABILITY OF LOCAL PHASE, WITH CONSTRUCTION (not needed for take-home assignment)

Let 
$$W = \mathbb{I} \oplus V = \begin{bmatrix} 4 \\ \hline V \end{bmatrix}$$

and () = A W B, A, W, B unitaries.

Then 
$$W \notin (10) | v \rangle + |1 \rangle | v \rangle = \frac{1}{n} (10) + e^{i\theta} | n \rangle | v \rangle.$$

Let 
$$W^{\ominus} = \mathbb{I} \oplus e^{i \oplus} V$$

$$W \stackrel{\phi}{=} \frac{1}{\sqrt{2}} (|0\rangle|V\rangle + \langle1\rangle|V\rangle) = \frac{1}{\sqrt{2}} (10\rangle + e^{i\varphi+\varphi}H\rangle)|V\rangle$$
  
Assume  $\varphi \neq \overline{\Pi}$ .

Now note 
$$|\langle +|(107 + e^{i\varphi} | 12)| \neq |\langle +|(105 + e^{i\varphi} | 12)|$$

$$\forall \Psi, \forall \Theta \neq \Pi$$
.  
Then measuring with respet any basis which includes  
the rec.  $A^{\dagger} | t \rangle | v \rangle$ , starting from  $B^{\dagger} | t \rangle | v \rangle$ 

detects (infuences outcome probabilities) between AWB & AW<sup>O</sup>B.

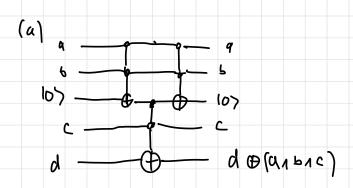
$$iF \quad \theta = \pi n, \quad \text{Pick} \quad \frac{1}{\sqrt{2}} (10) + e^{i\varphi} (10),$$

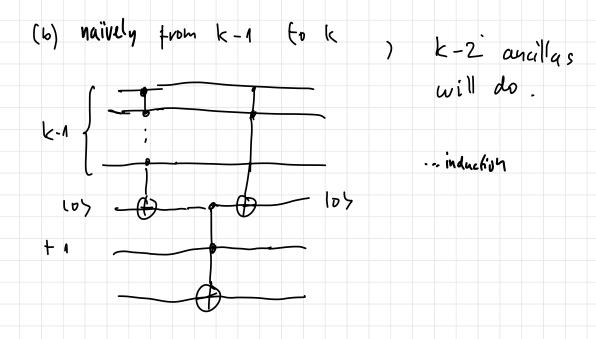
$$in \quad \text{place} \quad of \quad 1+\gamma.$$

The objective of the tack was to

- further evaluate capacity to evaluate circuits
- make you think about how phases multipling
- Unitaries percolak into amplitudes of g. stales
- make you note how measurement bases affect what aspects of amplitudes can be detected.
- develop a feeling for the "overcompletienness" of basic linear algebra description of RM

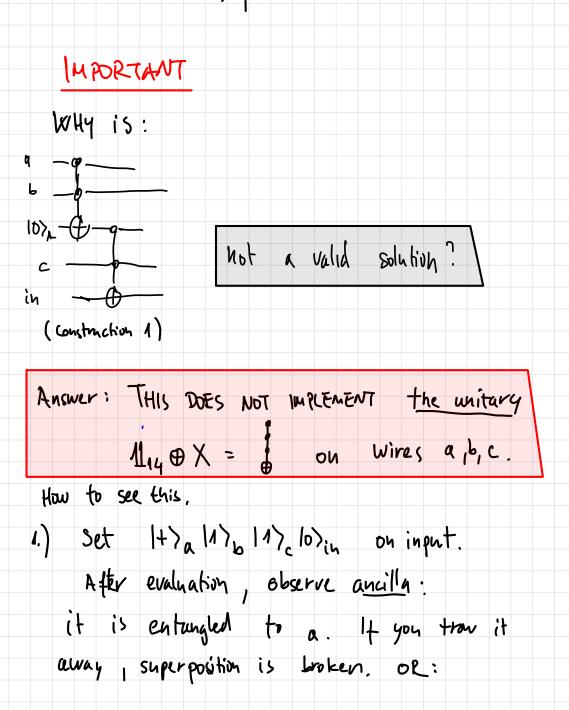
## PROBLEM 3





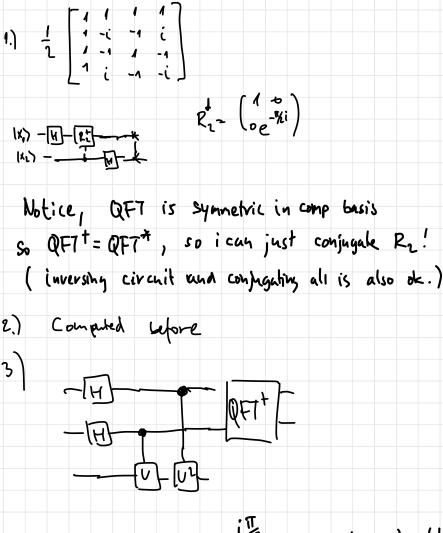
Much more space-efficient schemes possible utilizing <u>Phases</u> (think: Toffoli circuit does

a similar thing.



2) ctrl-Tottoli is SECF-INVERSE. The construction 1 suggests a new ancilla per call (it is set to 10>) q \_\_\_\_\_ 6 \_\_\_\_\_ 107A,--С \_\_\_\_\_ ίμ\_\_\_\_\_ This is not 11 on a, b, c, in, and it must be Set 1+7 M76 M2 loring to see ... But then the construction is Not ctr-Toffoli ... not a "dirty trick"... you were advised to "see page 183 of N&C"... LESSON: GARBAGE COLLECTION (Uncomputing ancilla gubits) is becessary for Q.C.

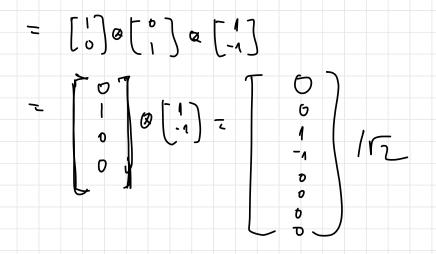
PROBLEM 4



4) 1-7 is an e<sup>itz</sup>-eigenvector of U!

i= exp("/2i) = exp(0.01 UTi)

=> By correctness of QPE, output is ID>IA> => =>



NB: it was only asked to produce the output state in vector form ...

5) しろして キ しょうしし も ==[+>(10>+i11>) 1-> → - (10>-11) (10>+i(1)) -> -> QFT property -D 1071171->

PROBLEM 5.1

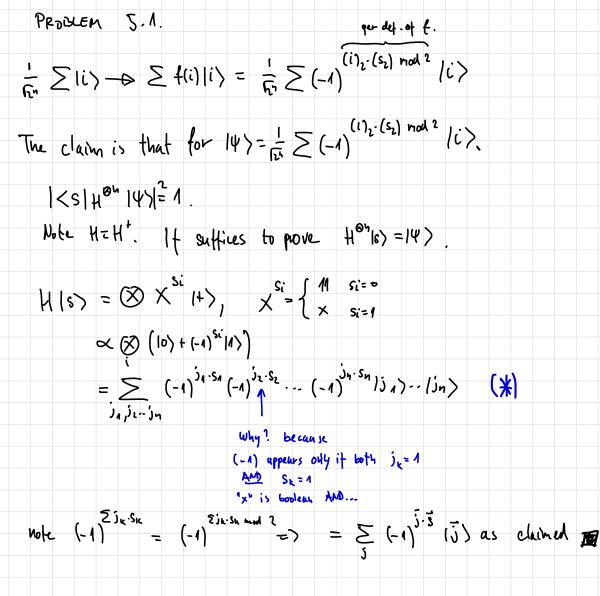
For PART (1) THE POINT WAS a) TO OBSERVE That Grover Can be applied, i.e. that the "databak" is the list of integers up to w, and to recall that given k "marked elements" the complexity is the i

b) to show that for w-smooth Square FREE numbers

(N is symme tree if all prime fadors occur at most once)
the number of divisors below w must be log w.
Since we forgot to specify square-tree-ness,
full credit was given tor a) alone

Proof. Let N= TTT; , P; #Pe # j#e, P; prime.

Claim:  $k \ge \lfloor \log_{W} N \rfloor$ Proof.  $k < \lfloor \log_{W} N \rfloor = > N < C_{U}^{\lfloor \log_{W} N \rfloor} \le W^{\log_{W} N} = N$ so N < N ~ 2 ~ 2. Note  $\log_{W} N = \frac{\log_{N} N}{\log_{W}}$ . Also note  $P_{3} \le W$ , so  $= > O(\sqrt{\frac{W}{k}}) = O(\sqrt{\frac{W \log_{W} W}{\log_{W}}})$ 



Note (\*) is the only non-tra observation ...