

# Logica (I&E)

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<http://liacs.leidenuniv.nl/~vlietrvan1/logica/>

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2. Predicate logic

2.3. Proof theory of predicate logic

*De hoogste opgave van het menselijk kennen is: te begrijpen  
dat hij niet begrijpen kan.*

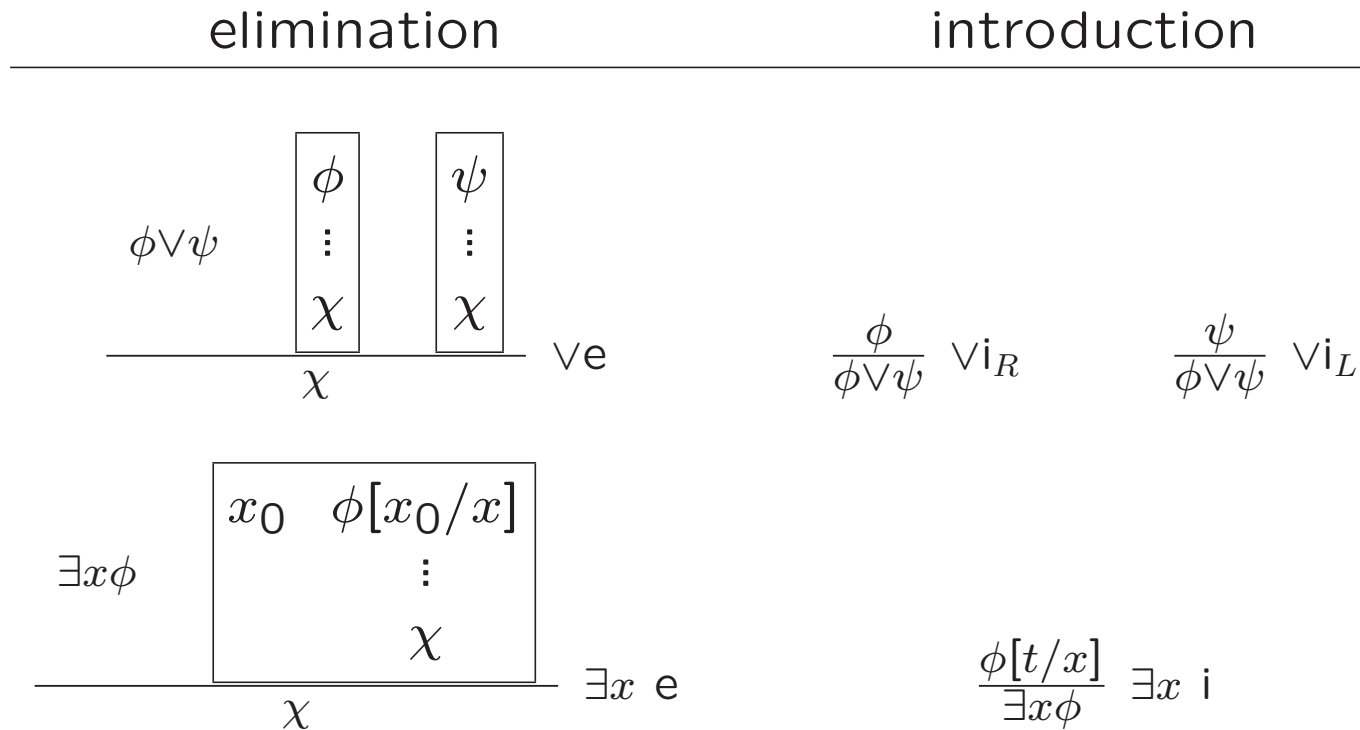
A slide from lecture 10:

# Analogy $\wedge$ and $\forall$

elimination	introduction
$\frac{\phi \wedge \psi}{\phi} \wedge e_R$	$\frac{\phi \quad \psi}{\phi \wedge \psi} \wedge i$
$\frac{\forall x \phi}{\phi[t/x]} \forall x e$	<div style="display: flex; align-items: center; justify-content: center;"> <div style="border: 1px solid black; padding: 10px; margin-right: 10px;"> <math>x_0</math>  <math>\vdots</math>  <math>\phi[x_0/x]</math> </div> <div style="text-align: center;"> <math display="block">\frac{}{\forall x \phi} \forall x i</math> </div> </div>

A slide from lecture 10:

# Analogy $\forall$ and $\exists$



**Example.**

$$\exists xP(x), \forall x\forall y(P(x) \rightarrow Q(y)) \vdash \dots$$

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$$\exists xP(x), \forall x\forall y(P(x) \rightarrow Q(y)) \vdash \forall yQ(y)$$

Proof...

# Why fresh variables

Example.

$$\exists xP(x), \forall x(P(x) \rightarrow Q(x)) \vdash \dots$$

# Why fresh variables

Example.

$$\exists xP(x), \forall x(P(x) \rightarrow Q(x)) \vdash \forall yQ(y)$$

'Proof'...

## 2.3.2. Quantifier equivalences

Is

$$\forall x \forall y \phi \dashv\vdash \forall y \forall x \phi$$

valid?

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Is

$$(\forall x \phi) \wedge (\forall x \psi) \dashv\vdash \forall x (\phi \wedge \psi)$$

valid?

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Is

$$(\forall x \phi) \wedge \psi \dashv\vdash \forall x (\phi \wedge \psi)$$

valid?



Is

$$\exists x(\phi \rightarrow \psi) \dashv\vdash \forall x\phi \rightarrow \psi$$

valid?

**Example 2.12.**

Not all birds can fly.

$$\neg \forall x (B(x) \rightarrow F(x))$$

$$\exists x (B(x) \wedge \neg F(x))$$

**Theorem 2.13.** Let  $\phi$  and  $\psi$  be formulas of predicate logic.

$$1.(a) \quad \neg\forall x\phi \dashv\vdash \exists x\neg\phi$$

$$(b) \quad \neg\exists x\phi \dashv\vdash \forall x\neg\phi$$

Proof 1(a)...

Left-to-right:

$$\begin{aligned} \neg(p_1 \wedge p_2) &\vdash \neg p_1 \vee \neg p_2 \\ \neg\forall xP(x) &\vdash \exists x\neg P(x) \\ \neg\forall x\phi &\vdash \exists x\neg\phi \end{aligned}$$

Right-to-left...

**Theorem 2.13.** Let  $\phi$  and  $\psi$  be formulas of predicate logic.

2. Assuming that  $x$  is not free in  $\psi$ :

$$(a) \quad \forall x\phi \wedge \psi \dashv\vdash \forall x(\phi \wedge \psi)$$

$$(b) \quad \forall x\phi \vee \psi \dashv\vdash \forall x(\phi \vee \psi)$$

(c) ... (h)

Proof 2(a)...

**Theorem 2.13.** Let  $\phi$  and  $\psi$  be formulas of predicate logic.

2. Assuming that  $x$  is not free in  $\psi$ :

$$(a) \quad \forall x(\phi \wedge \psi) \vdash \forall x\phi \wedge \psi$$

Proof...

**Theorem 2.13.** Let  $\phi$  and  $\psi$  be formulas of predicate logic.

$$3.(a) \quad \forall x\phi \wedge \forall x\psi \dashv\vdash \forall x(\phi \wedge \psi)$$

$$(b) \quad \exists x\phi \vee \exists x\psi \dashv\vdash \exists x(\phi \vee \psi)$$

Proof 3(b)...

**Theorem 2.13.** Let  $\phi$  and  $\psi$  be formulas of predicate logic.

$$4.(a) \quad \forall x \forall y \phi \dashv\vdash \forall y \forall x \phi$$

$$(b) \quad \exists x \exists y \phi \dashv\vdash \exists y \exists x \phi$$

Proof 4(b)...

Study this proof yourself.