



Focs

Lecture 6



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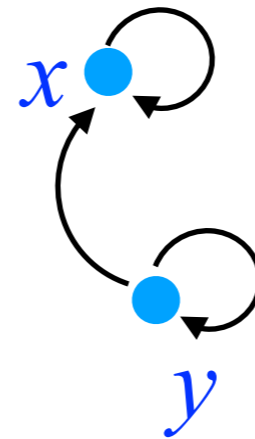
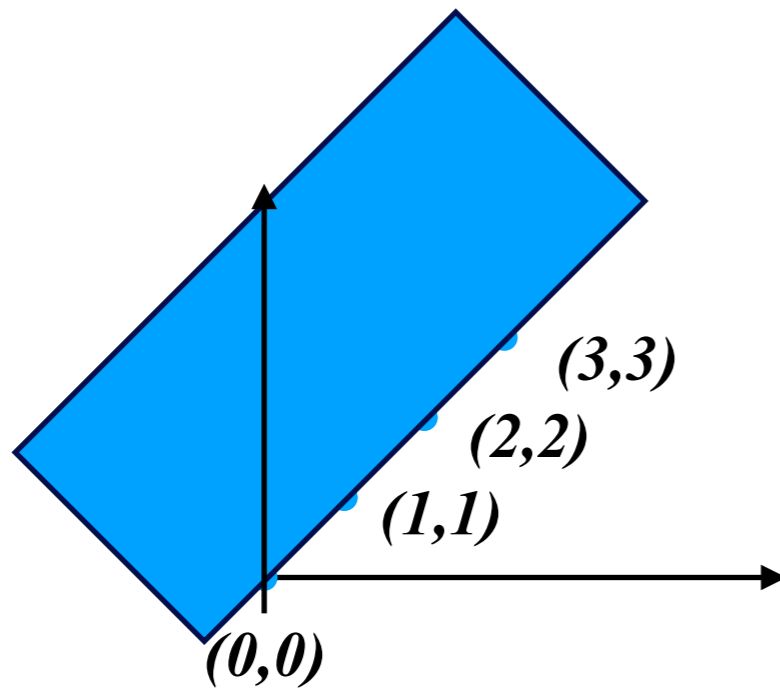
RELATIONS

CONTINUED

Properties of binary relations

Definition. Relation $R \in A \times A$ is **reflexive** if
For all x , xRx .

*If for all $x \in A$, $(x, x) \notin R$, then it is **irreflexive***

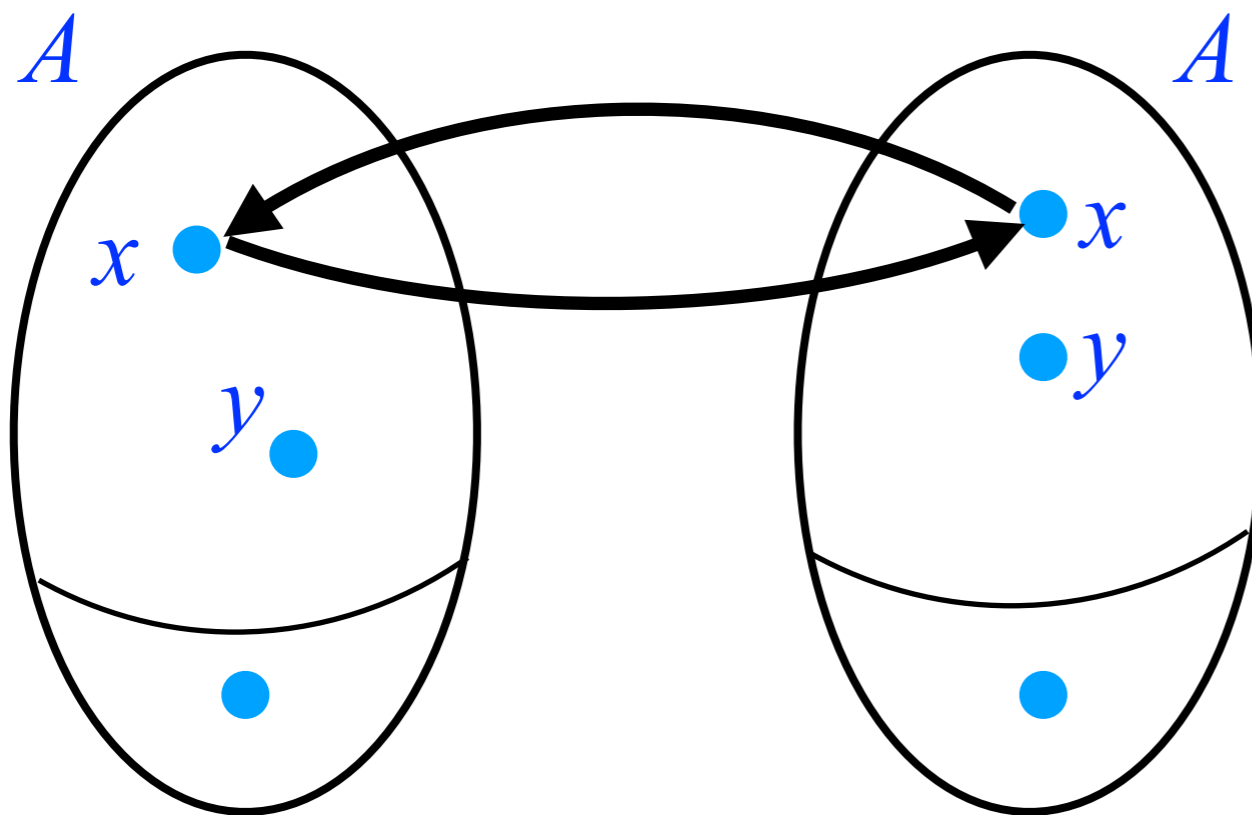


Properties of binary relations

Definition. Relation $R \in A \times A$ is **symmetric** if

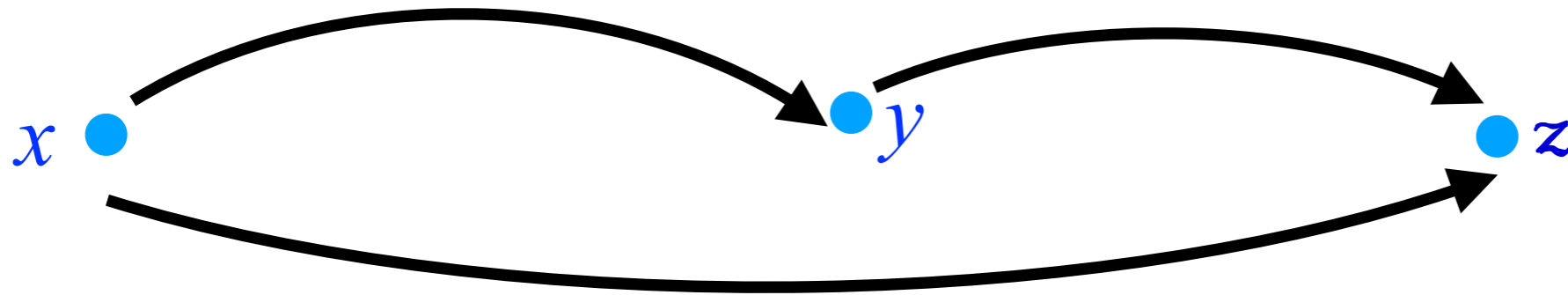
For all x, y , if xRy then yRx .

If xRy & yRx implies that $x=y$ then it is **antisymmetric**



Properties of binary relations

Definition. Relation $R \in A \times A$ is **transitive** if
For all x, y, z if xRy & yRz , then xRz





Extra clarification

not reflexive \neq irreflexive

not symmetric \neq antisymmetric

not reflexive: there exists x such that $(x, x) \notin R$

irreflexive: for all x $(x, x) \notin R$

Question: how would you express “not symmetric” formally, using “for all” etc...



Extra clarification

not reflexive \neq irreflexive

not symmetric \neq antisymmetric

not reflexive: there exists x such that $(x, x) \notin R$

irreflexive: for all x $(x, x) \notin R$

Question: how would you express “not symmetric” formally, using “for all” etc...

not symmetric: $\exists (x, y)$ s.t. $(x, y) \in R$ & $(y, x) \notin R$

N.B.: Symmetry does not mean ALL pairs-of-pairs $(x, y), (y, x)$ are in R .
It just means that if ONE of them is, then so is the other.



Extra clarification

not reflexive \neq irreflexive

not symmetric \neq antisymmetric

not symmetric: there exist x, y such that $(x, y) \in R$ & $(y, x) \notin R$

antisymmetric: for all x $(x, y) \in R$ & $(y, x) \in R$ then $x = y$

Not symmetric/reflexive: violates definition of ...

Antisymmetric: only symmetric pairs are reflexive ones

Extra clarification

Examples:

reflexive: \leq

not reflexive: “product of x and y is even” on integers

irreflexive: $<$, \subset

Question?

not reflexive $\Rightarrow \exists x$ s.t. $(x, x) \notin R$

xRy if $x \cdot y$ is even; xRx if x^2 is even. But $3^2 = 9 \dots$

IRREFLEXIVE: xRx never,

“Product of x, y is even” is neither reflexive nor irreflexive on integers

However, on evens it is reflexive (product of evens is even)



Extra clarification

Examples:

symmetric: =

not symmetric: <

antisymmetric: \leq

not antisymmetric: "is sibling of"; x, y such that $2x \geq y$ on \mathbb{N} (see below!)

IS $<$ ON \mathbb{Z} ANTISYMMETRIC?

Question?

YES, BY VACUOUS REASONS. NOTE, IMPLICATION "IF A THEN B" IS ONLY FALSE IF $A \ \& \ \text{NOT-}B$. IF A IS ALWAYS FALSE $A \rightarrow B$ IS TRUE.

$(x < y \ \& \ y < x) \Rightarrow x = y$ IS TRUE BECAUSE FOR NO x, y

IS $(x < y) \ \& \ (y < x)$ TRUE. THIS IS FALSE, SO $(x < y \ \& \ y < x) \Rightarrow (x = y)$ IS TRUE.



Extra clarification

Examples:

symmetric: =

not symmetric: <

antisymmetric: \leq

not antisymmetric: “is sibling of”; x, y such that $2x \geq y$ on \mathbb{N} (see below!)

IS $<$ ON \mathbb{Z} ANTISYMMETRIC?

Question?

$$xRy \Leftrightarrow 2x \geq y.$$

not antisymmetric \Rightarrow [IT IS NOT THE CASE THAT
IF xRy & yRx THEN $x=y$]

$$\Leftrightarrow \exists x, y \text{ st } xRy \text{ \& } yRx \text{ \underline{AND} } x \neq y.$$

NOTE $2 \cdot 3 \geq 4$ & $2 \cdot 4 \geq 3$ so $3R4$ & $4R3$ BUT $4 \neq 3$.



Extra example

Transitivity: $xRy \ \& \ yRz \Rightarrow xRz$

example (trivial): $<, =$

example (non-trivial) : empty relation (vacuous truth), IMPLICATION ITSELF.

counterexample: orthogonality

counterexample: two transitive relations

their composition is not transitive

$$R = \{(1,2), (3,4)\} \quad S = \{(2,3), (4,5)\}$$

WHY TRANSITIVE

Question?

$$R \circ S = \begin{matrix} 1 & 2 & 3 \\ 3 & 4 & 5 \end{matrix} = \{(1,3), (3,5)\} \quad (1,5) \text{ missing}$$

COMPOSITION OF TRANSITIVE RELATIONS MAY NOT BE TRANSITIVE

NOTE: ALL RELATIONS WHERE $(xRy \ \& \ yRz)$ NEVER HAPPENS ARE TRANSITIVE



Equivalence relations

*Definition. Binary relation $R \subseteq A^2$ is an **equivalence relation** if it is*

- *reflexive*
- *symmetric*
- *transitive*

Example: “=”

Slightly more complicated: aRb if $|a| = |b|$ [absolute value equality]

More complicated: “is congruent to (mod n)” [$a \equiv b \pmod{n} \Leftrightarrow (n \mid (b - a))$]

Counterexample: $<$

REMAINDER $a \bmod n = a - \lfloor \frac{a}{n} \rfloor$

$a \equiv b \pmod{n}$ if

$a \bmod n = b \bmod n$

EQUIV: $a \equiv b \pmod{n}$ if $(b - a) = k \times n$ for an integer k

↑
"Floor"
FIRST INTEGER LESS OR EQUAL TO a/n



Equivalence relations

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Example: “=”

Slightly more complicated: aRb if $|a| = |b|$ [absolute value equality]

More complicated: “is congruent to (mod n)” [$a \equiv b \pmod{n} \Leftrightarrow (n \mid (b - a))$]

Counterexample: $<$

$$aRb \Leftrightarrow a \equiv b \pmod{n}$$

$aRb \Leftrightarrow$ remainder of division of a & b
with n is the same.

$$a \equiv b \pmod{2} \quad \text{[“evenness”]}$$

$$10 \equiv 13 \pmod{3}$$



Equivalence relations

Definition. Binary relation $R \subseteq A^2$ is an equivalence relation if it is

- *reflexive*
- *symmetric*
- *transitive*

Equivalence relations = always capture some notion of equality
“the same up to ‘irrelevant’ properties”

Partial orders

Definition. Binary relation $R \subseteq A^2$ is a **partial order** if it is

- reflexive
- antisymmetric
- transitive

Example: \leq

More complicated: “divisibility” $b \mid a$, if $\frac{a}{b} \in \mathbb{Z}$ (*b divides a*)

Counterexample: $<$

$$\underline{a R b \Leftrightarrow b \mid a .}$$

$$(1) \quad a \mid a \quad \checkmark$$

$$(2) \quad a \mid b \ \& \ b \mid a \Rightarrow a = b \quad \checkmark$$

$$(3) \quad a R b, \ b R c \Rightarrow a R c$$

$$b \mid a, \ c \mid b \Rightarrow c \mid a$$

$$30 \mid 60 \quad 15 \mid 30 \Rightarrow 15 \mid 60 \quad \checkmark$$



Characterizations (1)

UNDERSTANDING
COMPOSITION

MAIN PROPERTIES U1A

- (1) R is reflexive $\Leftrightarrow id \subseteq R$
- (2) R is symmetric $\Leftrightarrow R^{-1} \subseteq R$
- (3) Transitive $\Leftrightarrow R \circ R \subseteq R$
(then also $R^{*n} \subseteq R$)

recall $id_A = \{ (a, a) \mid a \in A \}$
= set of all pairs (a, a)

Let's prove these

$id \in R$ means $\begin{cases} (x, y) \in id \Rightarrow (x, y) \in R \\ (x, y) \in id \Rightarrow (y, x) \in R \end{cases}$

(1) Reflexive $\Rightarrow \forall x (x, x) \in R \Rightarrow id \in R$

$id \in R \Rightarrow \forall x (x, x) \in R \Rightarrow R$ is reflexive

Characterizations (1)

- (1) R is reflexive $\Leftrightarrow id \subseteq R$
- (2) R is symmetric $\Leftrightarrow R^{-1} \subseteq R$
- (3) Transitive $\Leftrightarrow R \circ R \subseteq R$
(then also $R^{\circ n} \subseteq R$)

Let's prove these

(2) \Rightarrow

R is symm \Rightarrow $\underbrace{[xRy \Rightarrow yRx]}_{(c1)}$; need $(x,y) \in R^{-1} \Rightarrow (x,y) \in R$.

(2) $\Leftarrow R^{-1} \subseteq R$

$\Leftrightarrow \underbrace{(y,x) \in R \Rightarrow (x,y) \in R}_{(c1) \checkmark}$

\Rightarrow if $(x,y) \in R^{-1} \Rightarrow (x,y) \in R$ &
 $(y,x) \in R \Rightarrow (x,y) \in R$

Characterizations (1)

- (1) R is reflexive $\Leftrightarrow id \subseteq R$
- (2) R is symmetric $\Leftrightarrow R^{-1} \subseteq R$
- (3) Transitive $\Leftrightarrow R \circ R \subseteq R$
(then also $R^{*n} \subseteq R$)

Let's prove these

(3) \Rightarrow $\underbrace{[xRy \ \& \ yRz \Rightarrow xRz]}_{(C_1)}$;

\Leftarrow ASSUME THAT $[(a,c) \in R \circ R \Rightarrow \exists y (a,y) \in R, (y,c) \in R] \Rightarrow (a,c) \in R$

Need: $\underbrace{(a,y) \ \& \ (y,c) \in R}_{(C_1)} \Rightarrow (a,c) \in R$

But $(a,c) \in R \circ R$ so $\Rightarrow (a,c) \in R !!$

need $\underbrace{(a,c) \in R \circ R}_{(C_1)} \Rightarrow (a,c) \in R$
 $\exists y (a,y) \in R \ \& \ (y,c) \in R$
 Let $a \ \& \ c$ are such that
 that $\exists y$ for which
 $(a,y) \ \& \ (y,c)$ are in R
 But then by (C_1) $(a,c) \in R$. \checkmark

Characterizations (2)

- (1) R is irreflexive $\Leftrightarrow id \cap R = \emptyset$
 (2) R is antisymmetric $\Leftrightarrow R^{-1} \cap R \subseteq id$

Let's prove these

(1) \Rightarrow

$$\forall x \underbrace{(x, x) \notin R}_{(c1)}$$

$$(a, b) \in id \cap R \Rightarrow (c1)$$

$$(a, b) \in id \Rightarrow (a = b) \Rightarrow (a, b) \notin R$$

$$\Rightarrow (a, b) \notin id \cap R \quad \forall (a, b)$$

$$\Rightarrow \emptyset = id \cap R$$

$$\Leftarrow \underbrace{id \cap R = \emptyset}_{(c1)}, \quad \forall (a, a) \in id \stackrel{(c1)}{\Rightarrow} \underbrace{(a, a) \notin R}_{\Leftarrow \text{irreflexive}}$$



Characterizations (2)

- (1) R is irreflexive $\Leftrightarrow id \cap R = \emptyset$
- (2) R is antisymmetric $\Leftrightarrow R^{-1} \cap R \subseteq id$

Let's prove these

$$(2) \Leftarrow \text{IF } (x,y) \in R \text{ \& } (y,x) \in R \Rightarrow (x,y) \in id \Rightarrow x=y. \quad \checkmark$$
$$\Rightarrow yRx \text{ \& } xRy \Rightarrow x=y ; \quad (x,y) \in R \text{ \& } (y,x) \in R \Rightarrow x=y \Rightarrow (x,y) = (x,x) \in id.$$

Characterizations (3)

- (1) $R \subseteq A \times B$ is functional $\Leftrightarrow R^{-1} \circ R \subseteq id_B$
 (2) $R \subseteq A \times B$ is surjective $\Leftrightarrow id_B \subseteq R^{-1} \circ R$
 (3) $R \subseteq A \times B$ is injective $\Leftrightarrow R \circ R^{-1} \subseteq id_A$
 (4) $R \subseteq A \times B$ is total $\Leftrightarrow id_A \subseteq R \circ R^{-1}$ } *via inverse*

Maybe one

$$(1) \Rightarrow [xRy \ \& \ xRz \Rightarrow y=z] \quad (C_1)$$

$$(a,b) \in R^{-1} \circ R, \exists y \underbrace{(a,y) \in R^{-1}} \ \& \ (y,b) \in R$$

$$\Rightarrow \exists y \ (y,a) \in R \ \& \ (y,b) \in R \quad (\text{by } C_1)$$

$$\Rightarrow a=b \Rightarrow (a,b) = (b,b) \in id_B$$

\Leftarrow if (a,b) st. $(y,a) \ \& \ (y,b) \in R \Rightarrow a=b$. THIS IS FUNCTIONALITY!

Characterizations (3)

- (1) $R \subseteq A \times B$ is functional $\Leftrightarrow R^{-1} \circ R \subseteq id_B$
 (2) $R \subseteq A \times B$ is surjective $\Leftrightarrow id_B \subseteq R^{-1} \circ R$
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 (4) $R \subseteq A \times B$ is total $\Leftrightarrow id_A \subseteq R \circ R^{-1}$ } *via inverse*

Maybe one

(3) via inverse : R is injective $\Leftrightarrow R^{-1}$ is functional

$$aR^{-1}b \ \& \ aR^{-1}c \Rightarrow b=c$$

$$bR_a \ \& \ cR_a \Rightarrow b=c \text{ \{ injectivity \}}$$

(3) \Rightarrow $R^{-1} \subseteq B \times A$ is functional
 iff $R \circ R^{-1} \subseteq id_A$

\hookrightarrow Relabel ! R exchange with R^{-1} , B with A , get (1). ALREADY PROVEN TRUE.



Closure

For property P (reflexivity, symmetricity, transitivity)...

P -closure of relation $R =$ “smallest” relation with property P

But what does smallest mean??



Closure

For property P (reflexivity, symmetricity, transitivity)...

P -closure of relation $R =$ “smallest” relation with property P

But what does smallest mean??

Intuition... start adding pairs that are missing... but is this process always ending with the same relation...

is “ P -closure” well-defined?



Closure

Easy ones: symmetric and reflexive closure

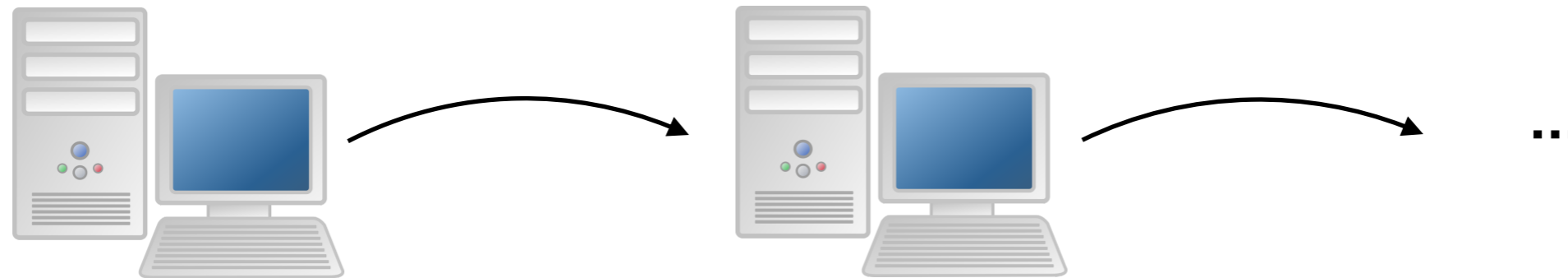
- (1) R is reflexive $\Leftrightarrow id \subseteq R$
- (2) R is symmetric $\Leftrightarrow R^{-1} \subseteq R$

Given R , its symmetric closure is $S = R \cup id_A$

Given R , its reflexive closure is $S = R \cup R^{-1}$

Closure

Transitive closure: super important



Your PC ran into a problem and needs to restart. We're just collecting some error info, and then we'll restart for you. (0% complete)

A problem has been detected and Windows has been shut down to prevent damage to your computer.

UNMOUNTABLE_BOOT_VOLUME

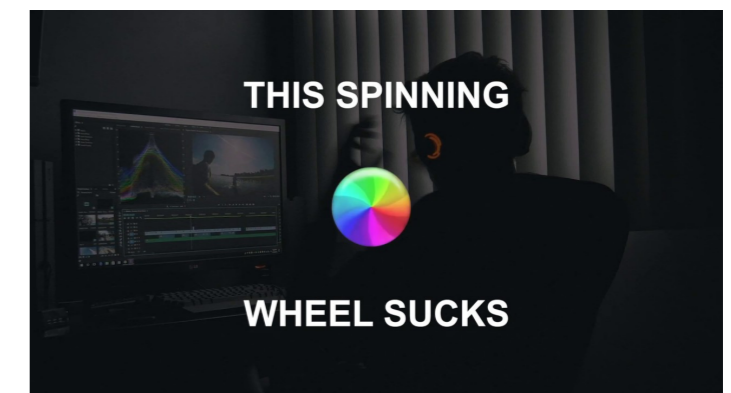
If this is the first time you've seen this error screen, restart your computer. If this screen appears again, follow these steps:

Check to make sure any new hardware or software is properly installed. If this is a new installation, ask your hardware or software manufacturer for any Windows updates you might need.

If problems continue, disable or remove any newly installed hardware or software. Disable BIOS memory options such as caching or shadowing. If you need to use Safe Mode to remove or disable components, restart your computer, press F8 to select Advanced Startup Options, and then select Safe Mode.

Technical Information:

*** STOP: 0x000000ED (0x80F128D0, 0xc000009c, 0x00000000, 0x00000000)



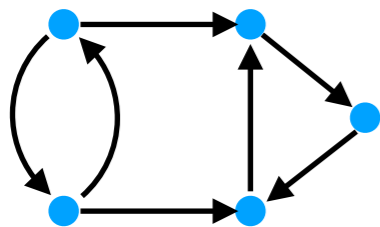
Closure: transitive

- (1) Transitive $\Leftrightarrow R \circ R \subseteq R$
 (then also) $R^{*n} \subseteq R$

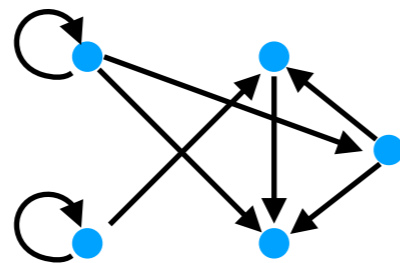
Suppose $R \circ R \not\subseteq R$

How about $R' = R \circ R \cup R$. Are we done?

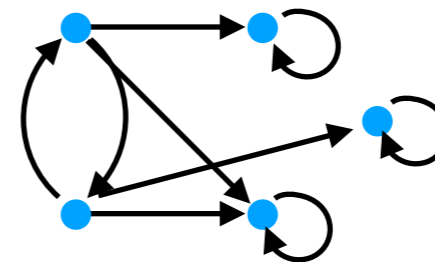
R



$R \circ R$



R^{*3}

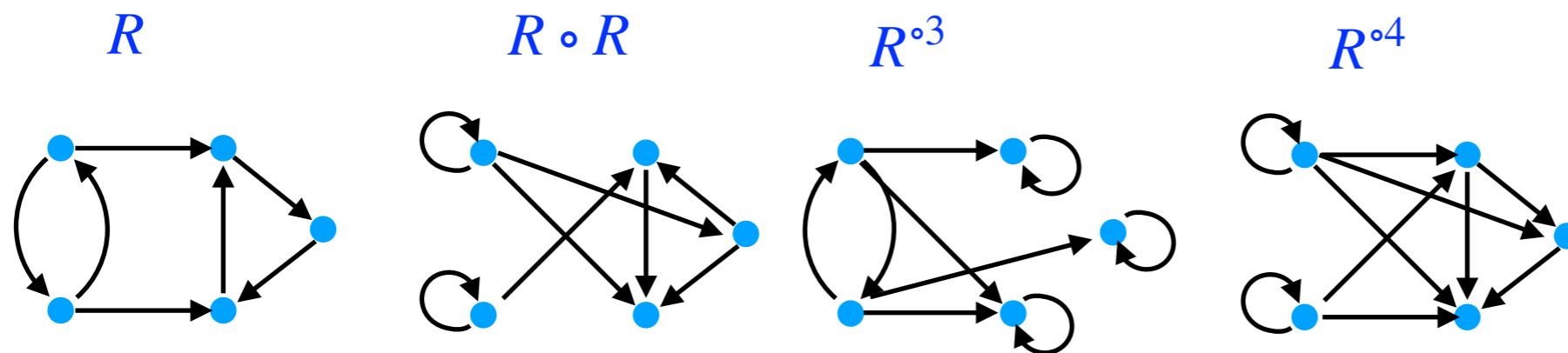


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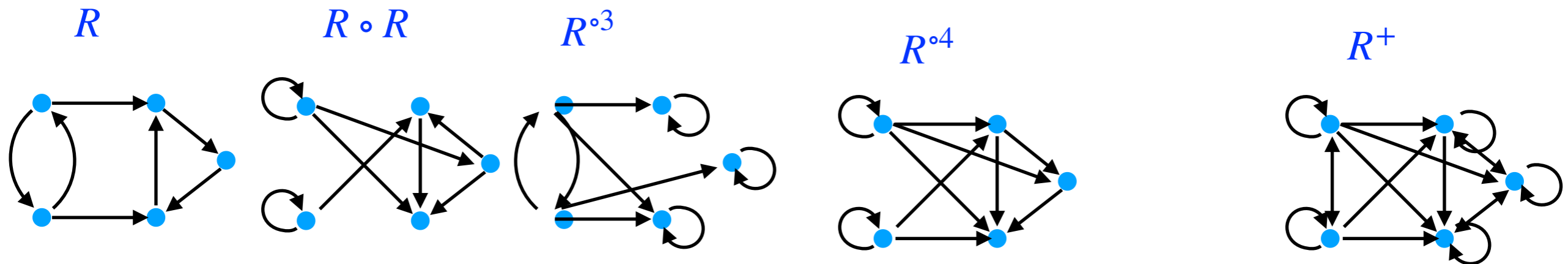


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Suppose $R \circ R \not\subseteq R$

How about $R' = R \circ R \cup R$. Are we done?





Closure: transitive

- (1) Transitive $\Leftrightarrow R \circ R \subseteq R$
(then also) $R^n \subseteq R$

The transitive closure R^+ of R is given with $R^+ = \bigcup_{k=1}^{\infty} R^{\circ k}$

Domain can be infinite...



Closure: transitive

The transitive closure R^+ of R is given with $R^+ = \bigcup_{k=1}^{\infty} R^{\circ k}$

Example: $R \in \mathbb{Z} \times \mathbb{Z}; aRb$ iff $b = a + 1$

What is R^+ ?

Work it!

$$R^+ = \bigcup R^{\circ k}$$



Closure: transitive

The transitive closure R^+ of R is given with $R^+ = \bigcup_{k=1}^{\infty} R^{\circ k}$

Example: $R \in \mathbb{Z} \times \mathbb{Z}; aRb$ iff $b = a + 1$

math jargon; iff means
“if and only if”

Work it!

$$R^+ = \bigcup R^{\circ k}$$

$$R^{\circ k} \Leftrightarrow a R^{\circ k} b \Leftrightarrow b = a + k$$

$$a R b \dots \text{it } \exists k \text{ st } b = a + k \dots$$

$$\Leftrightarrow a < b. \quad R^+ \cong <$$



Closure: transitive

Property: intersection and union of transitive relations is transitive.

Proofs! direct and by contradiction

Intersection

$$R \text{ st } R \circ R \subseteq R \quad (1)$$

$$S \text{ st } S \circ S \subseteq S \quad (2)$$

$$R \cap S = \{ (x, y) \text{ st } xRy \ \& \ xSy \}$$

ASSUME

$$(x, y) \in R \cap S \quad \& \quad (x, z) \in R \cap S.$$

$$\Rightarrow \overset{(1)}{(x, z) \in R} \quad \& \quad \overset{(2)}{(x, z) \in S}$$

$$\Rightarrow (x, z) \in R \cap S. \quad \text{done.}$$

5

Closure: transitive

Property: intersection and union of transitive relations is transitive.

Proofs! direct and by contradiction

BY CONTRADICTION: (intersection)

Assume R, S transitive & $R \cap S$ is not.
(*)

$\Rightarrow \exists \overbrace{(x,y), (y,z)} \in R \cap S$ st $(x,z) \notin R \cap S$.

\Rightarrow Either $(x,z) \notin R$ or $(x,z) \notin S$. Assume $(x,z) \notin R$.

But then $(x,y) \in R$ & $(y,z) \in R$ (by (*)) & $(x,z) \notin R$

$\Rightarrow R$ is NOT TRANSITIVE. CONTRADICTION.

IF $(x,z) \notin S$ WE GET $[(x,y) \in S \ \& \ (y,z) \in S \ \& \ (x,y) \notin S] \Rightarrow$ contradiction $\checkmark \square$