Compilerconstructie

najaar 2016

http://www.liacs.leidenuniv.nl/~vlietrvan1/coco/

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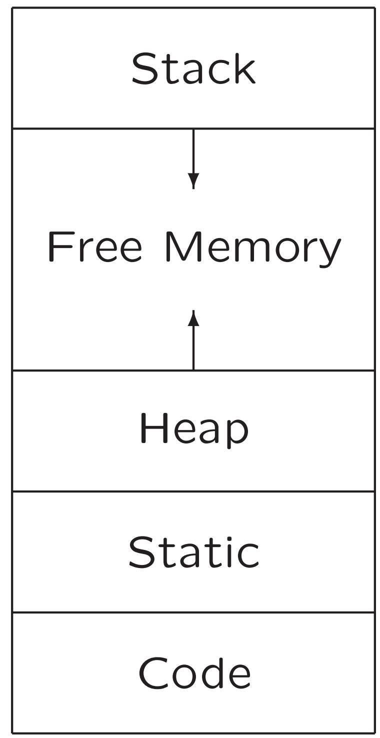
college 8, woensdag 16 november 2016
+ 'werkcollege'

Storage Organization

Code Generation

7.1 Storage Organization



Typical subdivision of run-time memory into code and data areas

7.1.1 Static Versus Dynamic Storage Allocation

- Static: compile time
- Dynamic: run time

Dynamic storage allocation:

- Stack storage: for data local to procedure
- Heap storage: for data that outlives procedure

Garbage collection to support heap management

7.2 Stack Allocation of Space

Possible because procedure calls are nested

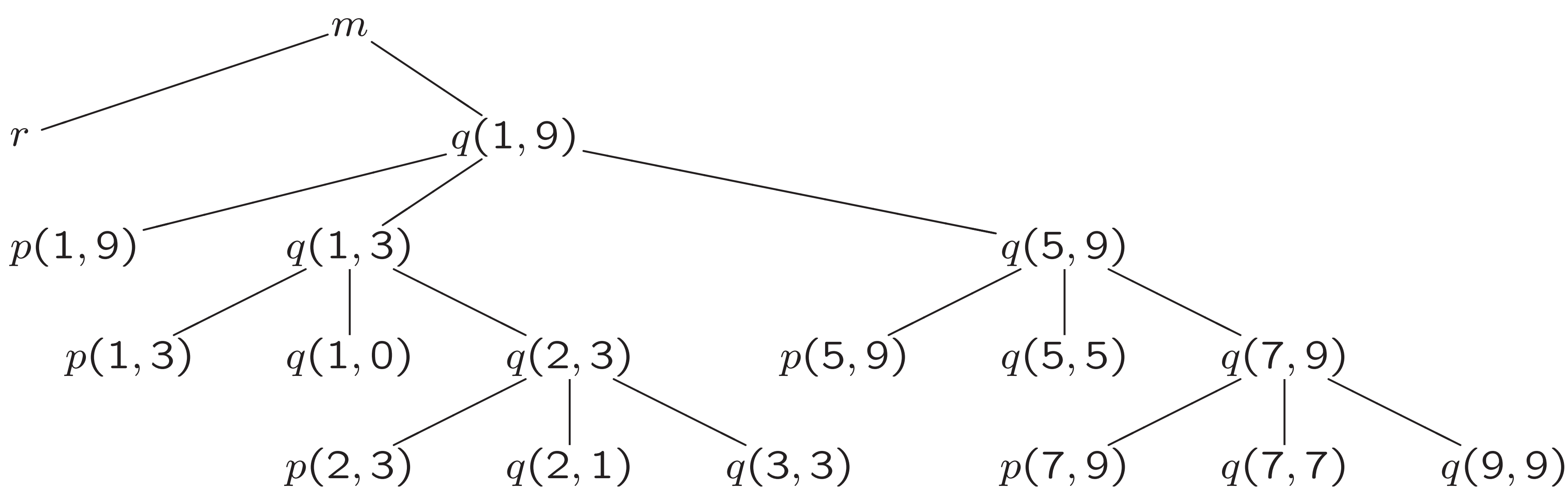
7.2 Stack Allocation of Space

```
int a[11];
void readArray()  /* Reads 9 integers into a[1],...a[9]. */
{ int i;
  ...
}
int partition (int m, int n)
{ /* Picks a separator value v, and partitions a[m..n] so that
     a[m..p-1] are less than v, a[p]=v, and a[p+1..n} are
     equal to or greater than v. Returns p. */
  ...
}
void quicksort (int m, int n)
{ int i;
  if (n > m)
  { i = partition(m, n);
    quicksort(m, i-1);
    quicksort(i+1, n);
  }
}
main ()
{ readArray();
  a[0] = -9999;
  a[10] = 9999;
  quicksort(1,9);
}
```

Possible Activations

```
enter main()
    enter readArray()
    leave readArray()
    enter quicksort(1,9)
        enter partition(1,9)
        leave partition(1,9)
        enter quicksort(1,3)
           ...
        leave quicksort(1,3)
        enter quicksort(5,9)
           ...
        leave quicksort(5,9)
    leave quicksort(1,9)
leave main()
```

7.2.1 Activation Trees



Traversal of Activation Tree

1. Sequence of procedure *calls* $\approx \ldots$ traversal

2. Sequence of procedure *returns* $\approx \ldots$ traversal

3. When control lies at particular node ($\approx$ activation), the ‘open’ (*live*) activations are . . .

Traversal of Activation Tree

1. Sequence of procedure *calls* $\approx$ preorder traversal

2. Sequence of procedure *returns* $\approx$ postorder traversal

3. When control lies at particular node ($\approx$ activation), the 'open' (*live*) activations are on path from root

7.2.2. Activation Records

(= stack frames)

Actual parameters
Returned values
Control link
Access link
Saved machine status
Local data
Temporaries

Possible (order of) elements of activation record

7.2.3 Calling Sequences

- Code to allocate (and fill) activation record on stack
- Divided between caller (at every location) and callee
- Return sequences analogous

8 Code Generation

- Output code must
  - be correct
  - use resources of target machine effectively

- Code generator must run efficiently

Generating optimal code is undecidable problem
Heuristics are available

8.1 Issues in Design of Code Generator

- Input to the code generator
- The target program
- Instruction selection
- Register allocation and assignment
- Evaluation order

8.1.1 Input to the Code Generator

- Intermediate representation of source program
  - Three-address representations (e.g., quadruples)
  - Virtual machine representations (e.g., bytecodes)
  - Postfix notation
  - Graphical representations (e.g., syntax trees and DAGs)
- Information from symbol table to determine run-time addresses
- Input is free of errors
  - Type checking and conversions have been done

8.1.2 The Target Program

- Common target-machine architectures
  - RISC: reduced instruction set computer
  - CISC: complex instruction set computer
  - Stack-based

- Possible output
  - Absolute machine code (executable code)
  - Relocatable machine code (object files for linker)
  - Assembly-language

8.1.3 Instruction Selection

- Given IR program can be implemented by many different code sequences
- Different machine instruction speeds
- Naive approach: statement-by-statement translation, with a code template for each IR statement

  Example: $x = y + z$ Now, $a = b + c \quad d = a + e$

```
LD  R0, y
LD  R1, z
ADD R0, R0, R1
ST  x, R0
```

```
LD  R0, b
LD  R1, c
ADD R0, R0, R1
ST  a, R0
LD  R0, a
LD  R1, e
ADD R0, R0, R1
ST  d, R0
```

8.2 The Target Language

- Designing code generator requires understanding of target machine and its instruction set

- Our machine model
  - byte-addressable
  - has n general purpose registers $\mathtt{R0}, \mathtt{R1}, \ldots, \mathtt{R}n-1$
  - assumes operands are integers

Instructions of Target Machine

- Load operations: `LD` *dst*, *addr*
  e.g., `LD` r, x or `LD` r_1, r_2

- Store operations: `ST` x, r

- Computation operations: *OP dst*, *src*$_1$, *src*$_2$
  e.g., *SUB* r_1, r_2, r_3

- Unconditional jumps: `BR` L

- Conditional jumps: `B`*cond* r, L
  e.g., `BLTZ` r, L

Addressing Modes of Target Machine

Form	Address	Example
r	r	`LD R1, R2`
x	x	`LD R1, x`
$a(r)$	$a + \mathit{contents}(r)$	`LD R1, a(R2)`
$c(r)$	$c + \mathit{contents}(r)$	`LD R1, 100(R2)`
$*r$	$\mathit{contents}(\mathit{contents}(r))$	`LD R1, *R2`
$*c(r)$	$\mathit{contents}(c + \mathit{contents}(r))$	`LD R1, *100(R2)`
$\#c$		`LD R1, #100`

Addressing Modes (Examples)

```
b = a[i]:

  LD  R1, i
  MUL R1, R1, #8
  LD  R2, a(R1)
  ST  b, R2
```

```
x = *p

  LD  R1, p
  LD  R2, 0(R1)
  ST  x, R2
```

```
a[j] = c

  LD  R1, c
  LD  R2, j
  MUL R2, R2, #8
  ST  a(R2), R1
```

```
if x < y goto L

  LD   R1, x
  LD   R2, y
  SUB  R1, R1, R2
  BLTZ R1, M
```

8.2.2 Program and Instruction Costs

- Costs associated with compiling / running a program
  - Compilation time
  - Size, running time, power consumption of target program

- Finding optimal target problem: undecidable

- (Simple) cost per target-language instruction:
  - 1 + cost for addressing modes of operands
    $\approx$ length (in words) of instruction

  Examples:

instruction	cost
`LD R0, R1`	1
`LD R0, x`	2
`LD R1, *100(R2)`	2

8.4 Basic Blocks and Flow Graphs

1. Basic block: maximal sequence of consecutive three-address instructions, such that

   (a) Flow of control can only enter through first instruction of block

   (b) Control leaves block without halting or branching

2. Flow graph: graph with
   nodes: basic blocks
   edges: indicate flow between blocks

8.4.1 Determining Basic Blocks

- Determine leaders

  1. First three-address instruction is leader

  2. Any instruction that is target of goto is leader

  3. Any instruction that immediately follows goto is leader

- For each leader, its basic block consists of leader and all instructions up to next leader (or end of program)

Determining Basic Blocks (Example)

Determine leaders

Pseudo code

for $i = 1$ to 10 **do**
    for $j = 1$ to 10 **do**
        $a[i, j] = 0.0$;
for $i = 1$ to 10 **do**
    $a[i, i] = 1.0$;

Three-address code

```
 1)   i = 1
 2)   j = 1
 3)   t1 = 10 * i
 4)   t2 = t1 + j
 5)   t3 = 8 * t2
 6)   t4 = t3 - 88
 7)   a[t4] = 0.0
 8)   j = j + 1
 9)   if j <= 10 goto (3)
10)   i = i + 1
11)   if i <= 10 goto (2)
12)   i = 1
13)   t5 = i - 1
14)   t6 = 88 * t5
15)   a[t6] = 1.0
16)   i = i + 1
17)   if i <= 10 goto (13)
```

Determining Basic Blocks (Example)

Determine leaders

Pseudo code

for $i = 1$ to 10 **do**
  for $j = 1$ to 10 **do**
    $a[i, j] = 0.0;$
for $i = 1$ to 10 **do**
  $a[i, i] = 1.0;$

Three-address code

```
-->   1)   i = 1
-->   2)   j = 1
-->   3)   t1 = 10 * i
      4)   t2 = t1 + j
      5)   t3 = 8 * t2
      6)   t4 = t3 - 88
      7)   a[t4] = 0.0
      8)   j = j + 1
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-->  10)   i = i + 1
     11)   if i <= 10 goto (2)
-->  12)   i = 1
-->  13)   t5 = i - 1
     14)   t6 = 88 * t5
     15)   a[t6] = 1.0
     16)   i = i + 1
     17)   if i <= 10 goto (13)
```

8.4.3 Flow Graphs

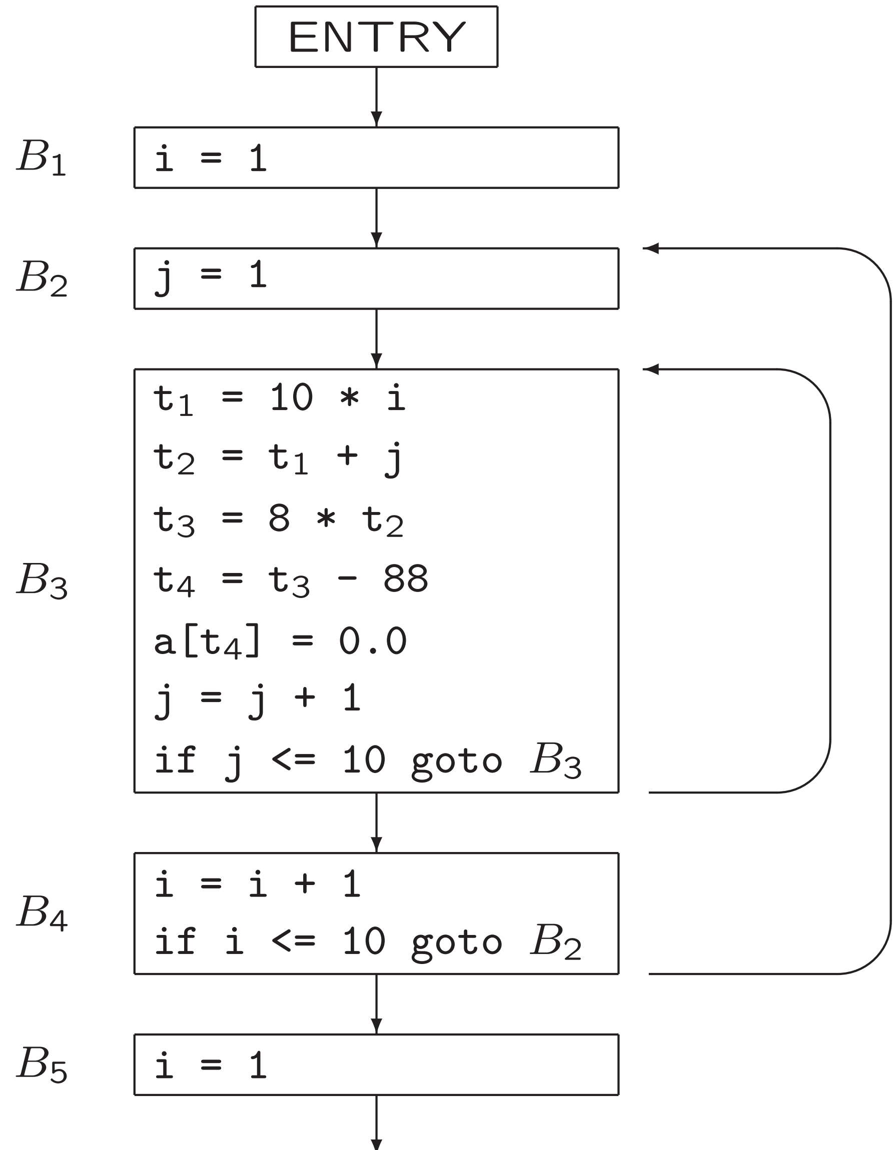
Edge from block B to block C

- if there is (un)conditional jump from end of B to beginning of C
- if C immediately follows B in original order, and B does not end in unconditional jump

Flow Graph (Example)

Three-address code

```
⟶   1)  i = 1
⟶   2)  j = 1
⟶   3)  t1 = 10 * i
     4)  t2 = t1 + j
     5)  t3 = 8 * t2
     6)  t4 = t3 - 88
     7)  a[t4] = 0.0
     8)  j = j + 1
     9)  if j <= 10 goto (3)
⟶  10)  i = i + 1
    11)  if i <= 10 goto (2)
⟶  12)  i = 1
⟶  13)  t5 = i - 1
    14)  t6 = 88 * t5
    15)  a[t6] = 1.0
    16)  i = i + 1
    17)  if i <= 10 goto (13)
```



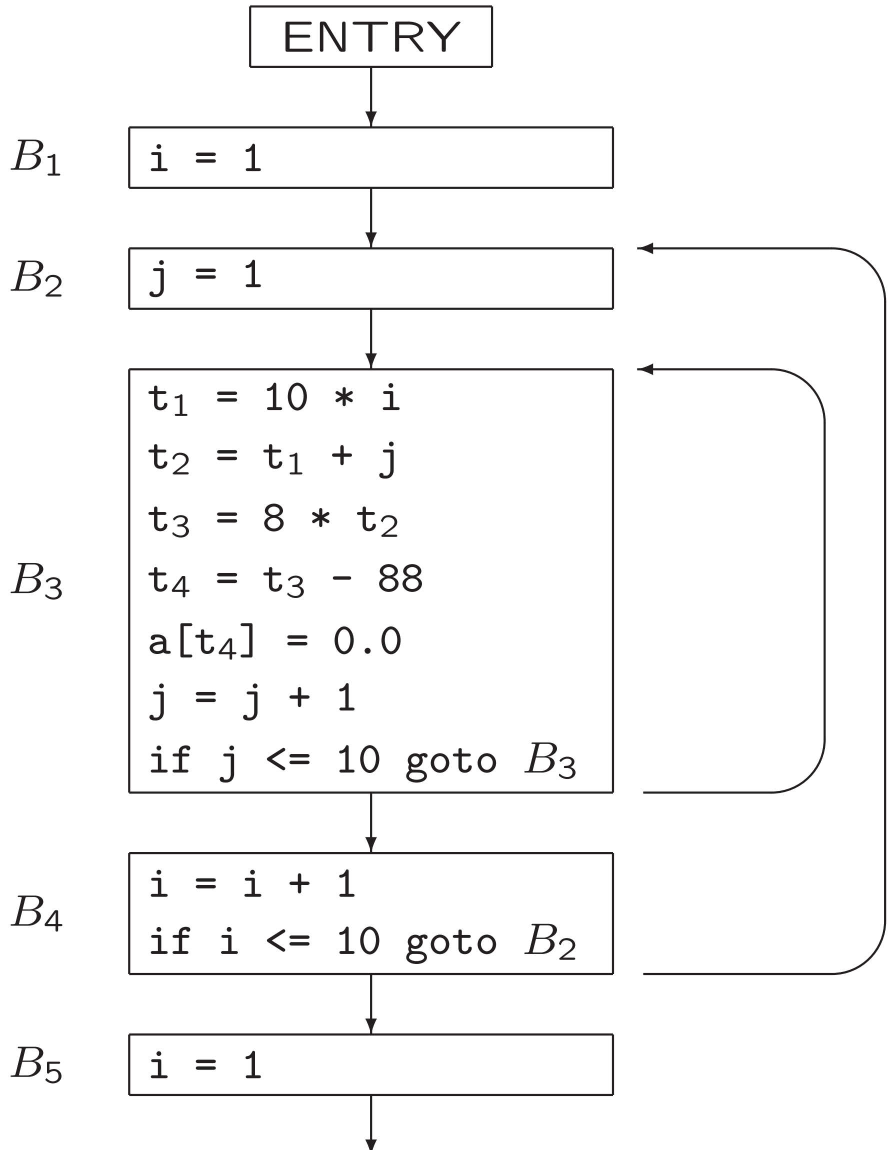
8.4.5 Loops

Loop is set of nodes in flow graph

- With unique loop entry e
- Every node in L has nonempty path in L to e

Example

- $\{B_3\}$, with loop entry B_3
- $\{B_2, B_3, B_4\}$, with loop entry B_2
- $\{B_6\}$, with loop entry B_6



8.4.2 Next-Use Information

- Next-use information is needed for dead-code elimination and register assignment

```
(i)  x = a * b
...
(j)  z = c + x
```

Instruction j uses value of x computed at i
x is live at i,
i.e., we need value of x later

- For each three-address statement $x = y\ op\ z$ in block, record next-uses of x, y, z

Determining Next-Use Information

For single basic block

- Assume all non-temporary variables are live on exit (stored in symbol table)

- Make backward scan of instructions in block

- For each instruction i: $x = y$ *op* z

  1. Attach to i current next-use- and liveness information of x, y, z

  2. Set x to 'not live' and 'no next use'

  3. Set y and z to 'live'
     Set 'next uses' of y and z to i

Determining Next-Use Information (Example)

1)	t = a - b	NU(t) = ...	NU(a) = ...	NU(b) = ...
2)	u = a - c	NU(u) = ...	NU(a) = ...	NU(c) = ...
3)	v = t + v	NU(v) = ...	NU(t) = ...	
4)	a = d	NU(a) = ...	NU(d) = ...	
5)	d = v + u	NU(d) = ...	NU(v) = ...	NU(u) = ...

Assume all variables are non-temporary, and thus are live on exit

Next-Use information in symbol table:

	a	b	c	d	t	u	v
after line 5 (on exit)	·	·	·	·	·	·	·
before line 5				...			

·	=	live, but next use is not known
—	=	not live
i	=	next use in line i

Determining Next-Use Information (Example)

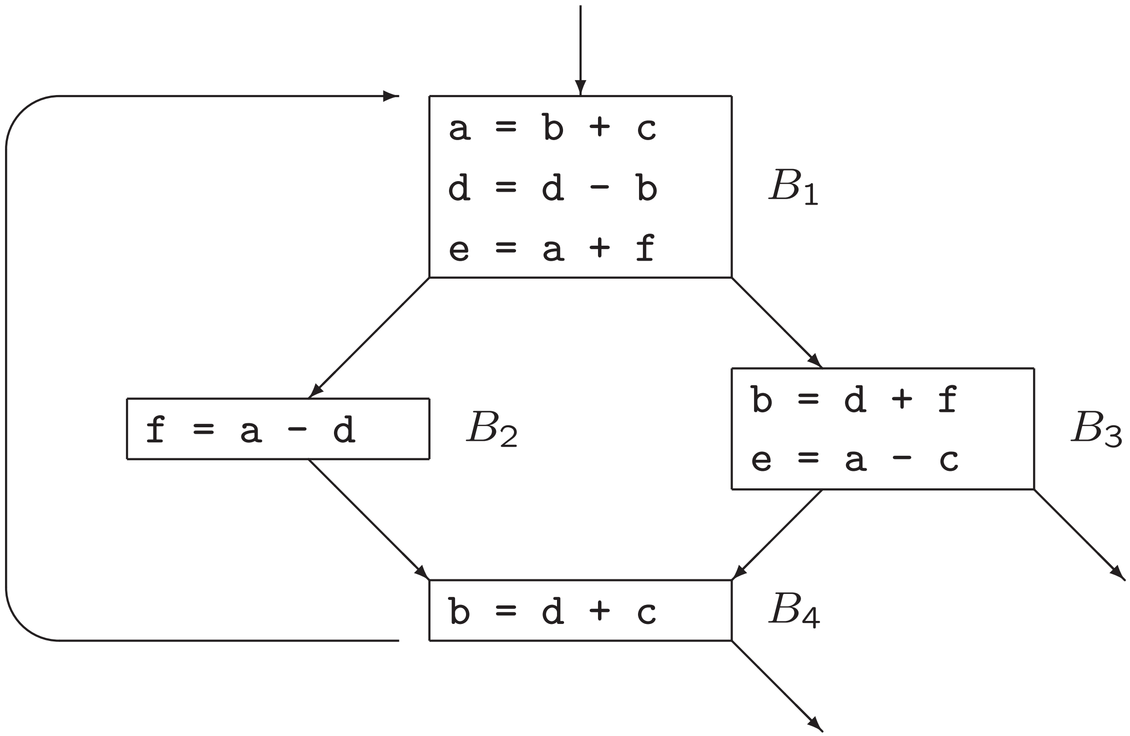
1)	t = a - b	NU(t) = 3 NU(a) = 2 NU(b) = ·
2)	u = a - c	NU(u) = 5 NU(a) = – NU(c) = ·
3)	v = t + v	NU(v) = 5 NU(t) = ·
4)	a = d	NU(a) = · NU(d) = –
5)	d = v + u	NU(d) = · NU(v) = · NU(u) = ·

	a	b	c	d	t	u	v
after line 5 (on exit)	·	·	·	·	·	·	·
before line 5	·	·	·	–	·	5	5
before line 4	–	·	·	4	·	5	5
before line 3	–	·	·	4	3	5	3
before line 2	2	·	2	4	3	–	3
before line 1 (on entry)	1	1	2	4	–	–	3

· = live, but next use is not known
– = not live
i = next use in line i

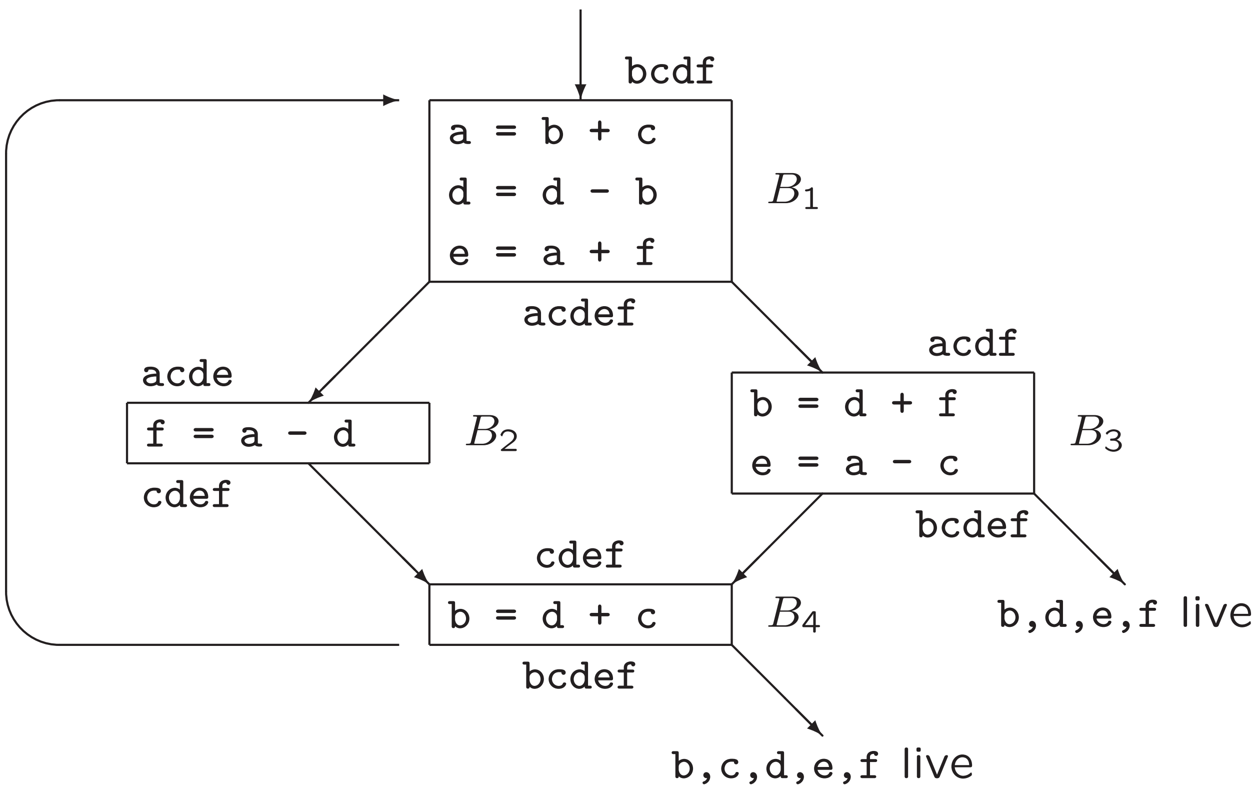
8.8.2 Passing Liveness Information over Blocks

Example of loop



Passing Liveness Information over Blocks

Example of loop



8.6 A Simple Code Generator

Use of registers

- Operands of operation must be in registers
- To hold values of temporary variables
- To hold (global) values that are used in several blocks
- To manage run-time stack

Assumption: subset of registers available for block

Machine instructions of form

- LD *reg*, *mem*
- ST *mem*, *reg*
- *OP* *reg*, *reg*, *reg*

8.6.1 Register and Address Descriptors

- Register descriptor keeps track of what is currently in register
  - Example:

$$\texttt{LD}\ R, x \qquad \rightarrow \text{register } R \text{ contains } x$$

  - Initially, all registers are empty

- Address descriptor keeps track of locations where current value of a variable can be found
  - Example:

$$\texttt{LD}\ R, x \qquad \rightarrow x \text{ is (also) in } R$$

  - Information stored in symbol table

8.6.2 The Code-Generation Algorithm

For each three-address instruction $x = y$ *op* z

1. Use *getReg*($x = y$ *op* z) to select registers R_x, R_y, R_z

2. If y is not in R_y, then issue instruction `LD` R_y, y',
   where y' is a memory location for y
   (according to address descriptor)

3. If z is not in R_z, . . .

4. Issue instruction *OP* R_x, R_y, R_z

Special case: $x = y$ . . .

At end of block: store all variables that are live-on-exit and not in their memory locations (according to address descriptor)

Managing Register / Address Descriptors

1. For the instruction `LD` R, x, ...

2. For the instruction `ST` x, R, ...

3. For an operation like `ADD` R_x, R_y, R_z, implementing $x = y + z$,

   (c) Remove R_x from addr. descr. of other variables

   (d) Remove x from reg. descr. of other registers

   (a) Change reg. descr. for R_x: only x

   (b) Change addr. descr. for x: only in R_x (not in x itself!)

4. For the copy statement $x = y$, ...

Managing Register / Address Descriptors

Example: $d = (a - b) + (a - c) + (a - c)$ $\quad a = \ldots$ old value of d

```
t = a - b
    LD  R1, a
    LD  R2, b
    SUB R2, R1, R2
u = a - c
    LD  R3, c
    SUB R1, R1, R3
v = t + u
    ADD R3, R2, R1
a = d
    LD  R2, d
d = v + u
    ADD R1, R3, R1

exit
    ST  a, R2
    ST  d, R1
```

R1	R2	R3

a	b	c	d	t	u	v
a	b	c	d			

Managing Register / Address Descriptors

Example: $d = (a - b) + (a - c) + (a - c)$ $\quad a = \ldots$old value of d

```
t = a - b
    LD  R1, a
    LD  R2, b
    SUB R2, R1, R2
u = a - c
    LD  R3, c
    SUB R1, R1, R3
v = t + u
    ADD R3, R2, R1
a = d
    LD  R2, d
d = v + u
    ADD R1, R3, R1

exit
    ST  a, R2
    ST  d, R1
```

R1	R2	R3
d	a	v

a	b	c	d	t	u	v
a,R2	b	c	d,R1			R3

8.6.3 Design of Function *getReg*

For each instruction $x = y$ *op* z

- To compute R_y
  1. If y is in register, $\longrightarrow R_y$
  2. Else, if empty register available, $\longrightarrow R_y$
  3. Else, select occupied register
     For each register R and variable v in R
     (a) If v is also somewhere else, then OK
     (b) If v is x, and x is not z, then OK
     (c) Else, if v is not used later, then OK
     (d) Else, `ST` v, R is required

     Take R with smallest number of stores

In fact, . . .

8.6.3 Design of Function *getReg*

For each instruction $x = y$ *op* z

- To compute R_y
  1. If y is in register, $\longrightarrow R_y$
  2. Else, if empty register available, $\longrightarrow R_y$
  3. Else, select occupied register
     For each register R and variable v in R
     (a) If v is also somewhere else, then OK
     (b) If v is x, and x is not z, then OK
     (c) Else, if v is not used later, then OK
     (d) Else, `ST` v, R is required

     Take R with smallest number of stores

- To compute R_x, similar with few differences (which?)

8.6.3 Design of Function *getReg*

For each instruction $x = y \; op \; z$

- To compute R_x
  1. If x is only value in register, $\longrightarrow R_x$
     (also if x is y or z)
  2. Else, if empty register available, $\longrightarrow R_x$
  3. Else, select occupied register
     For each register R and variable v in R
     (a) If v is also somewhere else, then OK
         (e.g., if v is y or z, just loaded)
     (b) If v is x (also if x is y or z), then OK
     (c) Else, if v is not used later, then OK
         (v might also be y or z)
     (d) Else, ST v, R is required

     Take R with smallest number of stores

Design of Function *getReg*

For each instruction $x = y$, choose $R_x = R_y$

Exercise 1

Addressing Modes of Target Machine

Form	Address	Example
x	x	`LD R1, x`
$a(r)$	$a + \mathit{contents}(r)$	`LD R1, a(R2)`
$\#c$		`LD R1, #100`

8.8 Register Allocation and Assignment

So far, live variables in registers are stored at end of block

Use of registers

- Operands of operation must be in registers
- To hold values of temporary variables
- To hold (global) values that are used in several blocks
- To manage run-time stack

8.8.2 Usage counts

With x in register during loop L

- Save . . . for . . . use of x that is not preceded by assignment in same block

- Save . . . for each block, where x is assigned a value and x is live on exit

-

$$\text{Total savings} \approx \sum_{\text{blocks } B \in L} \ldots$$

Choose variables x with largest savings

Usage counts

With x in register during loop L

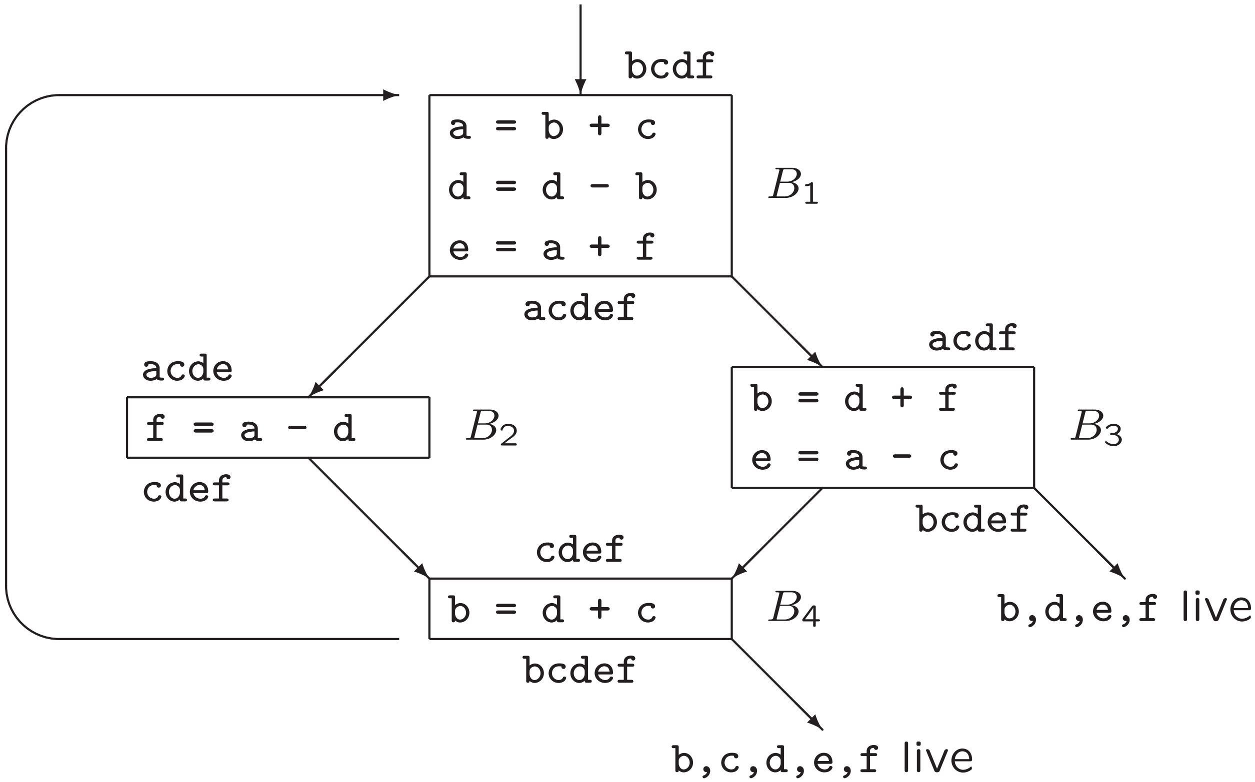
- Save 1 for each use of x that is not preceded by assignment in same block

- Save 2 for each block, where x is assigned a value and x is live on exit

- $$\text{Total savings} \approx \sum_{\text{blocks } B \in L} \mathit{use}(x,B) + 2 * \mathit{live}(x,B)$$

Choose variables x with largest savings

Usage counts (Example)



Savings for a are $1 + 1 + 1 * 2 = 4$

Komende week

- Vrijdag 18 november, 11:15–13:00: practicum

- Dinsdag 22 november: inleveren opdracht 3

- Woensdag 23 november, 11:15–13:00: hoorcollege + introductie opdracht 4 (inleveren 8 december)

- Woensdag 23 november, 13:45–. . . : werkcollege

Compiler constructie

college 8

Storage Organization

Code Generation

Chapters for reading:

7.1, 7.2–7.2.3

8.intro, 8.1, 8.2, 8.4, 8.6, 8.8–8.8.2