PEBBLES at STACS





Marseille Feb 06

Automata with **Nested Pebbles** capture

FO Logic with **Transitive Closure**

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two classic characterizations

strings: [non]deterministic logarithmic space

Immerman

Multi-Head Automata re (two-way)

$$\varphi^*(\underline{x},\underline{y})$$

our STACS paper

on strings, trees, grids, toruses, mazes, ...

First-Order Logic Multi-Head Automata + transitive closure + 'nested pebbles'

 $\varphi^*(\underline{x},\underline{y})$

arity k k heads

starting with trees, on single head First-Order Logic Multi-Head Automata + 'nested pebbles' + transitive closure unary $\phi^*(x,y)$ trees k=1arity k k heads

background

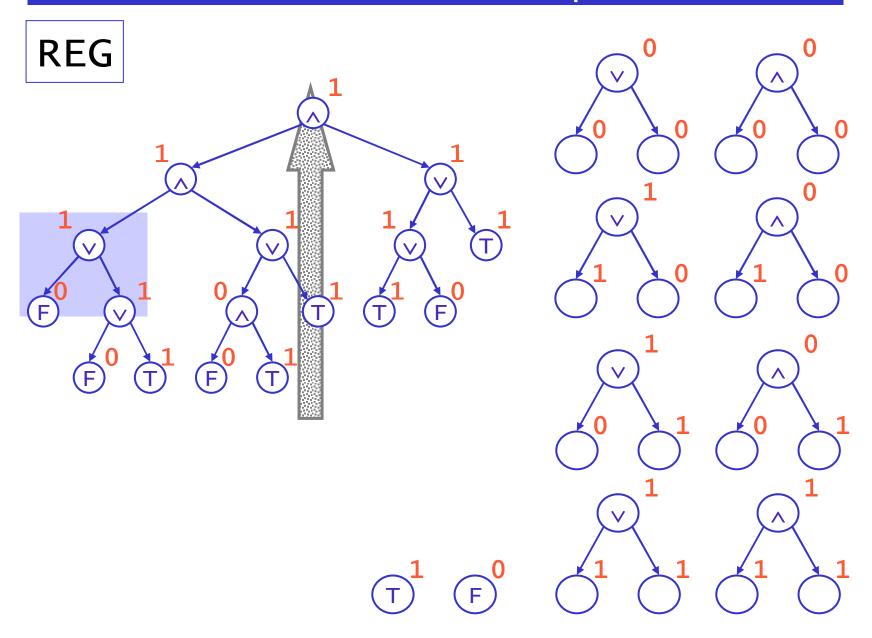
transitive closure
descriptive complexity

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descriptive complexity strings, trees, n-dim grids, ...
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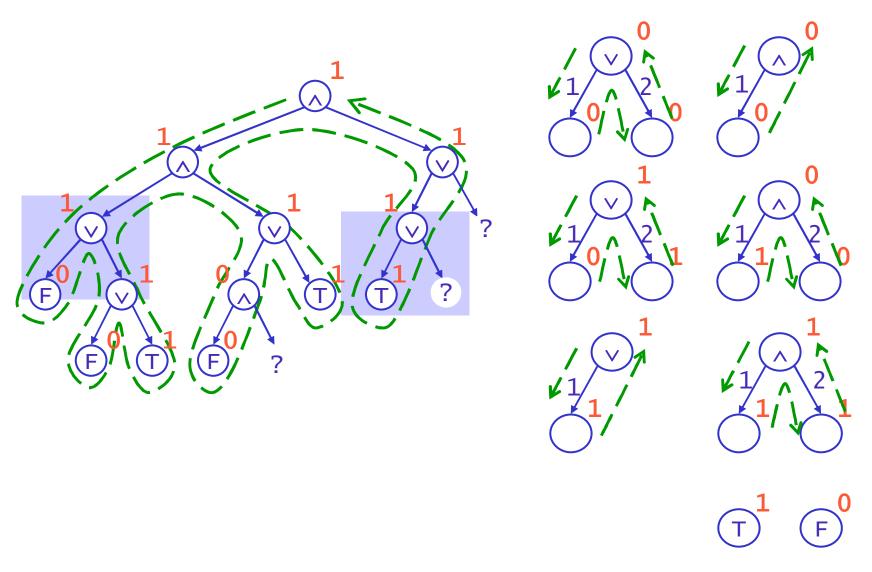
- XML document transformation single head on (unranked) trees
- graph exploration

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many heads on graphs 'robots' grids, toruses, mazes, ...
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bottom-up tree automata



walking along the tree

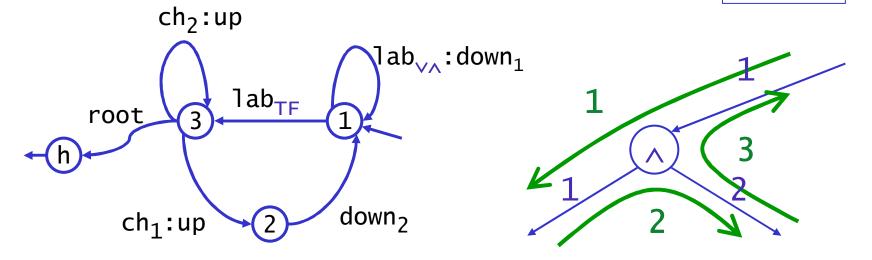


cf. two-way finite state automaton

tree walking automaton

example: tree traversal

TW

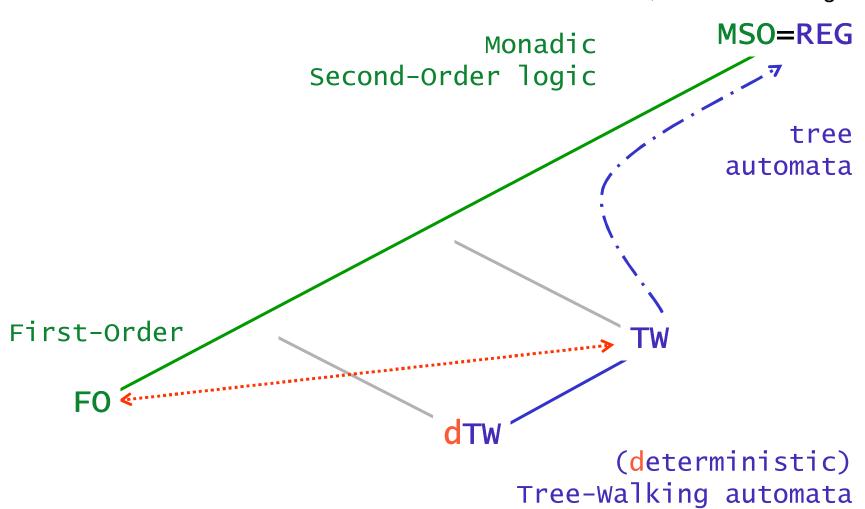


walk along edges, moves based on

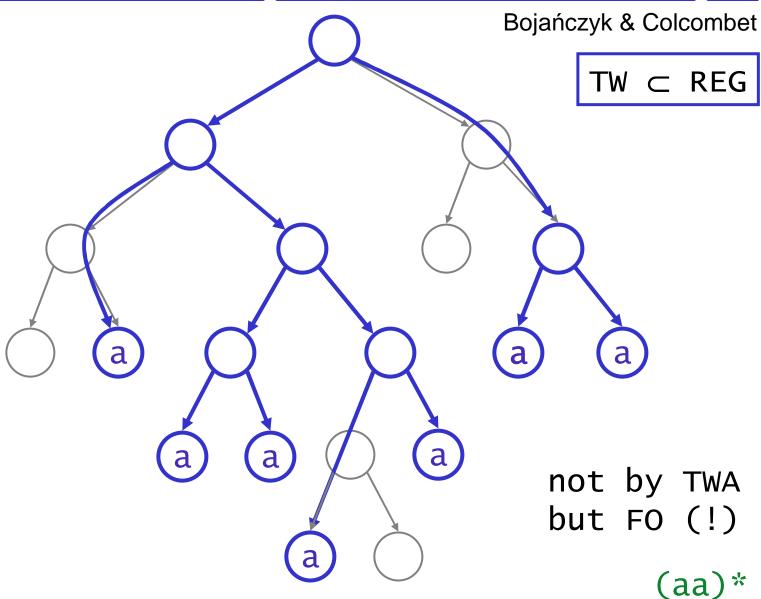
- state
- node label
- child number(= incoming edge)

single head on trees

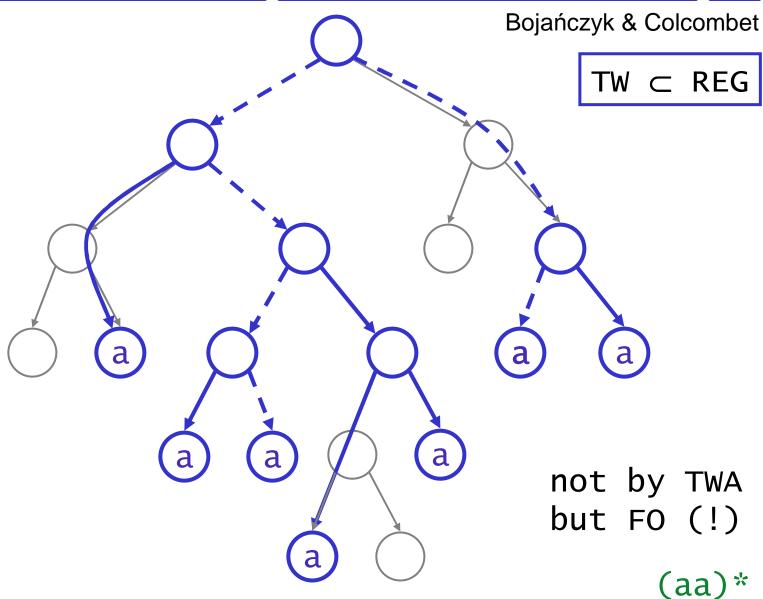
Doner; Thatcher & Wright



'branching structure' of even length

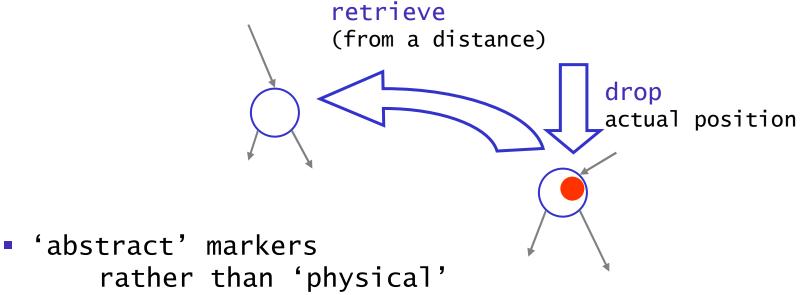


'branching structure' of even length



adding nested pebbles

pebble: marks a node

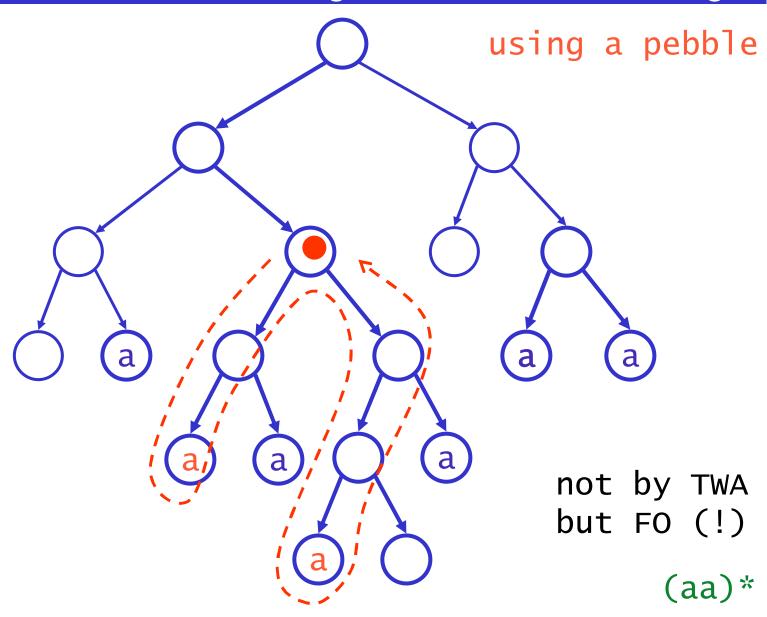


- nested lifetimes LIFO
- fixed number for automaton
- can be distinguished

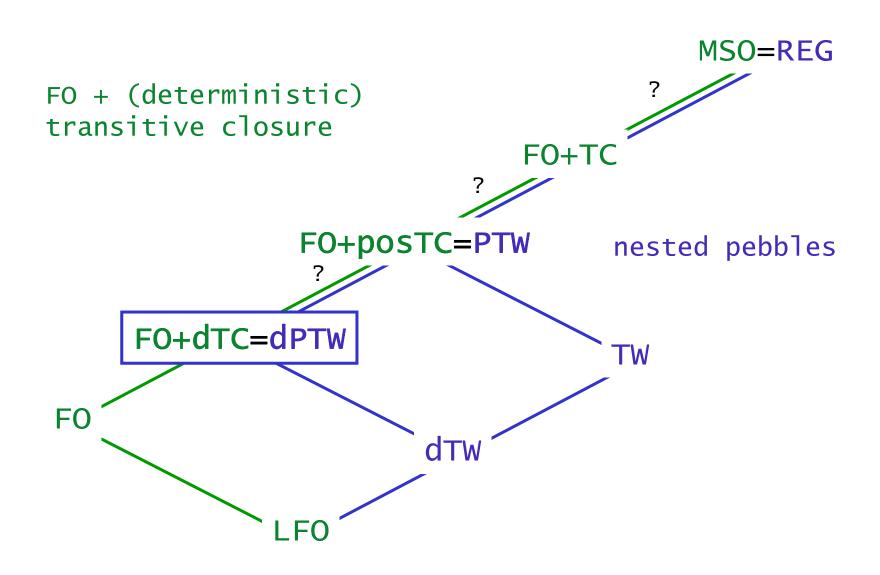
'regular' extension (for single head on trees)



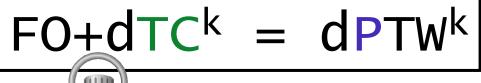
'branching structure' of even length

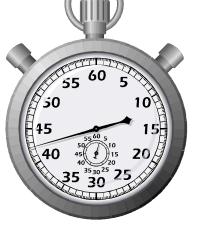


single head on trees



main result





proof summary
manager style

(1) logic to nested pebbles

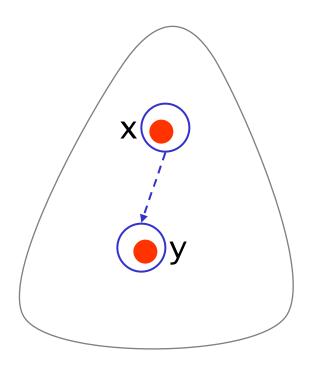
lab_a(x)
edg_i(x,y)

$$x \leq y$$

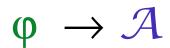
$$x = y$$

 $\neg \land \lor \\ \forall x \exists x$

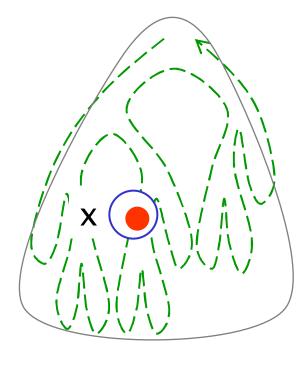
$$\phi^*(x,y)$$



$$x \leq y$$



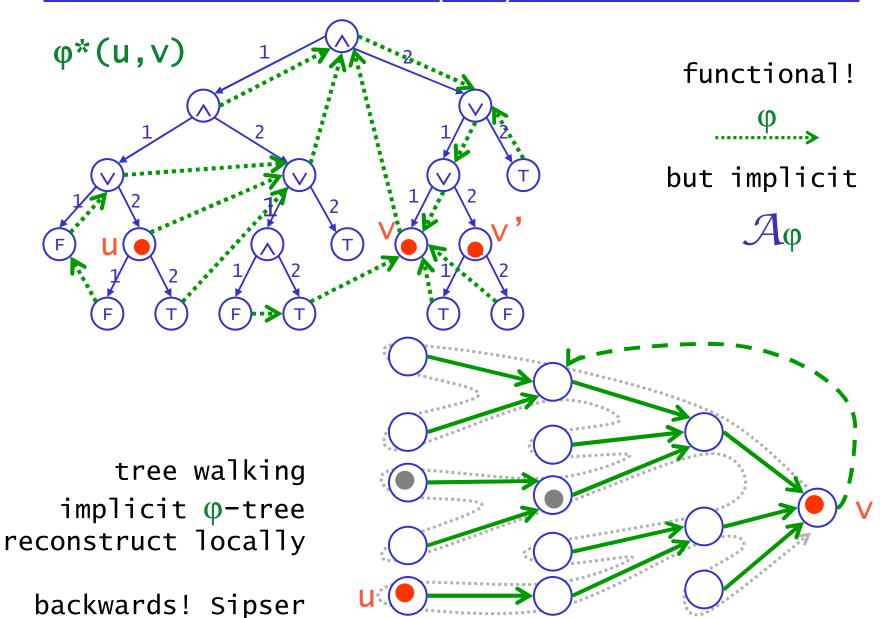
always halting free variables ~ fixed pebbles



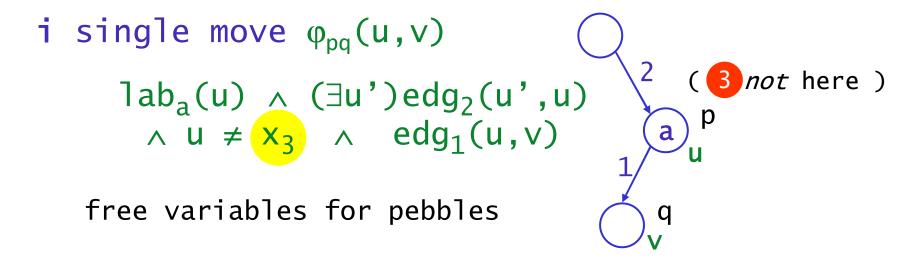
$$\forall x \ \phi(x)$$



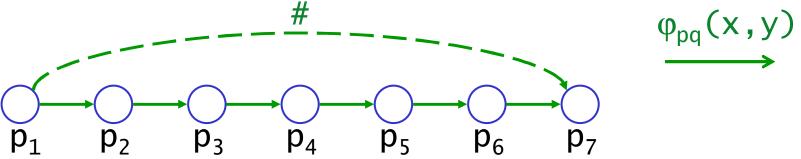
(1ctd) transitive closure



(2) nested pebbles to logic

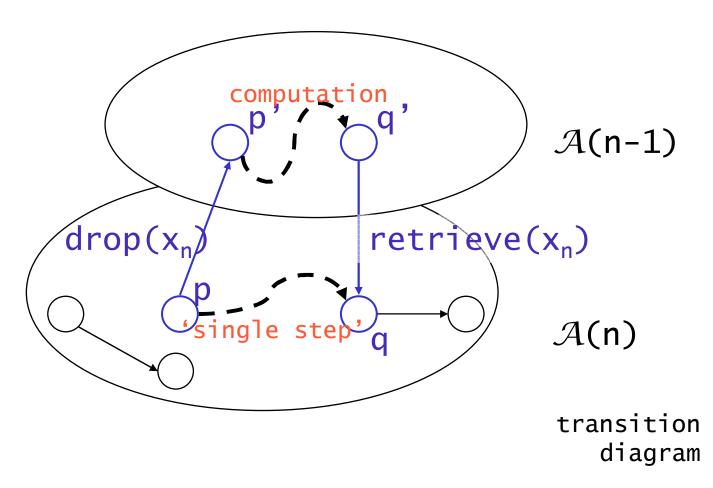


ii computation ~ tc with states



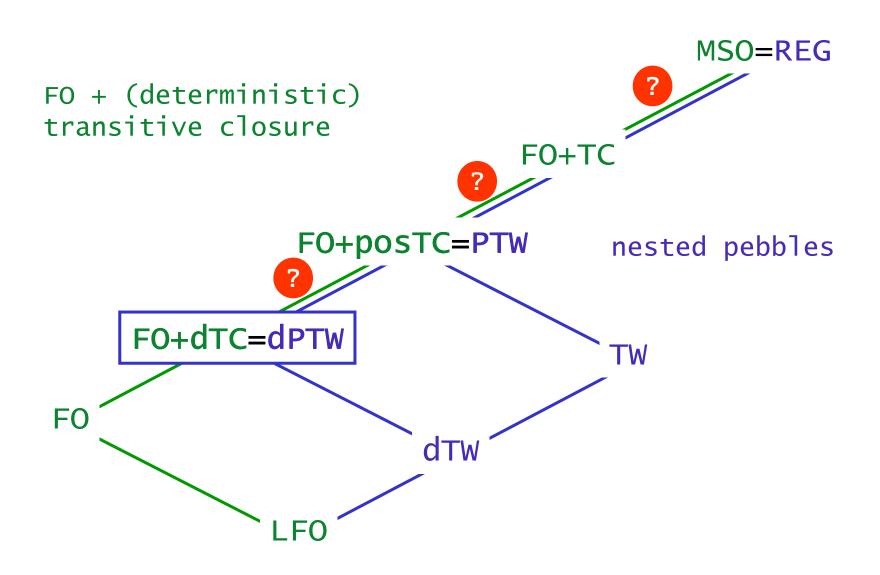
Kleene: removing states finite aut to reg expr

(2ctd) dropping pebbles

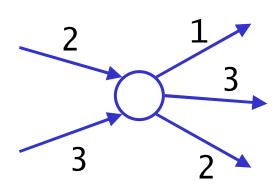


$$\phi_{pq}^{n}(u,v) = \phi_{p'q'}^{(n-1)\#}(u,v)$$
replacing x_n by u

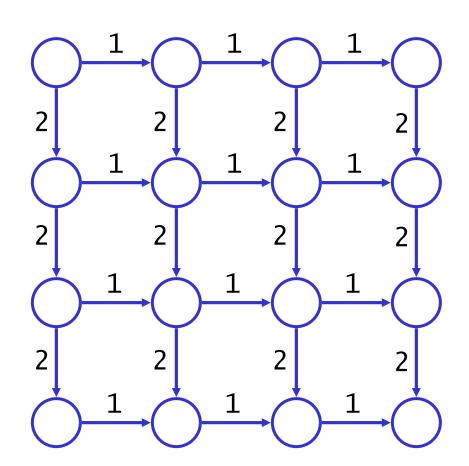
single head on trees



from trees to graphs



locally injective



grid, torus

nested pebbles to logic

$$X \leq Y$$

 $X = Y$

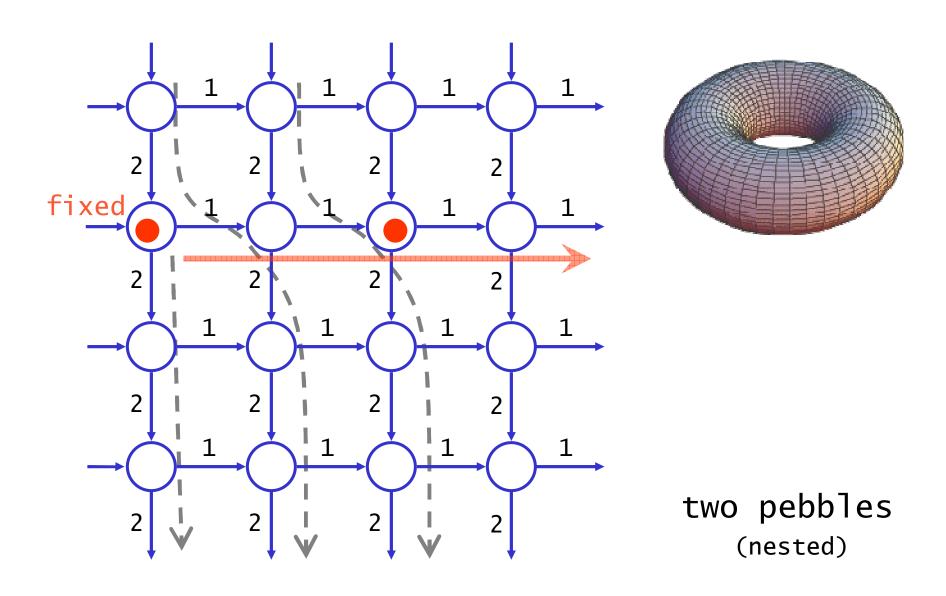
$$\neg \land \lor \\ \forall x \exists x$$

$$\phi^*(x,y)$$

$dPTW^k \subseteq FO+dTC^k$

for families of graphs (i.e. with fixed label alphabets)

walking the torus



graphs with a guide

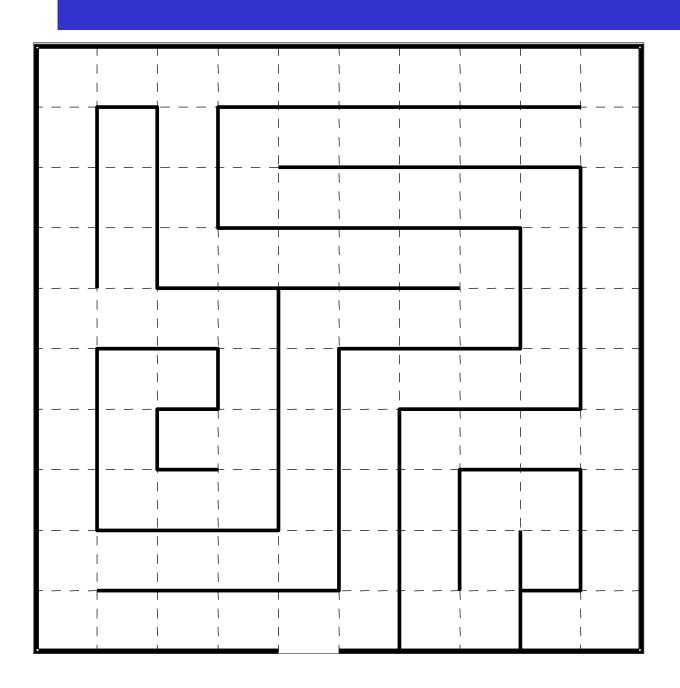
$FO+dTC^k = dPTW^k$

for families of *searchable* graphs with a 'guide'

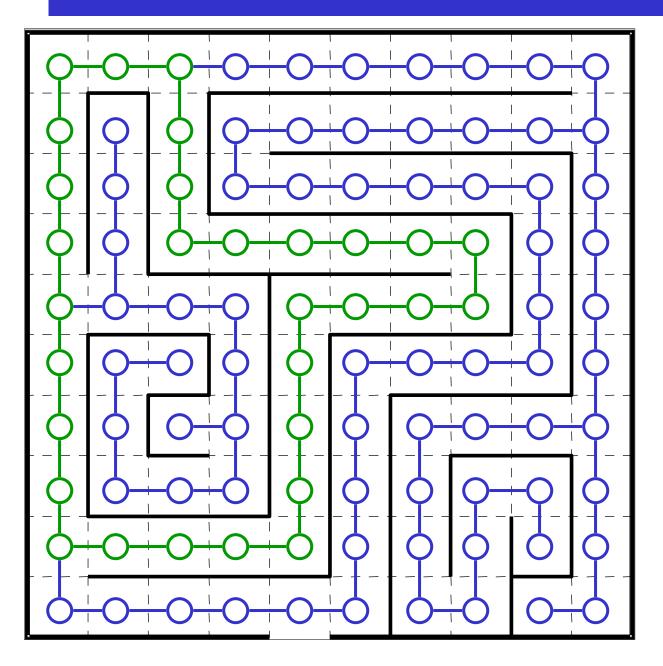
 $(\forall x) \ lab_0(x)$

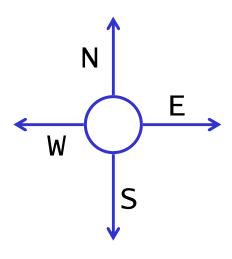
unranked trees, grids, toruses, ...
2 pebbles

mazes



mazes





Blum & Kozen

two heads!

(not nested)

searching with many heads

$$FO+dTC^k = dPTW^k$$

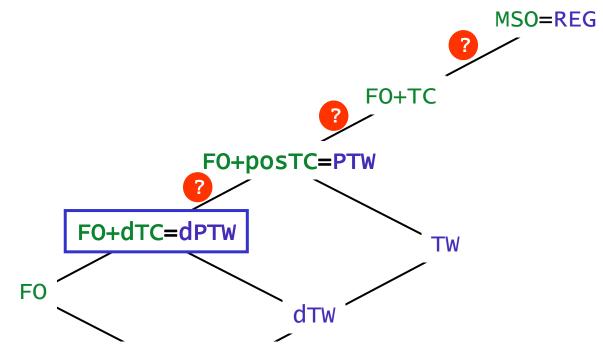
for families of *k-searchable* graphs

additional instruction
move head to pebble

Cook & Rackoff 'Jumping Automata' mazes
not all graphs

finally: work to do ...

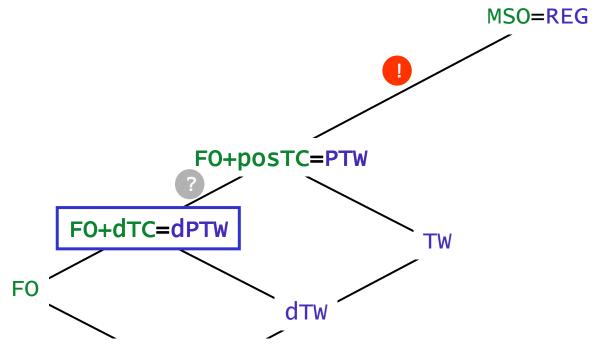
open for single head on trees:



- \bigcirc dPTW \subset PTW \subset REG
- \bigcirc F0+dTC \subset F0+posTC \subset F0+TC \subset MS0
- 🕜 pebble hierarchy
- type of pebbles physical vs. abstract
- 🕜 alternation

work to do ... until last week

Bojańczyk, Samuelides, Schwentick, Segoufin



- dPTW ⊆ PTW ⊂ REG
- \bigcirc F0+dTC \subseteq F0+posTC \subseteq F0+TC \subseteq MS0
- pebble hierarchy
- type of pebbles physical ⇔ abstract
- alternation

many heads? graphs?

