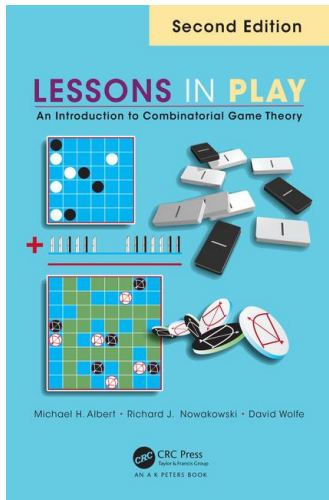


Lessons In Play

Canonical Form: Reducing Games

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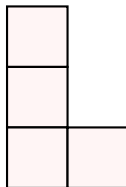
Left plays:



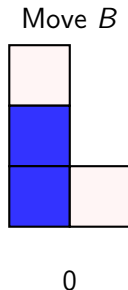
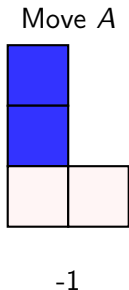
Right plays:



- ▶ Consider the board:
- ▶ Left moves first



Domination in Domineering



- ▶ Left always prefers B over A
- ▶ Move B *dominates* A
- ▶ Simplify a game through One-Hand-Tied principle

- ▶ Given the games:

$$G = \{A, B, C \dots \mid H, I, J, \dots\}$$

$$G' = \{B, C \dots \mid H, I, J, \dots\}$$

- ▶ If $B \geq A$ then $G = G'$
- ▶ Idem for Right: if $H \leq I$ remove I

Recall

For games G and H

$$G > H \iff G - H \in \mathcal{L}$$

$$G < H \iff G - H \in \mathcal{R}$$

Recall

For games G and H

$$G = H \iff G - H = 0$$

$$\iff G - H \in \mathcal{P}$$

Theorem 4.33

Proof

Show: if $B \geq A$, then $G - G' = 0 \iff \mathbf{G} + (-\mathbf{G}') = 0$

$$\underbrace{\{A, B, C, \dots \mid H, I, J, \dots\}}_{\mathbf{G}} + \underbrace{\{-H, -I, -J, \dots \mid -B, -C, \dots\}}_{-\mathbf{G}'}$$

Recall

For a game $G = \{ \mathcal{G}^L \mid \mathcal{G}^R \}$:
 $-G = \{ -\mathcal{G}^R \mid -\mathcal{G}^L \}$

Theorem 4.33

Proof

Show: if $B \geq A$, then $G - G' = 0 \iff \mathbf{G} + (-\mathbf{G}') = 0$

$$\overbrace{\{\overset{\checkmark}{A}, B, C, \dots \mid H, I, J, \dots\}}^{\mathbf{G}} \quad + \quad \overbrace{\{-H, -I, -J, \dots \mid -B, -C, \dots\}}^{-\mathbf{G}'}$$

- ▶ Left moves to A in G
Right moves to $-B$ in G'
 $A - B \leq 0$

Theorem 4.33

Proof

Show: if $B \geq A$, then $G - G' = 0 \iff \mathbf{G} + (-\mathbf{G}') = 0$

$$\overbrace{\{\check{A}, \check{B}, \check{C}, \dots \mid \check{H}, \check{I}, \check{J}, \dots\}}^{\mathbf{G}} \quad + \quad \overbrace{\{\check{-H}, \check{-I}, \check{-J}, \dots \mid \check{-B}, \check{-C}, \dots\}}^{-\mathbf{G}'}$$

- ▶ Left moves to A in G
Right moves to $-B$ in G'
 $A - B \leq 0$

- ▶ First player moves to m in G ,
Second player moves to $-m$ in G'
 $m - m = 0$

- ▶ Idem for any m in G'

Theorem 4.33

Proof

Show: if $B \geq A$, then $G - G' = 0 \iff \mathbf{G} + (-\mathbf{G}') = 0$

$$\overbrace{\{\overset{\checkmark}{A}, \overset{\checkmark}{B}, \overset{\checkmark}{C}, \dots \mid \overset{\checkmark}{H}, \overset{\checkmark}{I}, \overset{\checkmark}{J}, \dots\}}^{\mathbf{G}} \quad + \quad \overbrace{\{\overset{\checkmark}{-H}, \overset{\checkmark}{-I}, \overset{\checkmark}{-J}, \dots \mid \overset{\checkmark}{-B}, \overset{\checkmark}{-C}, \dots\}}^{-\mathbf{G}'}$$

- ▶ Left moves to A in G
Right moves to $-B$ in G'
 $A - B \leq 0$
- ▶ First player moves to m in G ,
Second player moves to $-m$ in G'
 $m - m = 0$
- ▶ Idem for any m in G'
- ▶ **All Second Player Wins**

Reversibility

Imagine:

- ▶ A game where Right has some advantage
- ▶ Left chooses some “good” move A
- ▶ But then Right has the move A^R , which reverses A and Right regains the advantage
- ▶ Left “actually” chooses A^R

Definition: Reversible Move

Let us formalize:

- ▶ Given a game $G = \{ A, \dots \mid \dots \}$
- ▶ $A = \{ \mathcal{A}^L \mid \mathcal{A}^R \}$
- ▶ We call A reversible $\iff \exists A^R \in \mathcal{A}^R : A^R \leq G$
- ▶ A^R is called a reversing move
- ▶ One-Hand-Tied: If Left moves to A , Right responds with A^R

Example: Reversible Move

We have a game $G = \{ \{ 2 \mid \{-8 \mid -9\} \} \mid -1 \}$

▶ $A = \{ 2 \mid \{-8 \mid -9\} \}$

▶ $A^R = \{-8 \mid -9\}$

Example: Reversible Move

We have a game $G = \{ \{ 2 \mid \{-8 \mid -9\} \} \mid -1 \}$

▶ $A = \{ 2 \mid \{-8 \mid -9\} \}$

▶ $A^R = \{-8 \mid -9\}$

$A^R \leq G?$

▶ $G - \{-8 \mid -9\} = G + \{9 \mid 8\}$

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$A^R \leq G?$

▶ $G - \{ -8 \mid -9 \} = G + \{ 9 \mid 8 \}$

▶ Left can win moving first:

$$\overset{L}{\mapsto} G + 9 \overset{R}{\mapsto} -1 + 9 > 0$$

Example: Reversible Move

We have a game $G = \{ \{ 2 \mid \{ -8 \mid -9 \} \} \mid -1 \}$

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$A^R \leq G$?

▶ $G - \{ -8 \mid -9 \} = G + \{ 9 \mid 8 \}$

▶ Left can win moving first:

$$\stackrel{L}{\mapsto} G + 9 \stackrel{R}{\mapsto} -1 + 9 > 0$$

▶ Right cannot win moving first:

$$\stackrel{R}{\mapsto} -1 + \{ 9 \mid 8 \} \stackrel{L}{\mapsto} -1 + 9 > 0$$

and

$$\stackrel{R}{\mapsto} G + 8 \stackrel{L}{\mapsto} \{ 2 \mid \{ -8 \mid -9 \} \} + 8 \stackrel{R}{\mapsto} \{ -8 \mid -9 \} + 8$$

$$\stackrel{L}{\mapsto} -8 + 8 = 0$$

Example: Reversible Move

We have a game $G = \{ \{ 2 \mid \{ -8 \mid -9 \} \} \mid -1 \}$

- ▶ $G - \{ -8 \mid -9 \} \in \mathcal{L}$
- ▶ $G - \{ -8 \mid -9 \} > 0$
- ▶ $\{ -8 \mid -9 \} < G$
- ▶ $\{ 2 \mid \{ -8 \mid -9 \} \}$ is *reversible*

Example: Reversible Move

We have a game $G = \{ \{ 2 \mid \{ -8 \mid -9 \} \} \mid -1 \}$

- ▶ $G - \{ -8 \mid -9 \} \in \mathcal{L}$
- ▶ $G - \{ -8 \mid -9 \} > 0$
- ▶ $\{ -8 \mid -9 \} < G$
- ▶ $\{ 2 \mid \{ -8 \mid -9 \} \}$ is *reversible*

We could then say

- ▶ If Left moves to $\{ 2 \mid \{ -8 \mid -9 \} \}$,
Right will move to $\{ -8 \mid -9 \}$
- ▶ When considering $\{ 2 \mid \{ -8 \mid -9 \} \}$,
we can replace it with -8
- ▶ $G = \{ -8 \mid -1 \}$

Given

- ▶ A game $G = \{ A, B, C, \dots \mid H, I, J, \dots \}$
- ▶ A reversible option A
- ▶ Some right option of A : $A^R = \{ W, X, Y, \dots \mid \dots \}$

Recall

$$A = \{ \mathcal{A}^L \mid \mathcal{A}^R \}$$

$$A^R \in \mathcal{A}^R : A^R \leq G$$

Given

- ▶ A game $G = \{ A, B, C, \dots \mid H, I, J, \dots \}$
- ▶ A reversible option A
- ▶ Some right option of A : $A^R = \{ W, X, Y, \dots \mid \dots \}$

Then $G = G'$

With $G' = \{ W, X, Y, \dots, B, C, \dots \mid H, I, J, \dots \}$

Recall

$$A = \{ \mathcal{A}^L \mid \mathcal{A}^R \}$$

$$A^R \in \mathcal{A}^R : A^R \leq G$$

Theorem 4.34

Proof

$$\overbrace{\{A, \check{B}, \check{C}, \dots \mid \check{H}, \check{I}, \check{J}, \dots\}}^G + \overbrace{\{\check{H}, \check{I}, \check{J}, \dots \mid -W, -X, -Y, \dots, -\check{B}, -\check{C}, \dots\}}^{-G'}$$



R

$G - W$

L wins:
 $\mathbf{G} \geq \mathbf{A}^R$

For the first move:

- ▶ Respond to all double moves with their counterpart
- ▶ If Right starts with $-W$, Left wins going second

Recall

$$A^R = \{ W, X, Y, \dots \mid \dots \}$$

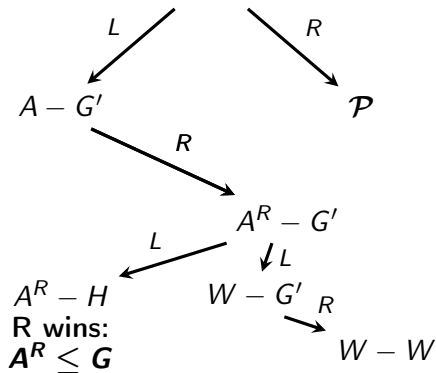
Note

$(-)W$ represents $(-)\{ W, X, Y, \dots \}$

Theorem 4.34

Proof

$$\overbrace{\{A, B, C, \dots \mid H, I, J, \dots\}}^G + \overbrace{\{\checkmark H, \checkmark I, \checkmark J, \dots \mid -\checkmark W, -\checkmark X, -\checkmark Y, \dots, -\checkmark B, -\checkmark C, \dots\}}^{-G'}$$



Recall

$$A^R = \{ W, X, Y, \dots \mid \dots \}$$

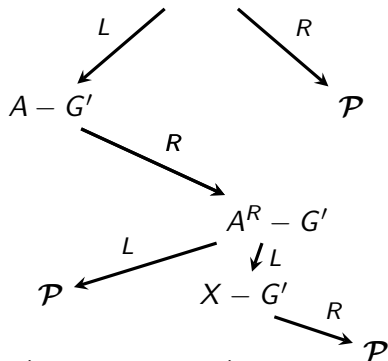
Note

$(-)H$ represents $(-)\{ H, I, J, \dots \}$

Theorem 4.34

Proof

$$\overbrace{\{A, B, C, \dots \mid H, I, J, \dots\}}^G + \overbrace{\{-H, -I, -J, \dots \mid -W, -X, -Y, \dots, -B, -C, \dots\}}^{-G'}$$



- ▶ $G - G' \in \mathcal{P}$, so $G - G' = 0$
- ▶ $\{A, B, C, \dots \mid H, I, J, \dots\} = \{W, X, Y, \dots, B, C, \dots \mid H, I, J, \dots\}$

Reduction Rules

Reduction Rule: Domination

*If we have some game \mathbf{G} where a move \mathbf{A} is **dominated** by another move from the same player.*

We may remove \mathbf{A} from \mathbf{G} .

Reduction Rule: Reversible Moves

*If we have some game \mathbf{G} where a player has a move \mathbf{A} that is **reversible**, where \mathbf{A}^R is a **reversing move**.*

\mathbf{A} may be replaced by the moves that this player would have in \mathbf{A}^R .

Example 4.42

$$G = \{ \{ 4 \mid 1 \}, 0 \mid 2 \}$$

Domination:

- ▶ $\{ 4 \mid 1 \} \in \mathcal{L}$, so $\{ 4 \mid 1 \} > 0$
- ▶ $\{ 4 \mid 1 \}$ dominates 0
- ▶ $G = \{ \{ 4 \mid 1 \} \mid 2 \}$

Example 4.42

$$G = \{ \{ 4 \mid 1 \} \mid 2 \}$$

Reversing option for Left:

▶ $A = \{ 4 \mid 1 \}$, $A^R = 1 = \{ 0 \mid \}$

▶ $G - A^R = \{ \{ 4 \mid 1 \} \mid 2 \} + \{ \mid 0 \}$

Example 4.42

$$G = \{ \{ 4 \mid 1 \} \mid 2 \}$$

Reversing option for Left:

- ▶ $A = \{ 4 \mid 1 \}$, $A^R = 1 = \{ 0 \mid \}$
- ▶ $G - A^R = \{ \{ 4 \mid 1 \} \mid 2 \} + \{ \mid 0 \}$
- ▶ $\stackrel{R}{\mapsto} \{ \{ 4 \mid 1 \} \mid 2 \} + 0 \Rightarrow$ Left wins going second
- ▶ $\stackrel{R}{\mapsto} 2 + \{ \mid 0 \} = 1 \Rightarrow$ Left wins going second

Example 4.42

$$G = \{ \{ 4 \mid 1 \} \mid 2 \}$$

Reversing option for Left:

- ▶ $A = \{ 4 \mid 1 \}$, $A^R = 1 = \{ 0 \mid \}$
- ▶ $G - A^R = \{ \{ 4 \mid 1 \} \mid 2 \} + \{ \mid 0 \}$
- ▶ $\overset{R}{\mapsto} \{ \{ 4 \mid 1 \} \mid 2 \} + 0 \Rightarrow$ Left wins going second
- ▶ $\overset{R}{\mapsto} 2 + \{ \mid 0 \} = 1 \Rightarrow$ Left wins going second
- ▶ Left wins going second at the least
- ▶ $\{ 0 \mid \} \leq G$, so $G = \{ 0 \mid 2 \}$

Example 4.42

$$G = \{ 0 \mid 2 \}$$

Reversing option, now for Right:

▶ $G = \{ 0 \mid \{ 1 \mid \} \}$

▶ $A = \{ 1 \mid \}, A^R = 1 = \{ 0 \mid \}$

Example 4.42

$$G = \{ 0 \mid 2 \}$$

Reversing option, now for Right:

- ▶ $G = \{ 0 \mid \{ 1 \mid \} \}$
- ▶ $A = \{ 1 \mid \}, A^R = 1 = \{ 0 \mid \}$
- ▶ Remember: $\{ 0 \mid \} \leq G$
- ▶ $G = \{ 0 \mid \} = 1$

Conclusion

- ▶ Through the application of these rules, we can reduce the expressions of a game.
- ▶ More complex game expressions might now actually be readable

Reduction Rule: Domination

*If we have some game \mathbf{G} where a move \mathbf{A} is **dominated** by another move from the same player.*

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Reduction Rule: Reversible Moves

*If we have some game \mathbf{G} where a player has a move \mathbf{A} that is **reversible**, where \mathbf{A}^R is a **reversing move**.*

\mathbf{A} may be replaced by the moves that this player would have in \mathbf{A}^R .