

Compilerconstructie

najaar 2016

<http://www.liacs.leidenuniv.nl/~vlietrvan1/coco/>

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college 9, woensdag 23 november 2016

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Code Optimization (1)

8.5 Optimization of Basic Blocks

To improve running time of code

- **Local** optimization: within block
- **Global** optimization: across blocks

Local optimization benefits from DAG representation of basic block

A slide from lecture 5

6.1 Variants of Syntax Trees

Directed Acyclic Graphs for Expressions

$$a + a * (b - c) + (b - c) * d$$

Syntax tree vs DAG...

Pros DAG...

8.5.1 DAG Representation of Basic Blocks

1. A node for initial value of each variable appearing in block
2. A node N for each statement s in block
Children of N are nodes corresponding to last definitions of operands used by s
3. Node N is labeled by operator applied at s
 N has list of variables for which s is last definition in block
4. *Output nodes* \approx live on exit

Example:

$a = b + c$

$b = a - d$

$c = b + c$

$d = a - d$

8.5.2 Finding Local Common Subexpressions

- Use value-number method to detect common subexpressions
- Remove redundant computations

Example:

$$a = b + c$$

$$b = a - d$$

$$c = b + c$$

$$d = a - d$$

Local Common Subexpression Elimination

- Use value-number method to detect common subexpressions
- Remove redundant computations

Example:

$$a = b + c$$

$$b = a - d$$

$$c = b + c$$

$$d = a - d$$

$$a = b + c$$

$$b = a - d$$

$$c = b + c$$

$$d = b$$

8.5.3 Dead Code Elimination

- Remove roots with no live variables attached
- If possible, repeat

Example:

$a = b + c$

$b = b - d$

$c = c + d$

$e = b + c$

No common subexpression

If c and e are not live. . .

Dead Code Elimination

- Remove roots with no live variables attached
- If possible, repeat

Example:

$a = b + c$

$b = b - d$

$c = c + d$

$e = b + c$

$a = b + c$

$b = b - d$

No common subexpression

If c and e are not live...

8.5.5 Representation of Array References

```
x = a[i]
```

```
y = x+z
```

```
z = a[i]
```

DAG...

Representation of Array References

```
x = a[i]  
a[j] = y  
z = a[i]
```

DAG...

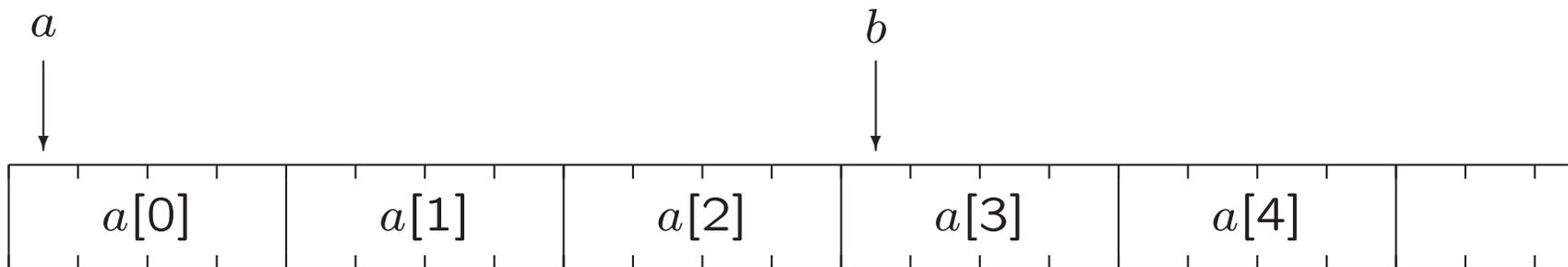
Representation of Array References

`b = 12 + a`

`x = b[i]`

`a[j] = y`

`z = b[i]`



DAG...

8.5.6 Pointer Assignments and Procedure Calls

a = b + c

e = a - d

c = b + c

b = a - d

DAG...

Pointer Assignments vs. Common Subexpressions

`p = &a`

`a = b + c`

`e = a - d`

`*p = y`

`c = b + c`

`b = a - d`

DAG...

Pointer Assignments vs. Common Subexpressions

a = b + c

e = a - d

*p = y

c = b + c

b = a - d

DAG...

Pointer Assignments vs. Dead Code

```
a = b + c  
b = b - d  
c = c + d  
e = b + c  
x = *p
```

DAG...

If c and e are not live...

To summarize:

`*q = y`

`x = *p`

Procedure calls...

8.5.4 The Use of Algebraic Identities

and other algebraic transformations

(cf. assignment 3)

Algebraic identities:

$$\begin{aligned}x + 0 &= 0 + x = x \\x * 1 &= 1 * x = x\end{aligned}$$

Reduction in strength:

$$\begin{aligned}x^2 &= x * x && \text{(cheaper)} \\2 * x &= x + x && \text{(cheaper)} \\x/2 &= x * 0.5 && \text{(cheaper)}\end{aligned}$$

Constant folding:

$$2 * 3.14 = 6.28$$

Algebraic Transformations

Common subexpressions resulting from commutativity / associativity of operators:

$$\begin{aligned}x * y &= y * x \\c + d + b &= (b + c) + d\end{aligned}$$

Common subexpressions generated by relational operators:

$$x > y \Leftrightarrow x - y > 0$$

8.7 Peephole Optimization

- Examines short sequence of instructions in a window (peep-hole) and replace them by faster/shorter sequence
- Applied to intermediate code or target code
- Typical optimizations
 - Redundant instruction elimination
 - Eliminating unreachable code
 - Flow-of-control optimization
 - Algebraic simplification
 - Use of machine idioms

8.7.1 Eliminating Redundant Loads and Stores

Naive code generator may produce

```
ST  a, R0  
LD  R0, a
```

N.B.: optimize only within basic block

8.7.2 Eliminating Unreachable Code

Example:

```
    if debug == 1 goto L1
    goto L2
L1: print debugging information
L2:
```

Jump over jump

Eliminating Unreachable Code

Example:

```
    if debug != 1 goto L2  
L1: print debugging information  
L2:
```

If debug is set to 0 at beginning of program, ...

8.7.3 Flow-of-Control Optimizations

Example 1:

```
    goto L1
    ...
L1: goto L2
```

Example 3:

```
    goto L1
    . . .
L1: if a < b goto L2
L3:
```

8.7.3 Flow-of-Control Optimizations

Example 1:

```
    goto L1
    ...
L1: goto L2
```

```
    goto L2
    ...
L1: goto L2
```

Example 3:

```
    goto L1
    . . .
L1: if a < b goto L2
L3:
```

```
    if a < b goto L2
    goto L3
    ...
L3:
```

9.1 The Principal Sources of Optimization

Causes of redundancy

- At source level
- Side effect of high-level programming language, e.g., $A[i][j]$

9.1.2 A Running Example: Quicksort

```
void quicksort (int m, int n)
    /* recursively sorts a[m] through a[n] */
{
    int i, j;
    int v, x;

    if (n <= m) return;

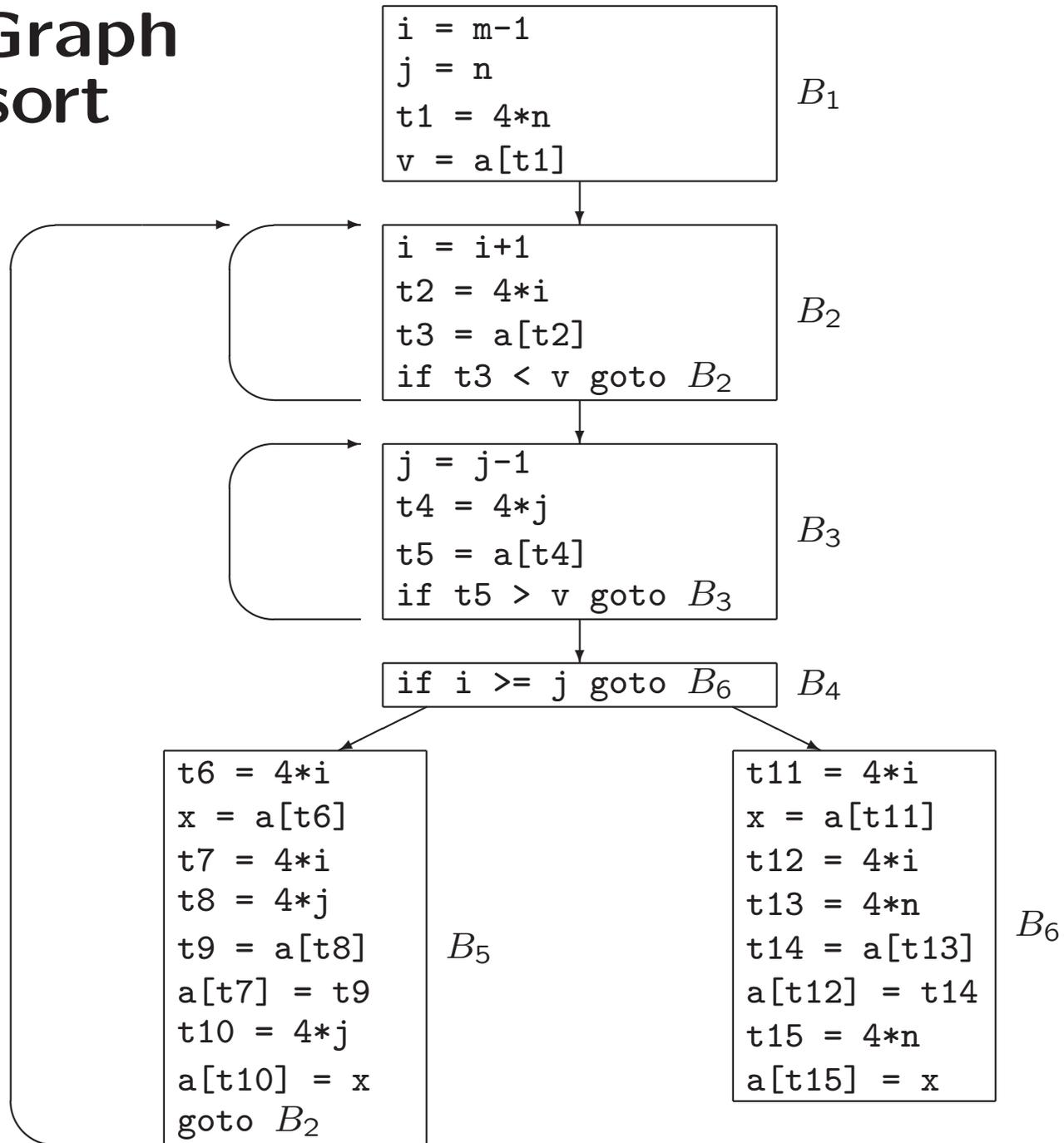
    i = m-1; j = n; v = a[n];
    while (1)
    {    do i = i+1; while (a[i] < v);
        do j = j-1; while (a[j] > v);
        if (i >= j) break;
        x = a[i]; a[i] = a[j]; a[j] = x; /* swap a[i], a[j] */
    }
    x = a[i]; a[i] = a[n]; a[n] = x; /* swap a[i], a[n] */

    quicksort(m,j); quicksort(i+1,n);
}
```

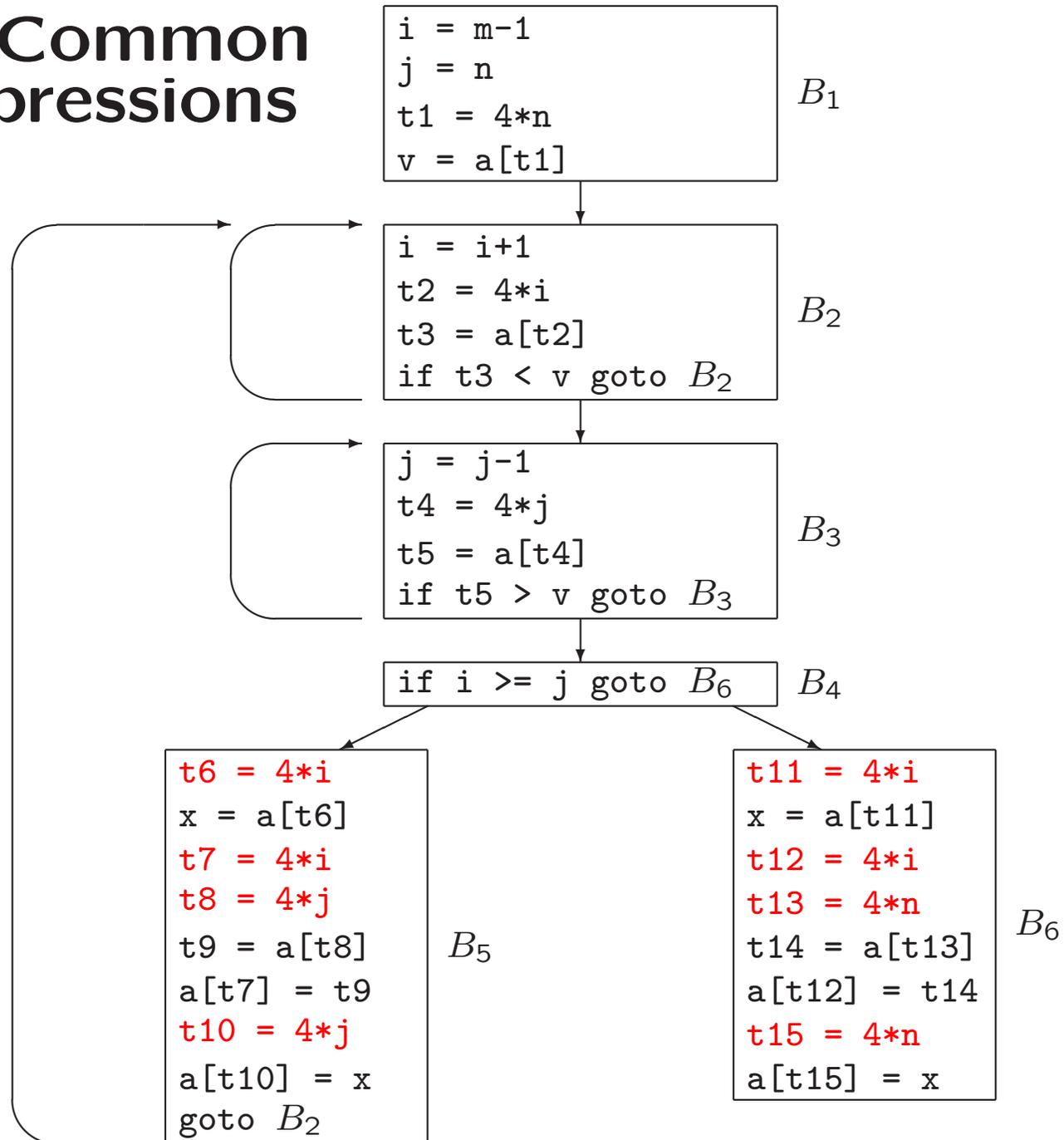
Three-Address Code Quicksort

```
→ (1) i = m-1
   (2) j = n
   (3) t1 = 4*n
   (4) v = a[t1]
→ (5) i = i+1
   (6) t2 = 4*i
   (7) t3 = a[t2]
   (8) if t3<v goto (5)
→ (9) j = j-1
   (10) t4 = 4*j
   (11) t5 = a[t4]
   (12) if t5>v goto (9)
→ (13) if i>=j goto (23)
→ (14) t6 = 4*i
   (15) x = a[t6]
   (16) t7 = 4*i
   (17) t8 = 4*j
   (18) t9 = a[t8]
   (19) a[t7] = t9
   (20) t10 = 4*j
   (21) a[t10] = x
   (22) goto (5)
→ (23) t11 = 4*i
   (24) x = a[t11]
   (25) t12 = 4*i
   (26) t13 = 4*n
   (27) t14 = a[t13]
   (28) a[t12] = t14
   (29) t15 = 4*n
   (30) a[t15] = x
```

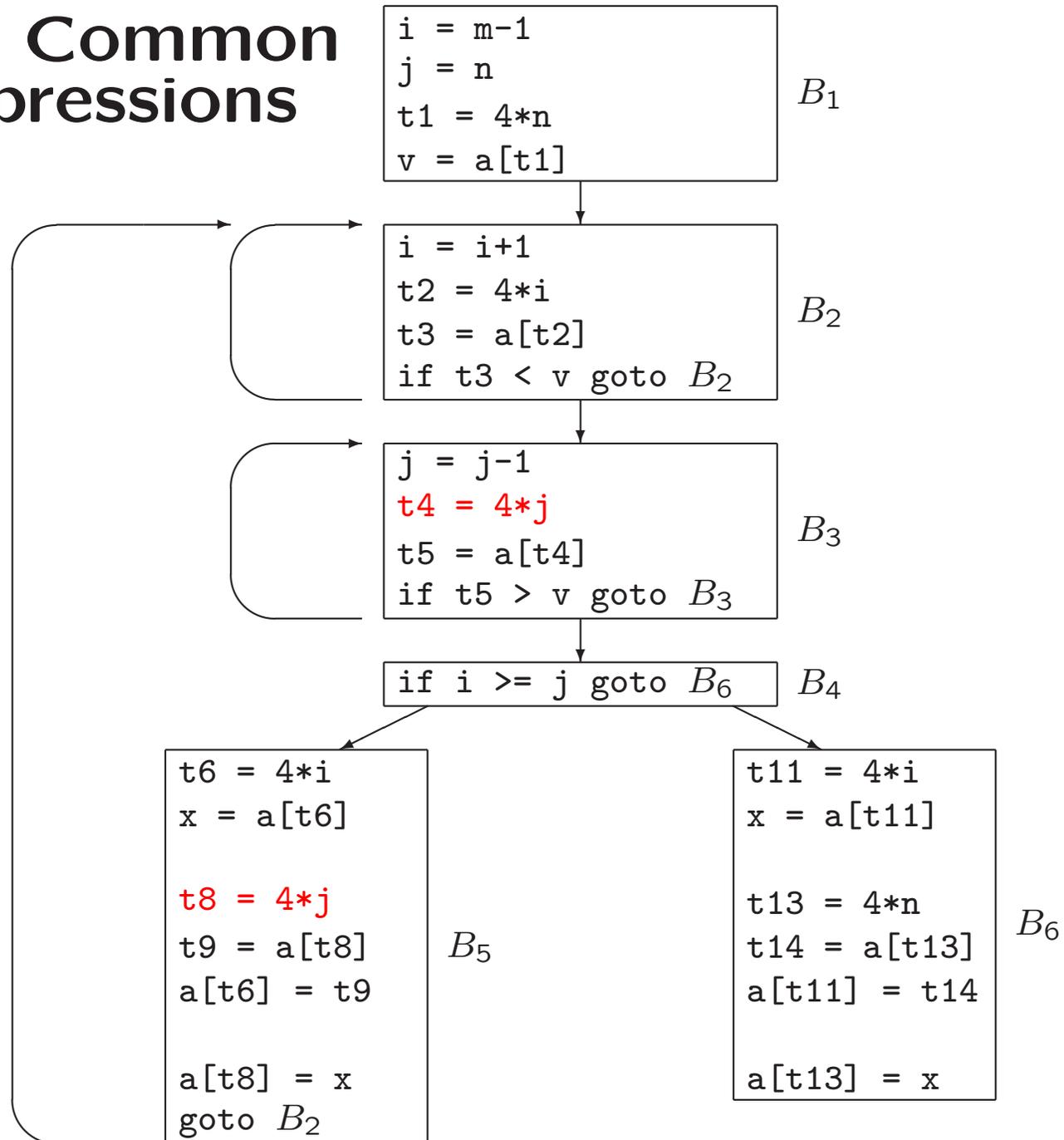
Flow Graph Quicksort



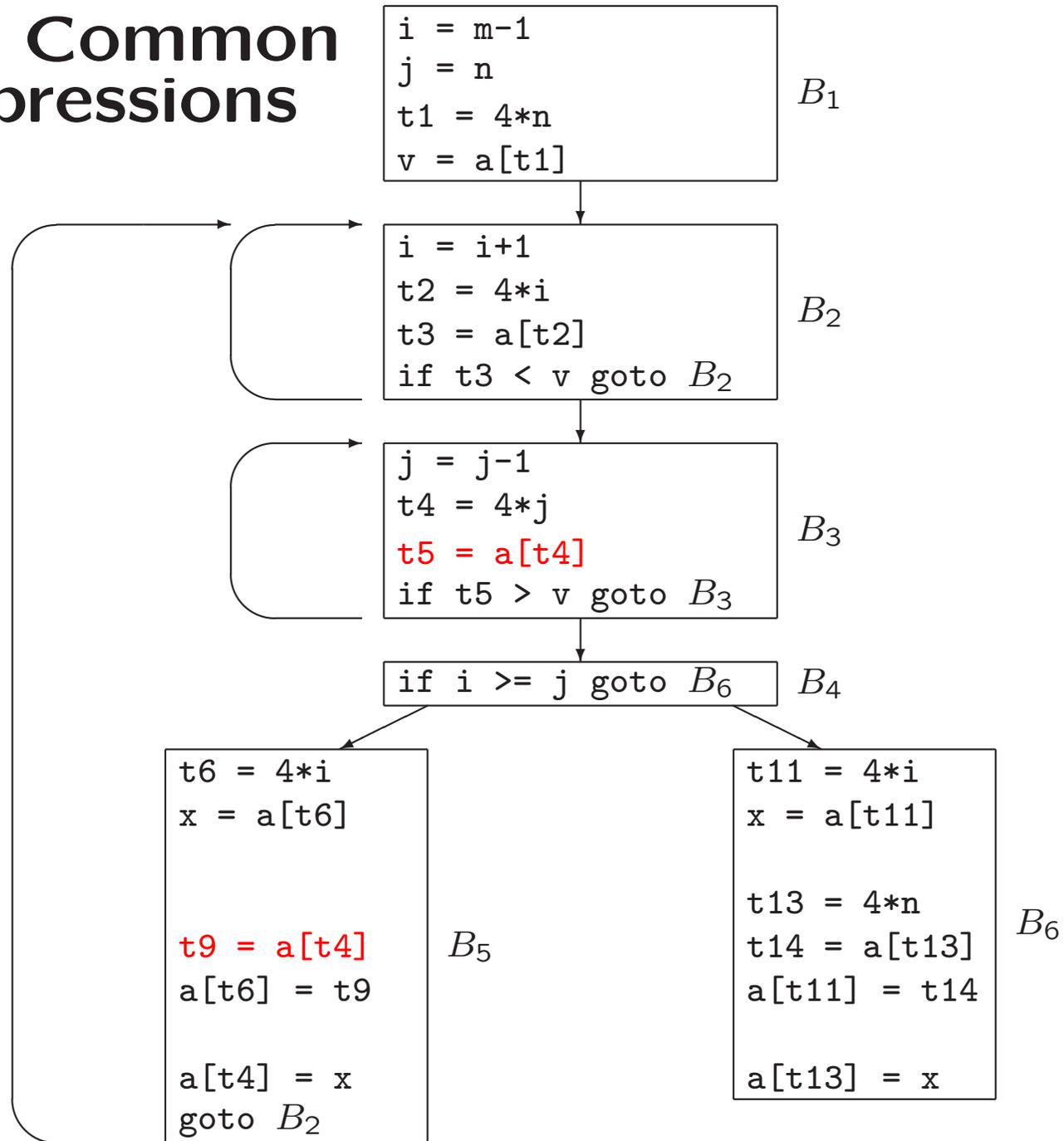
Local Common Subexpressions



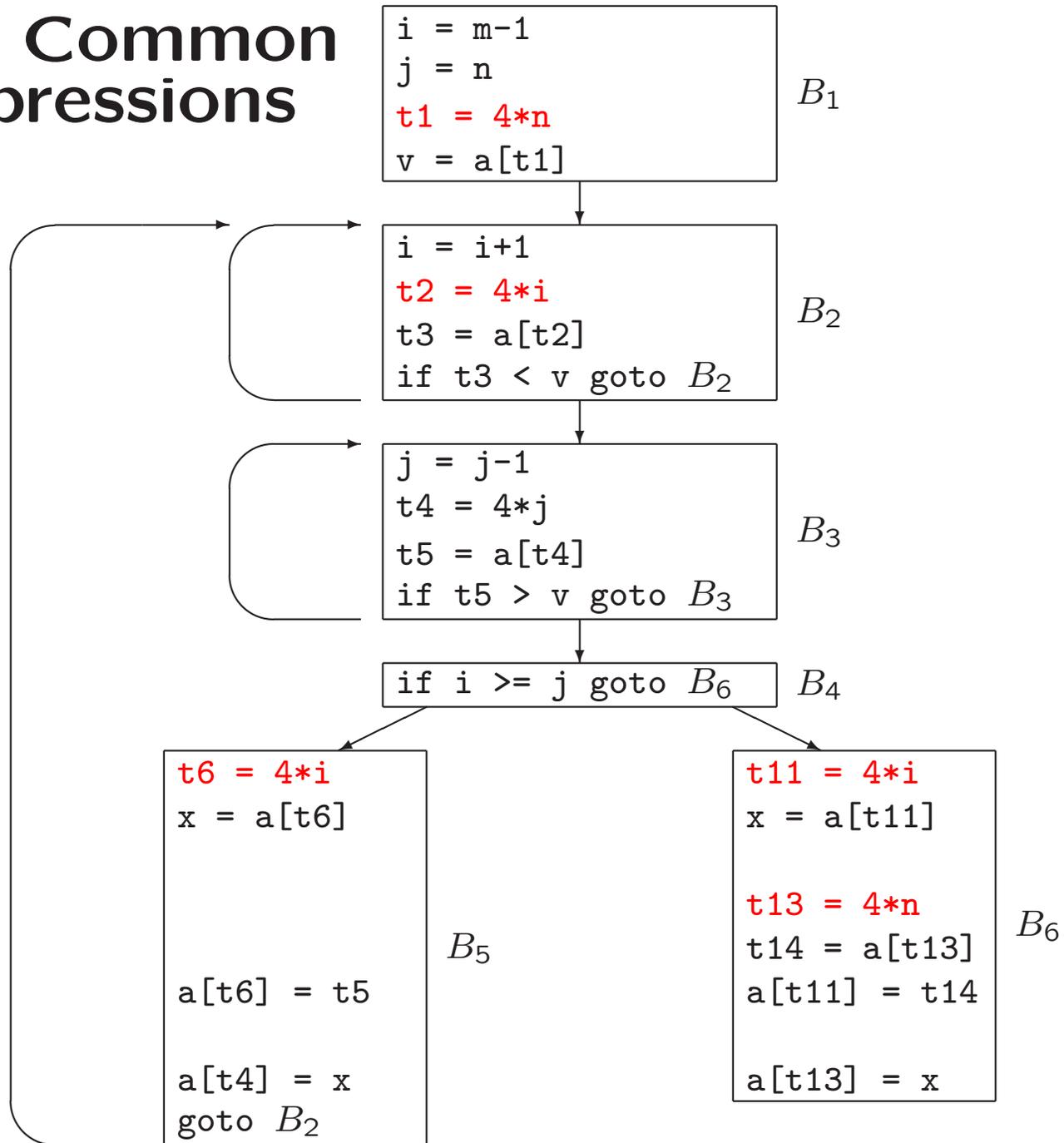
Global Common Subexpressions



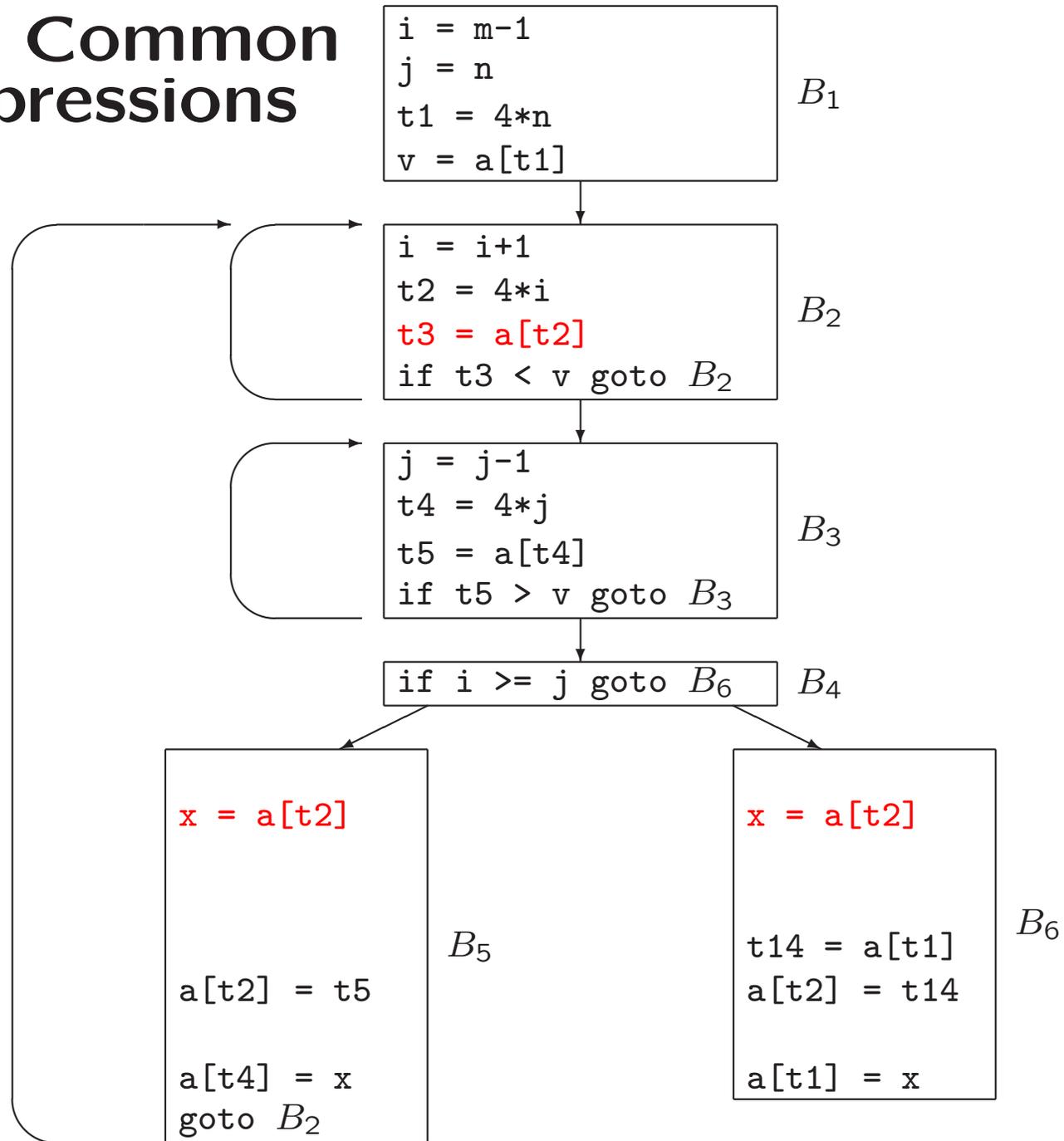
Global Common Subexpressions



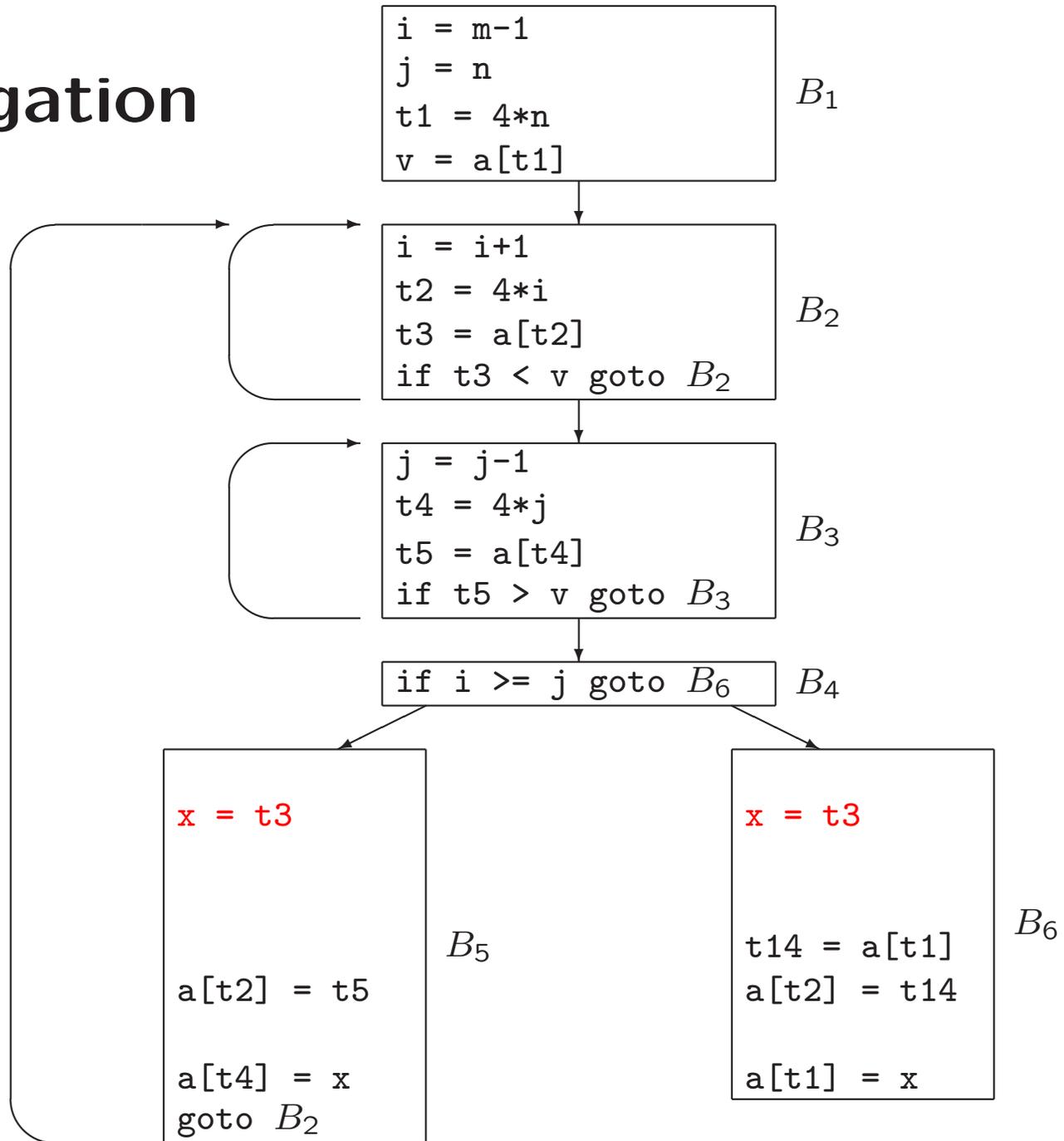
Global Common Subexpressions



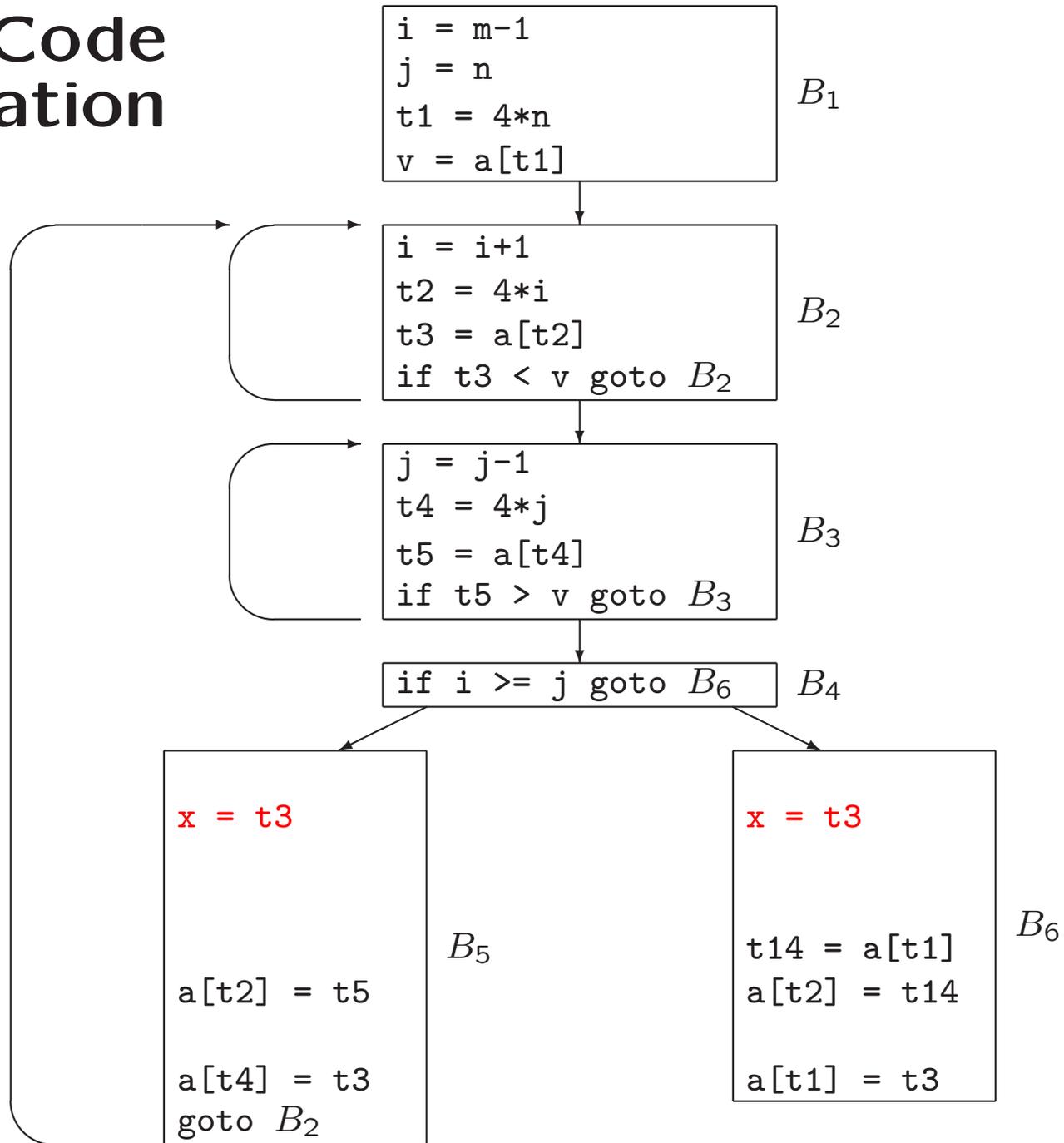
Global Common Subexpressions



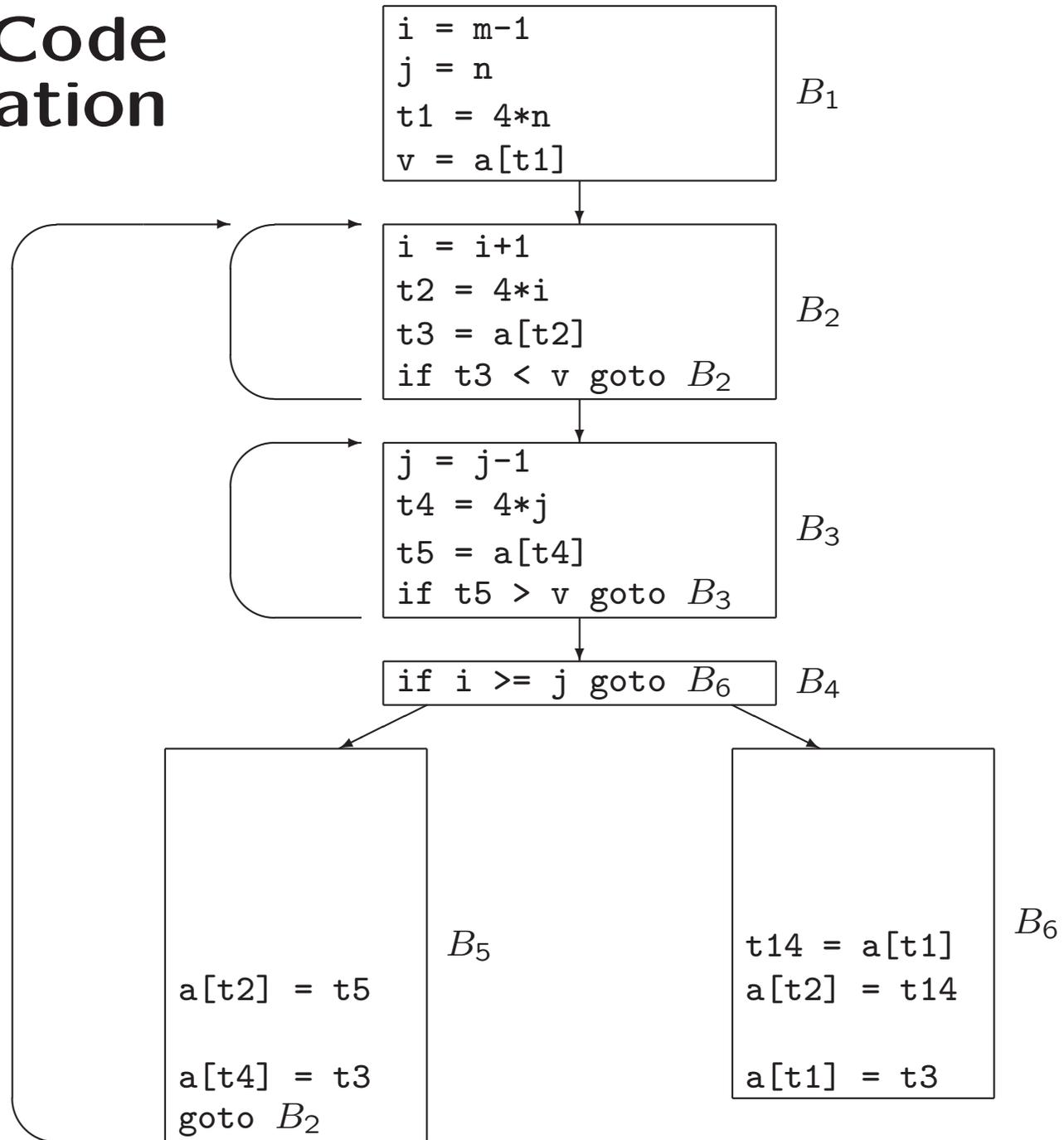
Copy Propagation



Dead-Code Elimination



Dead-Code Elimination



Code Motion

- loop-invariant computation
- compute **before** loop
- Example:

```
while (i <= limit-2) /* statement does not change limit */
```

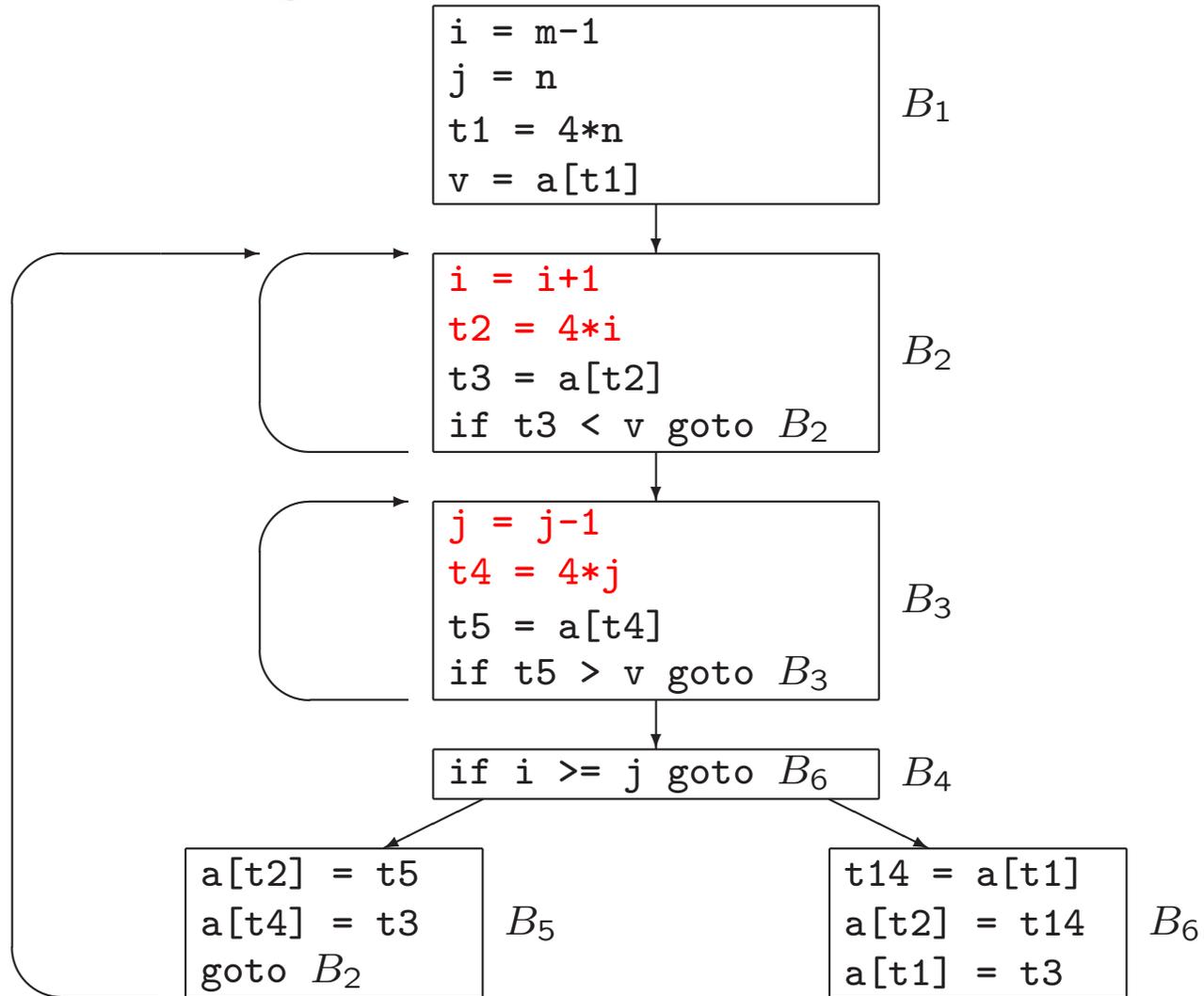
After code-motion

```
t = limit-2  
while (i <= t) /* statement does not change limit or t */
```

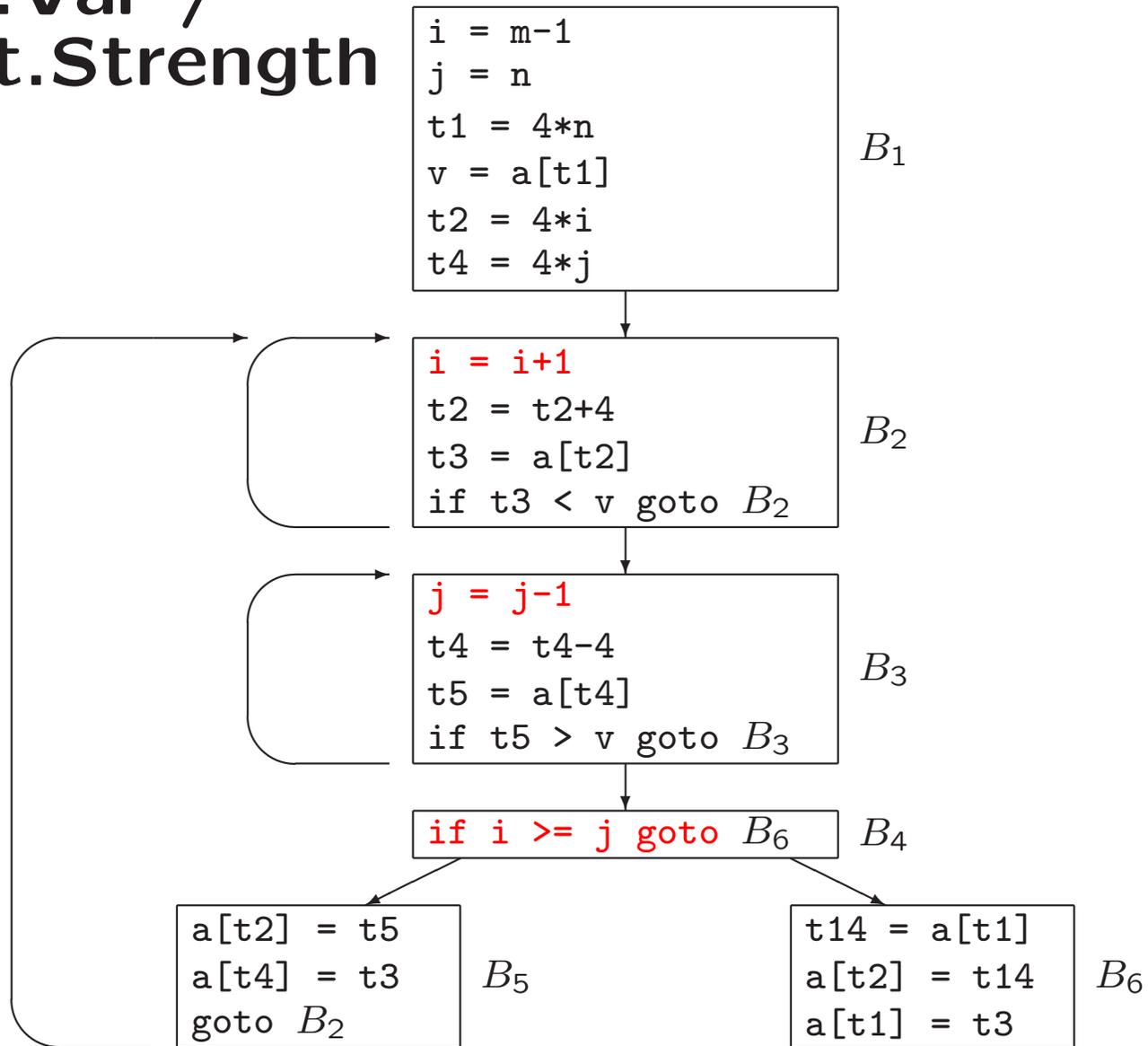
Induction Variables and Reduction in Strength

- **Induction variable:** each assignment to x of form $x = x + c$
- **Reduction in strength:** replace expensive operation by cheaper one

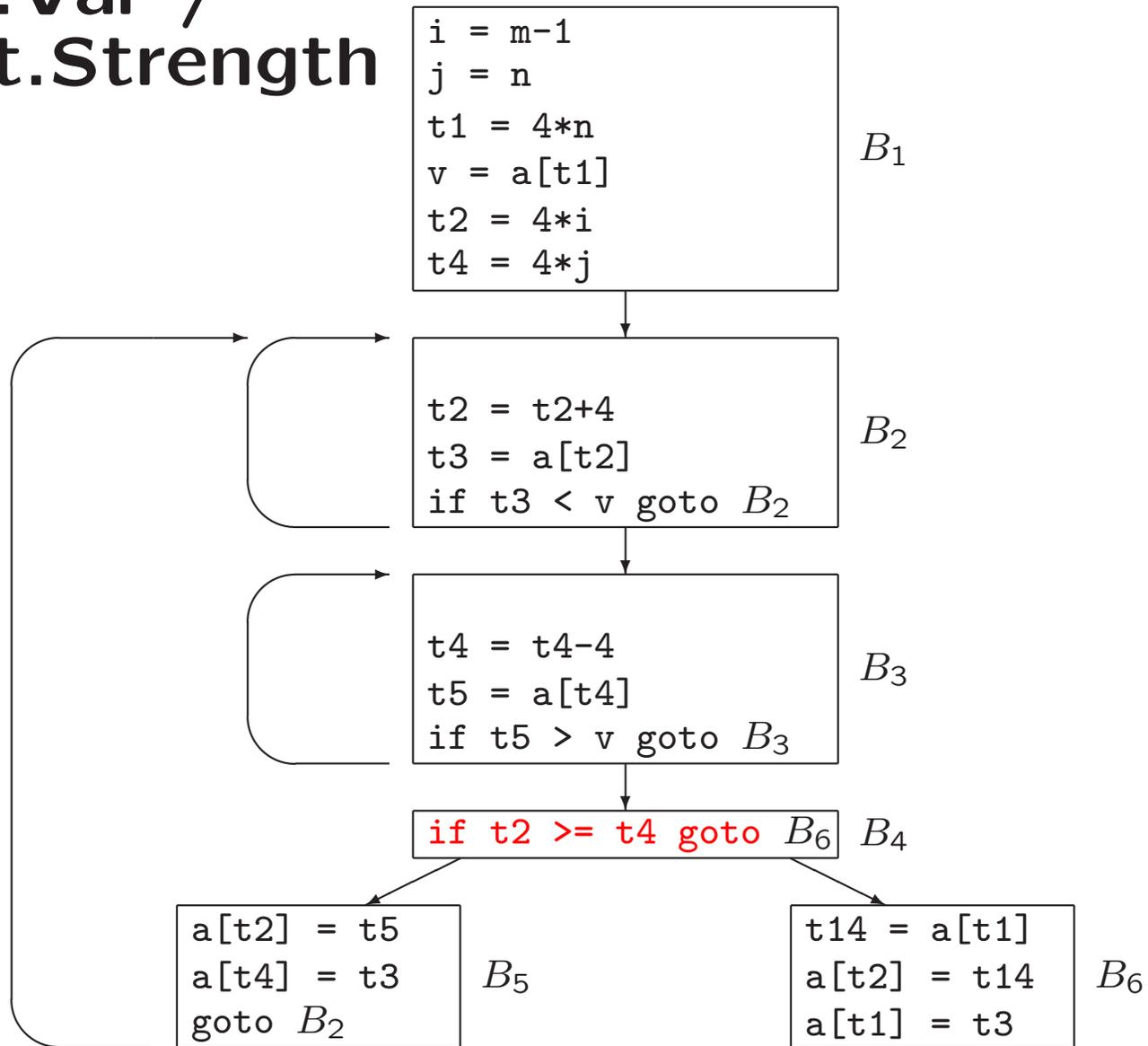
Induct.Var / Reduct.Strength



Induct.Var / Reduct.Strength

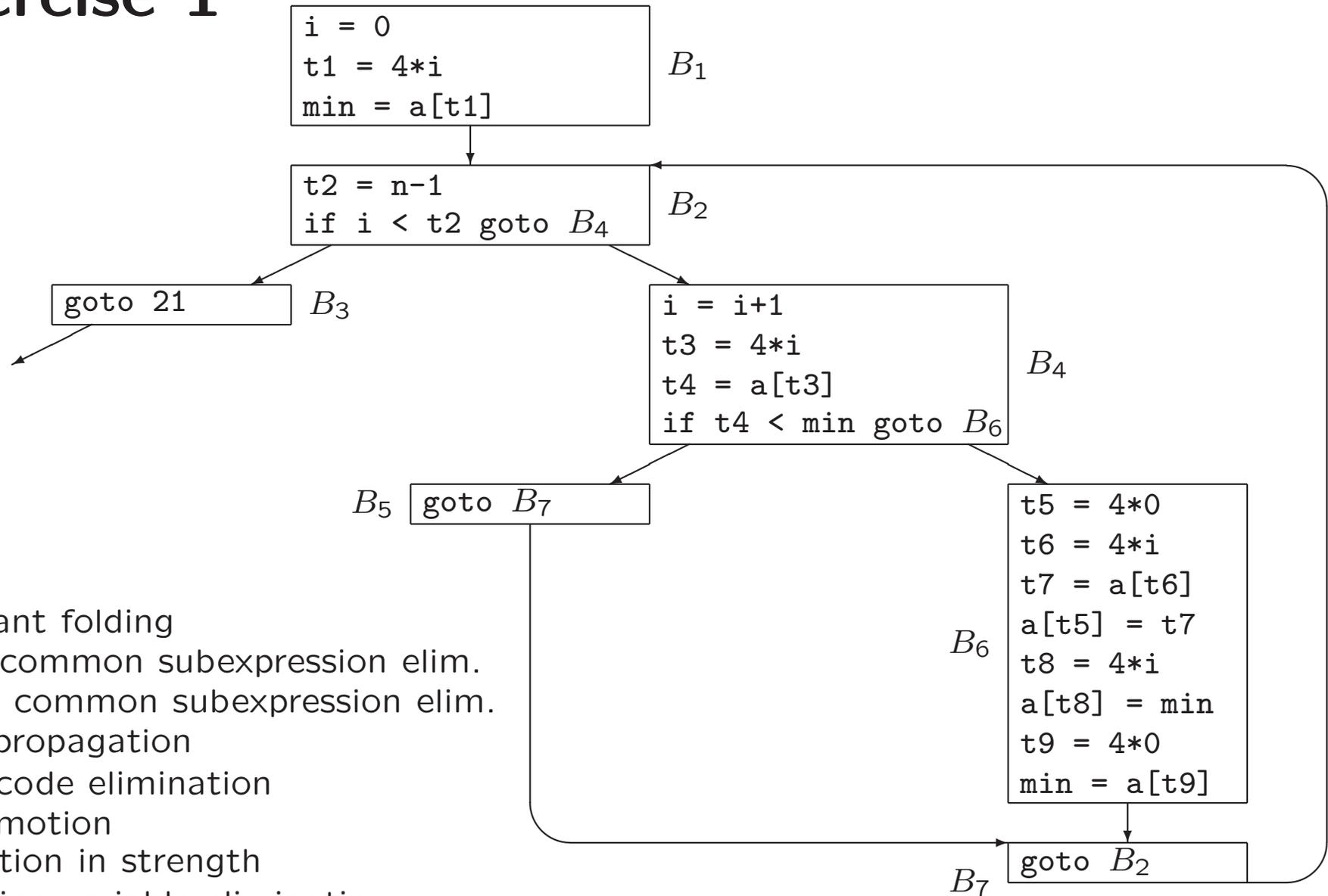


Induct.Var / Reduct.Strength



Exercise 1

Flow Graph Exercise 1



Constant folding
Local common subexpression elim.
Global common subexpression elim.
Copy propagation
Dead-code elimination
Code motion
Reduction in strength
Induction-variable elimination

Volgende week

- Practicum over opdracht 4
- Inleveren 13 december
- Woensdag 7 december: laatste hoor-/werkcollege

Compilerconstructie

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Code Optimization

Chapters for reading:

8.5–8.5.6, 8.7

9.intro, 9.1