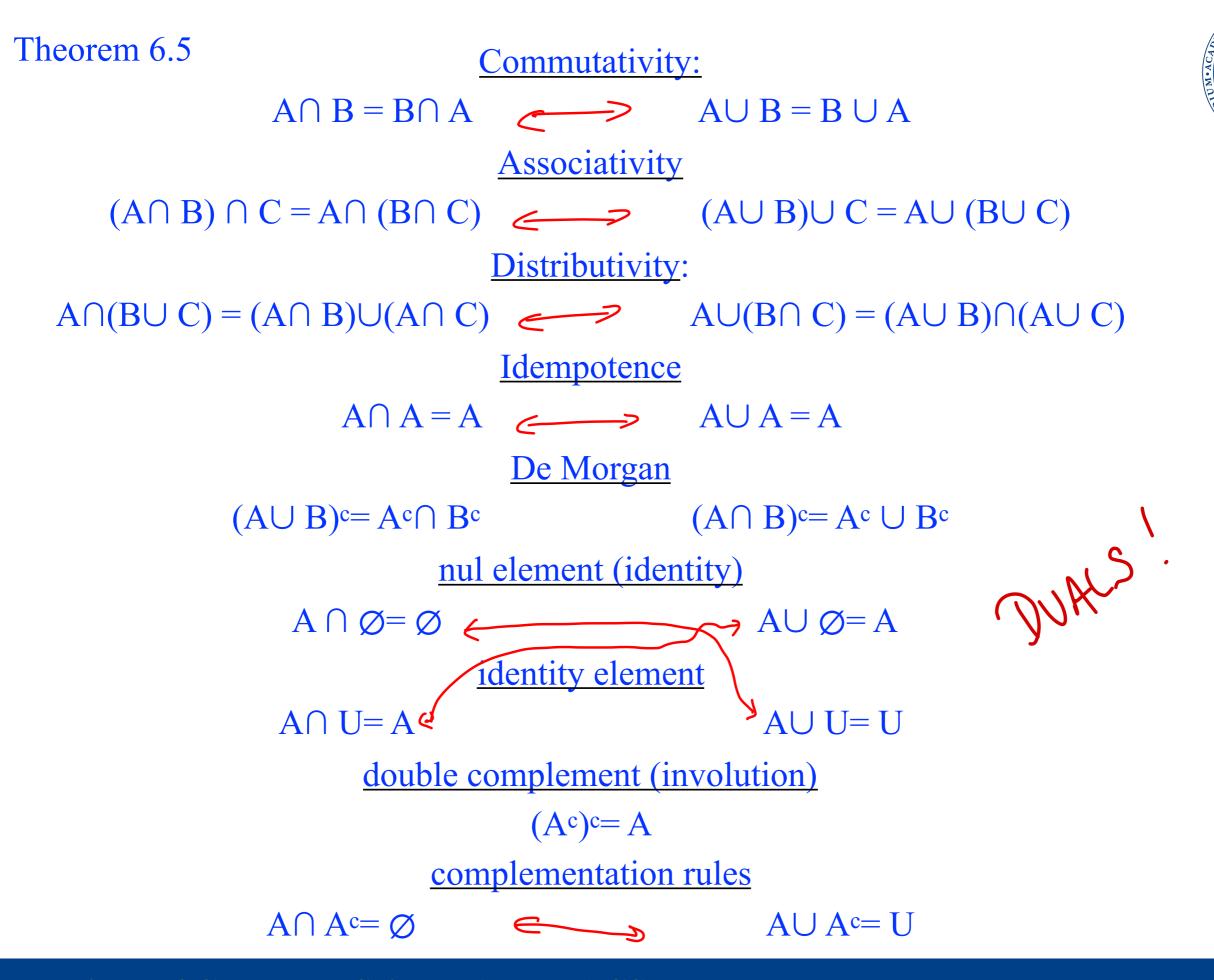


Duality of intersection and union "take any valid expression, and switch unions with intersections, and empty sets with U. It is still true."

Why, how, what?

-the duality holds for the basic rules (axioms).



Proofs in algebra of sets



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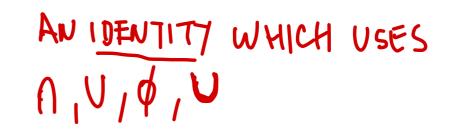
-the duality holds for the basic rules (axioms).
-any true statement <u>can be expanded to a sequential application</u> of the elementary rules.
-but we can apply duality to every step, so each step remains true after the substitute.

-so the first, and last remain true

EXAMPLE

- $A = A \cap (A \cup B)$ $A = A \cup (A \cap B)$
- $A = A U \phi \qquad A = A \cap U$
- $AU\phi = AU(BN\phi)$ $\leftarrow >$ AUU = AU(BUU)
- AU(BNØ) = (AUB) n (AUØ) AN(BUU) = (ANB) U(MU)
- $(AUA) \cup A = (UAA) \cup (ADA) \longrightarrow A \cap (BUA) = (\varphi \cup A) \cap (\partial \cup A)$

Proofs in algebra of sets





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Why, how, what?

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Digression!



How would we formally define "expressions"?

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NOUCTIVE

Digression!



Duality in mathematics

- *Duality*, generally speaking, translates concepts, theorems or mathematical structures into other concepts, theorems or structures
- in a one-to-one fashion
- often by an involution operation: if the dual of A is B, then the dual of B is A.
- Involutions sometimes have fixed points, so that the dual of A is A itself.

'In mathematical contexts, duality has numerous meanings although it is "a very pervasive and important concept in (modern) mathematics" and "an important general theme that has manifestations in almost every area of mathematics".'

Examples of using duality:

Prove, prove dual: $(U \cap A) \cup (A \cap B) = A$ **Find dual:** $A \cap (A \cap \emptyset)^c = A \cap U$

$$\frac{(unh)u(hnb)}{= A} \quad (\not 0 uh) \cap (A ub) = A$$

$$= A \quad U(hnb) \quad A \quad (A ub) = A$$

$$= A \quad U(hnb) \quad A \quad (A ub) = A$$

$$A \quad (A u b) \quad = A u \phi$$

$$A \quad (A \quad u) \quad = A u \phi$$



Sets as elements of sets



 $A \in B$

Example:

 $\{1,2\} \in \{\{1,2\},\emptyset\}$ $\emptyset \notin \{\{1,2\}\}$

(BUT. \emptyset ⊆ {{1,2}}!)



Definition: Given the set *S*, the powerset $\mathcal{P}(S)$ is the set of all subsets of *B*:

 $\mathcal{P}(S) := \{A \mid A \subseteq S\}$





Definition: Given the set *S*, the powerset $\mathscr{P}(S)$ is the set of all subsets of *B*:

 $\mathscr{P}(S) := \{A \mid A \subseteq S\}$

Powerset, aka: 2^S

Cardinality: |S|; $|\mathcal{P}(S)| = 2^{|S|}$



Definition: Given the set *S*, the powerset $\mathcal{P}(S)$ is the set of all subsets of *B*:

 $\mathscr{P}(S) := \{A \mid A \subseteq S\}$

Binary numbers, or *bitstrings*

how many?

Bitstrings		



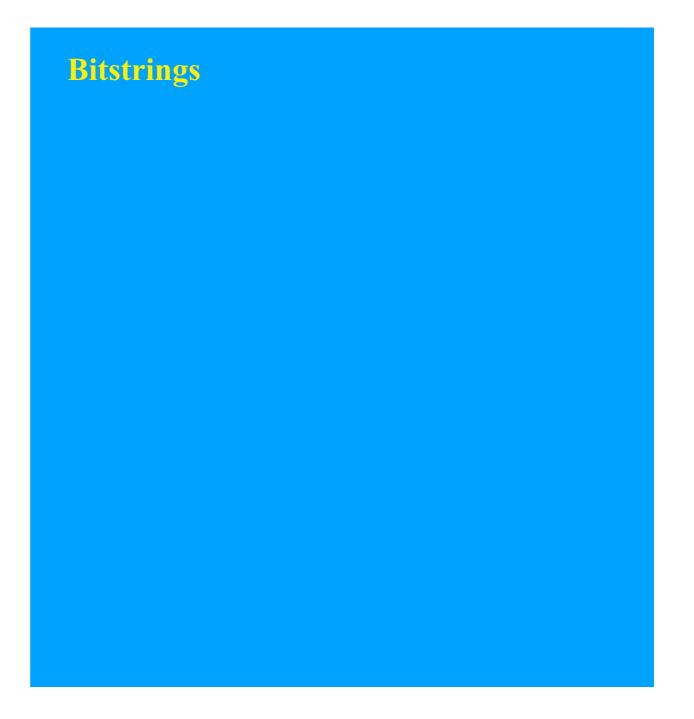
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Binary numbers, or *bitstrings*

how many?

"if I add one more bit, the count doubles"





Definition: Given the set *S*, the powerset $\mathscr{P}(S)$ is the set of all subsets of *B*:

 $\mathscr{P}(S) := \{A \mid A \subseteq S\}$

Why is

$$\begin{split} | \{ b_{n-1} b_{n-2} \dots b_0 \ | \ b_k \in 0, 1 \} | = \\ = | \mathcal{P}(\{ a_0, \dots, a_{n-1} \}) | \end{split}$$

EACH BITSTRING SPECIFIES EXACTLY ONE OVESET & F EACH SUBSET SPECIFIES EXACTLY ONE BITSTRING

Bitstrings v.s. powersets $\alpha_{\eta} \underline{\mathcal{Q}}$ $L > \{a_1, a_1, \dots\}$



Definition: Given the set *S*, the powerset $\mathscr{P}(S)$ is the set of all subsets of *B*:

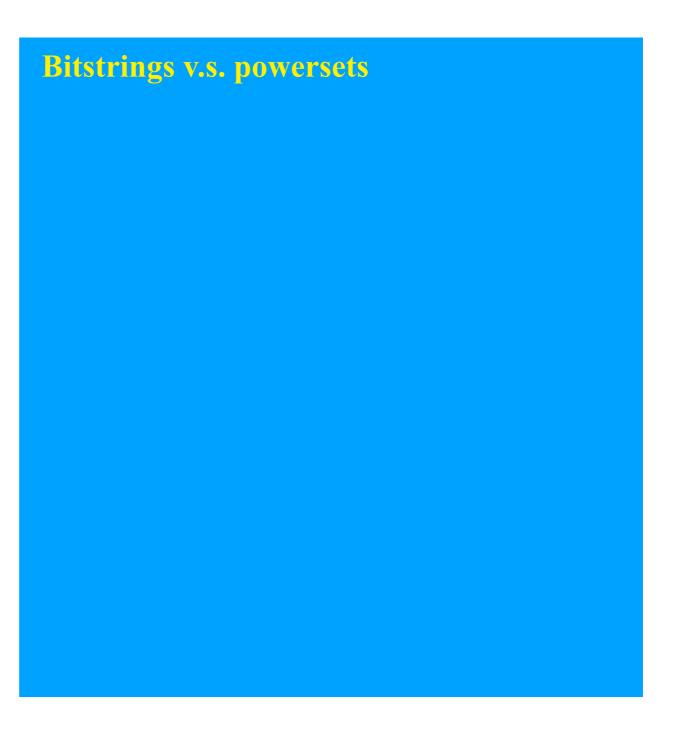
 $\mathscr{P}(S) := \{A \mid A \subseteq S\}$

Why is

$$|\{b_{n-1}b_{n-2}...b_0 | b_k \in 0,1\}| = |\mathcal{P}(\{a_0, ..., a_{n-1}\})|$$

bitstrings =

= H Subsets







- **Definition.** An <u>alphabet</u> is a non-empty, finite set of <u>letters</u>.
- Definition. A string (word) over the alphabet Σ is an ordered set of letters from the alphabet Σ
- Notation: Σ^* set of all strings from the *alphabet* Σ
- Definition. A <u>language</u> over the alphabet Σ is a set of strings (words) over Σ
- Notation: $\mathscr{P}(\Sigma^*)$ all the languages over Σ





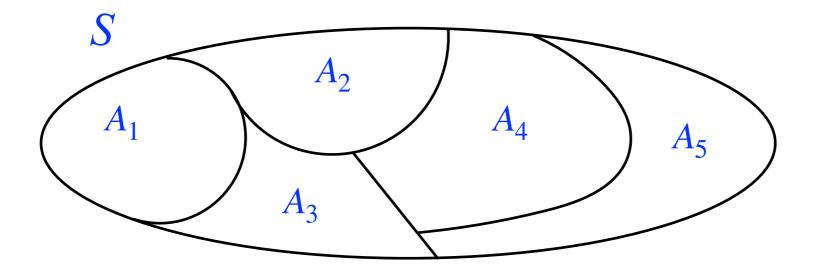
Definition: Given the set *S*, family of sets $P_s = \{A_1, A_2, A_3...\}$ is called a (countable) partition of *S* if

• $\emptyset \notin P_s$ • $S \in A_1 \cup A_2 \cup \dots = \bigcup_{i \in I} A_i$

 $[P_s \text{ covers } S]$

• For all $i, j, A_i \cap A_j = \emptyset$

[elements of P_s are pairwise disjoint]







Definition: Given the set *S*, family of sets $P_s = \{A_1, A_2, A_3...\}$ is called a (countable) partition of *S* if: *a*) no A_k is empty; *b*) they cover *P c*) pairwise disjoint

Examples:





- Defined relative to the universe \mathbb{Z}
- given $k \in \mathbb{N}$, and $l \in \mathbb{Z}$, the congruence class of l modulo k is:

 $\bar{l} = \{ \dots, l + nk, \dots l - 2k, l - k, l, l + k, l + 2k, \dots, l + nk, \dots \}$

• **Example**: congruence classes of 0,1,2 modulo 7:

 $\bar{0} = \{\dots - 14, -7, 0, 7, 14\dots\}$ $\bar{1} = \{\dots - 13, -6, 1, 8, 15\dots\}$ $\bar{2} = \{\dots - 12, -5, 2, 9, 16\dots\}$



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• **Example**: congruence classes of 0,1,2 modulo 7:

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- NB: congruence classes are sets. Recall when two sets are equal.
- work out a few other congruence classes modulo 7 (of some other number)
- for a given *k*, are there infinitely many classes? how many?

VERY! SUPPLEMENTAL



• **Example**: congruence classes of 0,1,2 modulo 7:

 $\bar{0} = \{\dots - 14, -7, 0, 7, 14\dots\}$ $\bar{1} = \{\dots - 13, -6, 1, 8, 15\dots\}$ $\bar{2} = \{\dots - 12, -5, 2, 9, 16\dots\}\dots$

• Note: for mod 7: $\overline{0} = \overline{7}$; $\overline{1} = \overline{8}$; for mod k: $\overline{l} = \overline{k+l}$; $\overline{l} = \overline{k+ml}$

Lets work this out:

• **Example**: congruence classes of 0,1,2 modulo 7:

 $\bar{0} = \{ \dots - 14, -7, 0, 7, 14 \dots \}$ $\bar{1} = \{ \dots - 13, -6, 1, 8, 15 \dots \}$ $\bar{2} = \{ \dots - 12, -5, 2, 9, 16 \dots \} \dots$ VERY! SUPPLEMENTAL



- Partition of \mathbb{Z} !
- $R_7 = \{\bar{0}, \bar{1}, ..., \bar{6}\}$ (check properties!)

Lets work this out:

Important partitions: non-trivial infinite partition

 $A_0 = \{1,3,5,7...\}$ $A_1 = \{2,6,10,14,...\}$ $A_2 = \{4,12,20,28,...\}$



Important partitions: non-trivial infinite partition



$$A_0 = \{1,3,5,7...\}$$
$$A_1 = \{2,6,10,14,...\}$$
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$$A_k = \{2^k \cdot 1, 2^k \cdot 3, 2^k \cdot 5, 2^k \cdot 7, \dots\}$$

Important partitions: non-trivial infinite partition



$$A_0 = \{1,3,5,7...\}$$
$$A_1 = \{2,6,10,14,...\}$$
$$A_2 = \{4,12,20,28,...\}$$

$$A_k = \{2^k \cdot 1, 2^k \cdot 3, 2^k \cdot 5, 2^k \cdot 7, \dots\}$$

Claim: this is a partition of the natural numbers. Prove!

Intermezzo! Binary numbers

Decimal num

Auxiliary

$$\begin{aligned} \text{hbers (base 10)} \quad & d_{n-1}d_{n-2}...d_{0}, \text{ with } d_{k} \in \{0,...,9\} \\ & 123 = 100 + 20 + 3 = 1 \cdot 10^{2} + 2 \cdot 10^{2} + 3 \cdot 10^{2} \\ & d_{n-1}d_{n-2}...d_{0} = d_{n-1} \cdot \underline{10^{n-1}} + \underline{d_{n-2}} \cdot 10^{n-3} + \cdot + \underline{d_{0}} \cdot \underline{10^{0}} \\ & b_{n-1}b_{n-2}...b_{0}, \text{ with } b_{k} \in \{0,1\} \\ & b_{n-1}b_{n-2} \times 2^{n-4} + b_{n-2} \times 2^{n-2} - b_{0} \cdot 2^{n-4} \\ & b_{n-4} \times 2^{n-4} + b_{n-2} \times 2^{n-2} - b_{0} \cdot 2^{n-4} \\ & b_{n-4} \times 2^{n-4} + b_{n-2} \times 2^{n-4} + b_{n-2} \times 2^{n-4} + b_{n-2} \\ & 0 \\ & 0 \\ & 1 \\ &$$





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