

- 1) a. Choose $p \hat{=}$ 'shines today', and $q \hat{=}$ 'shines tomorrow'. We get $p \rightarrow (\neg q)$.
 c. Choose $p \hat{=}$ 'shines today', and $r \hat{=}$ 'rains today'. We explicitly have to express *exclusive or*, $(p \vee r) \wedge \neg(p \wedge r)$.
 d. Both options, or only one of them?

- 2) 'In full' suggests to add also the outer parentheses we usually omit.

$$\begin{aligned} & ((\neg p) \rightarrow (p \wedge r)) \\ & ((p \rightarrow q) \rightarrow (r \rightarrow (s \vee t))) \\ & ((p \vee q) \rightarrow ((\neg p) \vee r)) \end{aligned}$$

$p \vee q \wedge r$ has no fixed meaning if we strictly follow the rules of the exercise. The binding order of \vee and \wedge is not specified, associativity seems to hold only for \rightarrow .

- 3) b. Sorry, I mean natural language 'and', to have two tables, to compare the two related formulas.

- c. We see from the table that $(q \wedge (r \rightarrow q))$ is actually equivalent to q .

p	q	r	$p \vee (\neg(q \wedge (r \rightarrow q)))$			
T	T	T	T	F	T	T
T	T	F	T	F	T	T
T	F	T	T	T	F	F
T	F	F	T	T	F	T
F	T	T	F	F	T	T
F	T	F	F	F	T	T
F	F	T	T	T	F	F
F	F	F	T	T	F	T
			4	3	2	1

- 4) a. $(\forall x)A(m, P(x))$ is not correct as ' $P(x)$ ' means the logical statement ' x is a professor' (true or false) and can not be written in place of a person 'professor x '.
 $(\forall x)(P(x) \rightarrow A(m, x))$ For every x , when professor, then Mary admires x
 This is *not* the same as $(\forall x)(P(x) \rightarrow A(m, x))$ which also states that all x are professors.
 We assume there are also students.

- b. $(\exists x)(P(x) \wedge A(x, m))$ There is a person who is professor and admires Mary.

- c. No student attended every lecture. Tricky. There is no single student that attended all lectures. $\neg(\exists x)(\forall y)(S(x) \wedge L(y) \wedge B(x, y))$
 Every student missed at least one lecture: $(\forall x)(S(x) \rightarrow (\exists y)(L(y) \wedge \neg B(x, y)))$

- 5) a. For every person there is a person that is the mother of the first. $(\forall x)(\exists y)M(y, x)$

- b. $(\forall x)(\exists y)M(y, x) \wedge (\forall x)(\exists y)F(y, x)$
 or $(\forall x)(\exists y)(\exists z)(M(y, x) \wedge M(z, x))$

- d. $(\exists x)(\exists y)(F(\text{Ed}, x) \wedge F(x, y))$

- e. Probably: $(\forall x)(\exists y)F(x, y) \rightarrow (\exists y)(F(x, y) \vee M(x, y))$

j. $H(\text{Ed}, \text{Patsy})$

k. Brother of husband of Monique, or husband of sister of Monique, assuming old-fashioned marriages.

6) In the first case ϕ is not true. If we choose $x = b$ and $y = a$ then $R(x, y)$ is true, as (b, a) is in the relation given. However we can not find $z \in A$ that makes $R(y, z)$ true, as there is no pair (b, z) in the relation. This makes $R(x, y) \rightarrow R(y, z)$ false (for this x , this y and all z). Consequently ϕ is false.

Second case. Assume x and y are chosen that make $R(x, y)$ true. We consider cases to find z .

$x = b, y = c$ then we choose $z = b$,

$x = a, y = b$ then we choose $z = c$,

$x = c, y = b$ then we choose $z = c$.

Hence we can always find z , which makes the statement true.